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**COMMON MEMORY
FOR THE PERSONAL COMPUTER**

August 11, 1988

By:

Siegfried Rybczynski



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Gaithersburg, Maryland

Common Memory for the Personal Computer

Siegfried Rybczynski

August 11, 1988

DISCLAIMER

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I. INTRODUCTION

1. ORGANIZATION OF THIS DOCUMENT

Section I serves as an introduction to the concept of common memory. It identifies the purpose of the accompanying software collection, and lists the personal computer (PC) configuration prerequisites of the PC common memory software library.

Section II provides an overview of the common memory architecture defined by the Automated Manufacturing Research Facility (AMRF) of the National Bureau of Standards. A short personal opinion of the drawbacks to common memory is provided at the end of the section.

Section III specifies the PC common memory architecture. The AMRF architecture allows for a range of solutions to meet specific management and access coordination problems. This section identifies which solutions (or alternatives) were implemented in the PC common memory architecture. Some minor extensions to the original architecture are presented.

Section IV is the programmer's reference section. It lists all PC common memory function calls and details their argument lists and return values. A functional description of each function is provided.

The various appendices provide detailed support information for the use of the PC common memory program library. Included are data structure definitions, a sample program and associated output, and the complete source listing of the PC common memory library.

2. HISTORY OF COMMON MEMORY

The Automated Manufacturing Research Facility (AMRF) at the National Bureau of Standards (NBS) is using an architecture called "common memory" for interprocess communication.

The AMRF common memory architecture originated as a consequence of an application that required real-time data reduction with concurrent robot control [5]. To meet that need, a multiprocessor configuration that included a physical common memory entirely contained within a single Multibus chassis was used. Each of the processors was a single-board computer. The memory designated for common use by all processors was resident on a separate board and had an address range that mapped into the address space of each of the processors. Hence, local common memory was defined to be a contiguous area of physical memory accessible to two or more distinct processes within a single computer system (Figure I-1).

The AMRF began work on an automated factory in 1981 [6]. The AMRF automated factory concept held that the old idea of a single huge computer controlling all machines in the factory was too inflexible. Instead, computer processes such as control programs would run on many different computers, of all sizes and models, and might possibly be located in different buildings.

This extended the local common memory concept of the robot control system in many ways. A major extension was the linking and coordination of local common memories using network services. That is, processes which had to communicate with each other were often in separate backplanes and used different operating systems. A single physical common memory was no longer possible or practical, so each computer system had to have its own local common memory. These common memories were connected using network services that were transparent to the application process. The linked local common memories established a global common memory (Figure I-2). The information contained therein was considered to represent a global memory-resident database.

A discussion of the AMRF network architecture is beyond the scope of this document. Detailed information about that architecture is available in reference [7].

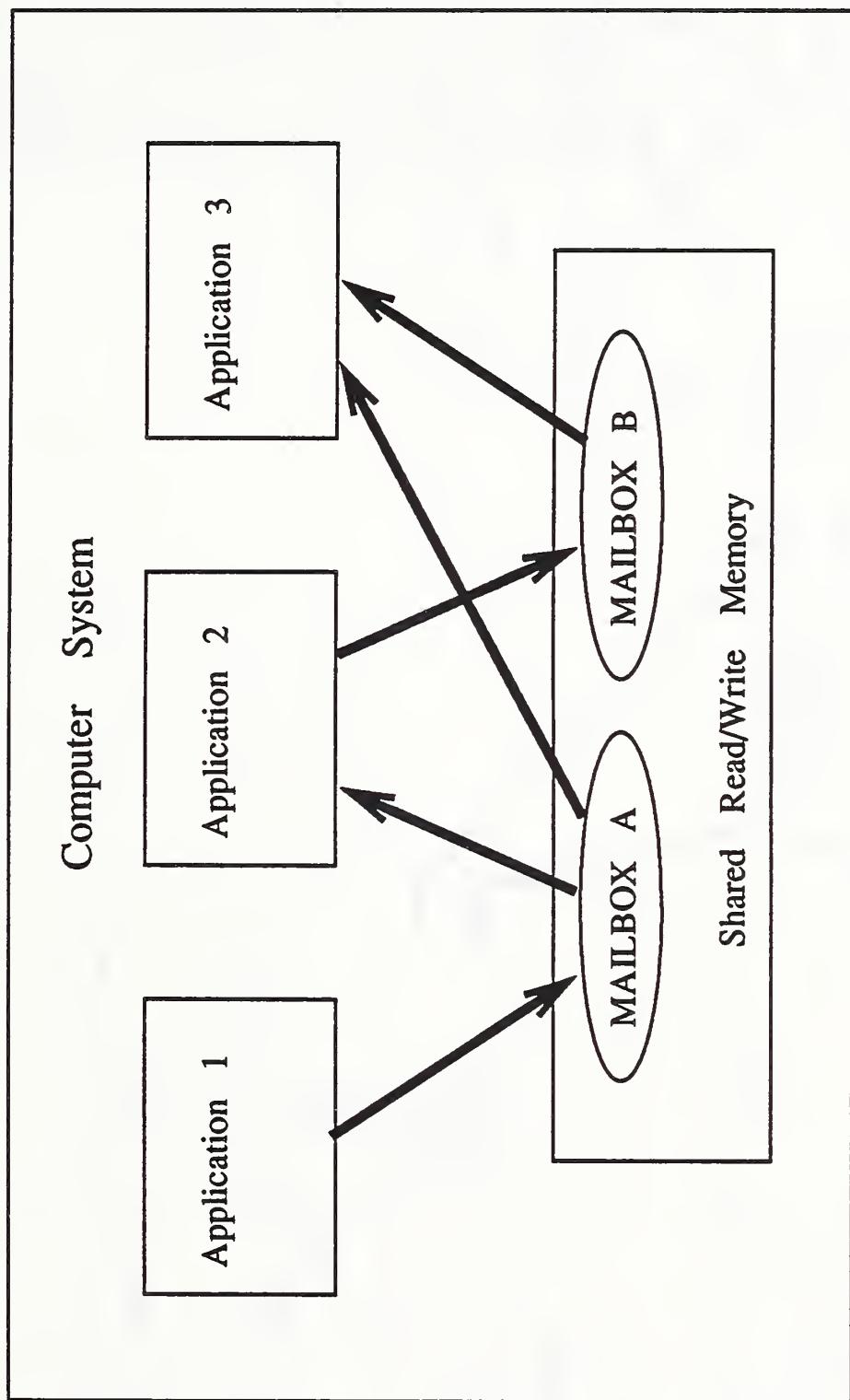


Figure I-1. Local Common Memory

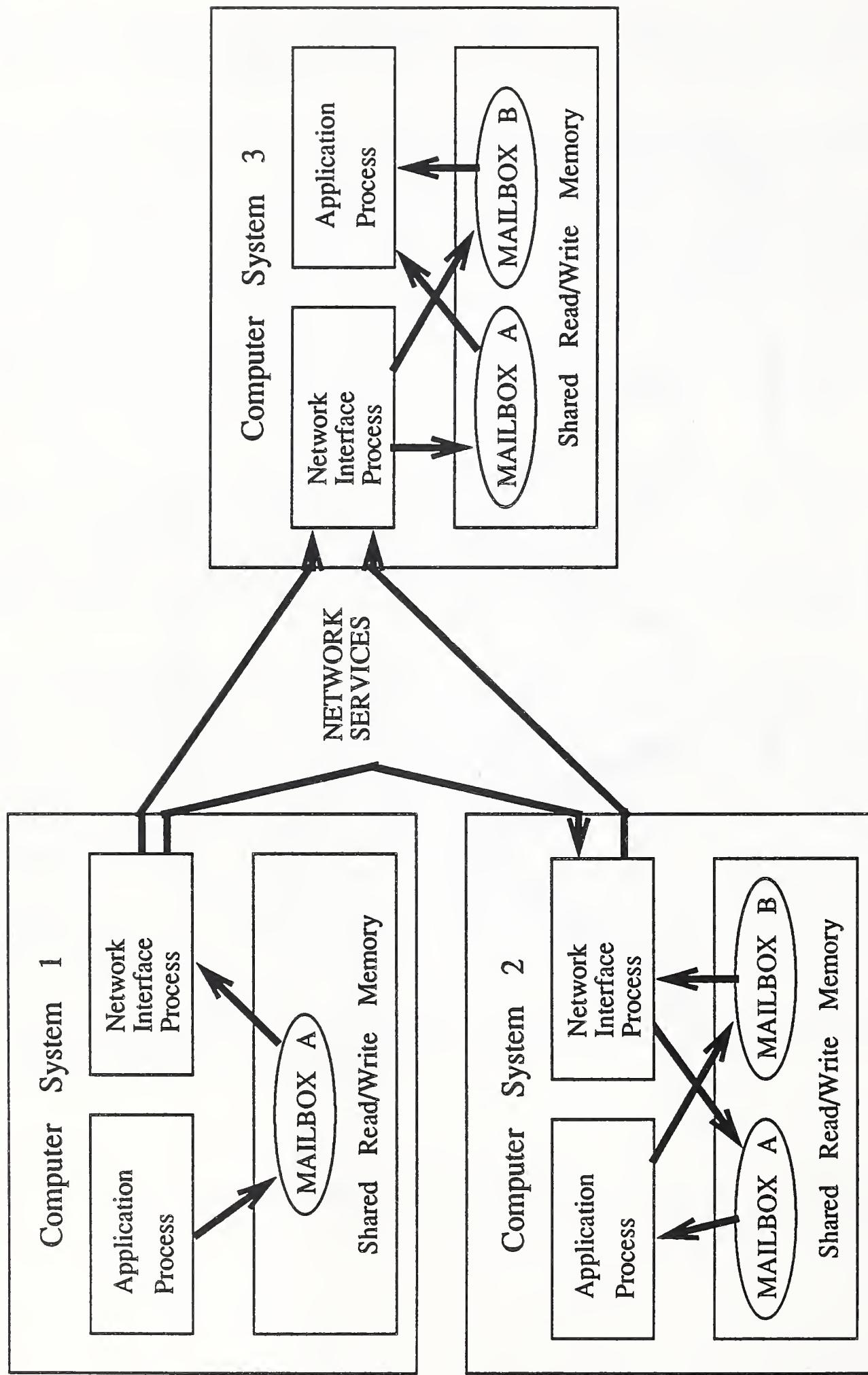


Figure I-2. Global Common Memory

3. PURPOSE OF THIS SOFTWARE COLLECTION

The intent of this software is to provide a local common memory environment for the IBM PC and compatibles. This generic group of personal computers will hereafter be referred to as PC's.

Given the increasing speed (e.g., 25 MHz 80386 personal computers) and capability (32 Mbytes of addressable space, multitasking operating systems, etc) of the PC and its rapidly decreasing price, it is feasible and desirable to consider the PC for implementation as a networked real-time controller. Providing the common memory for the PC is the first step in this process.

4. COMPUTER CONFIGURATION REQUIREMENTS

Computer configuration requirements are:

1. An IBM PC (or compatible) with at least 256 Kbytes of memory and DOS operating system.
2. The C programming language. Although the Turbo C (version 1.0) distribution was used during PC common memory development, Turbo C dependencies were studiously avoided -- with one exception, identified in Section IV.1.4. The use of C enabled the creation of object module libraries that can be used during the LINK process to incorporate common memory into user processes written in languages such as assembler, Prolog, FORTRAN, C and others.

II. OVERVIEW OF THE COMMON MEMORY ARCHITECTURE

The AMRF architecture of common memory and network communications is fully described in [7]. The following subsections discuss components of that architecture that are important to the understanding of PC common memory.

1. COMMON MEMORY COMPONENTS

1.1. Mailboxes and Mailgrams

All interprocess communication is accomplished through a mechanism called "mailboxes". Mailboxes are logical storage areas where messages (called "mailgrams") are placed by sender processes and picked up by receiver processes. From the point of view of the sender and receiver processes, the location of the correspondents does not affect their communication.

Mailboxes reside in a special area of memory, designated as local common memory. A common memory manager is responsible for assigning and managing the local common memory area.

Local common memories can be combined into a global common memory by implementing a unique common memory client process that interfaces with another similar process on a remote host for the purpose of exchanging common memory information. The physical connections can be point-to-point or utilize various local area network media, such as Ethernet or broadband. (Figure I-2.)

The user interface to local common memory is discussed below.

1.1.1. Mailgram Format

By convention, mailgrams are placed into a mailbox beginning at the lowest memory location allocated for the mailbox and continue occupying bytes until either the entire mailgram is in the mailbox or until all space allocated for the mailbox has been filled. Except for adherence to this convention, there is no standard format for a mailgram. There are, however, standard information units which must be associated with each mailbox. They are:

1. Length of the current mailgram in the mailbox,
2. Indication of a change in mailbox contents (discussed in Section II.2.3.), and
3. Some type of mailbox access control mechanism (discussed in Section II.2.)

The current level of the AMRF common memory architecture requires that these information units must be present either in the

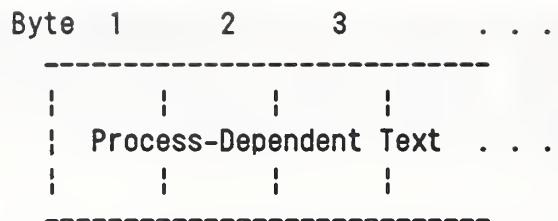
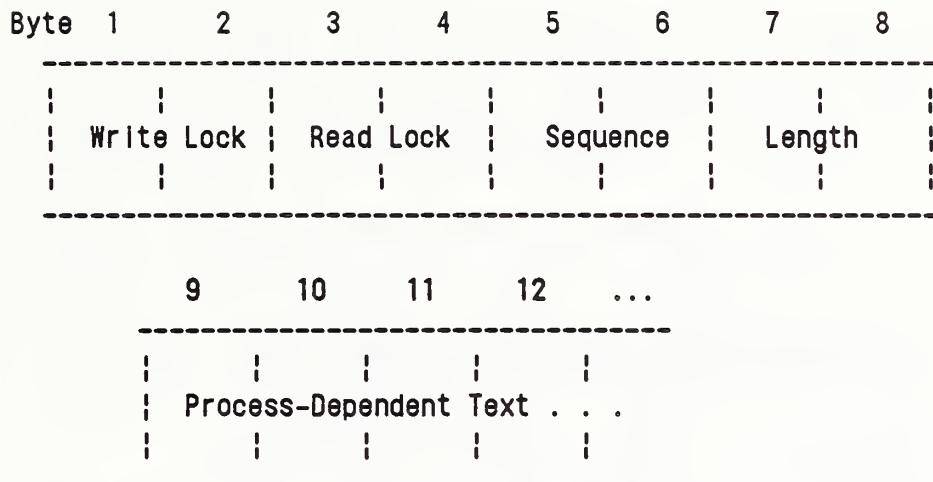


Figure II-1. Generic Mailbox Structure

mailgram or be maintained with the mailbox by the common memory manager. In general, if there is a common memory manager available in a particular common memory implementation, then the mailbox structure is shown in Figure II-1.

However, when no common memory manager is present and these entities are expressed in the mailbox area itself, then the mailbox has the structure shown in Figure II-2.



Write Lock is a semaphore indicating current writer activity.
(i.e., if the write lock is ON, then the mailbox is being written, and should not be read)

Read Lock is a semaphore indicating current reader activity.
(i.e., if the read lock is ON, then the mailbox is being read, and should not be updated)

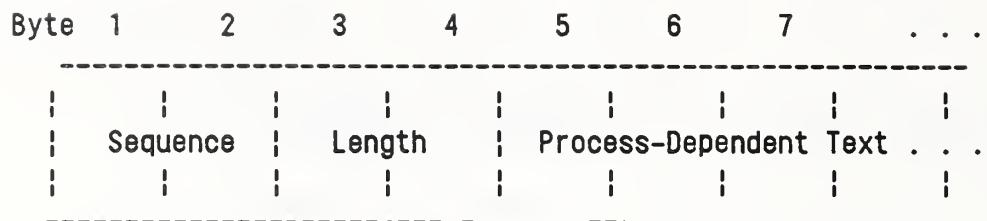
Sequence is a sequence number attached to the mailgram in the particular mailbox. Every time the text of the mailgram is changed, the sequence number is incremented. The update can be detected by examining only the sequence field.

Length is the length of the mailgram in bytes.

Text is the information portion of the mailgram that is defined entirely by the communicating processes.

Figure II-2. Special Mailbox Structure Used When A Common Memory Manager Is Not Present

With reference to Figure II-2, the read and write locks are considered to be part of the mailbox and not the mailgram. That is, when the mailgram is read or written, the lock bytes are manipulated in order to assure the integrity of the mailbox access, but are not transferred to or from a user data storage area. Furthermore, if a network interface process is to transfer the mailgram to another network interface process located on a different computer, the lock bytes are not transported.



Sequence Is a sequence number attached to the mailgram in the particular mailbox. Every time the text of the mailgram is changed, the sequence number is incremented so that the change can be detected by examining only the sequence field.

Length Is the length of the mailgram in bytes.

Text Is the information portion of the mailgram that is defined entirely by the communicating processes.

Figure II-3. General Mailgram Format Used In A Mixed Common Memory Environment

In a mixed global common memory environment, where some systems have a common memory manager and others do not, the mailgram is assigned the specific format depicted in Figure II-3. The system with the common memory manager receives a mailgram with the sequence number and the length byte included in what it considers to be "process-dependent information". It is the responsibility of the application reading from the mailbox to know the mailgram format being used.

1.1.2. Mailbox and Mailgram Properties

- (1) The mailbox must be created (or declared) before mailgrams can be deposited into it or an attempt can be made to read the mailbox contents.
- (2) Every mailbox has a unique global name. The name is assigned to the mailbox at the time it is created and uniquely identifies the mailbox to all processes participating in the global common memory. Remote systems desiring a copy of the mailbox contents must have a mailbox with the same name declared in their local common memory. Network services perform the mailgram transfers [7].

- (3) Every mailbox contains an initial value assigned by the creator when it is created. In some systems, this is a standard value (e.g., all zeros); in other systems, this value defaults to the contents of memory at the time of creation.
- (4) Every mailbox contains exactly one mailgram at any given time. A mailgram stays in the mailbox, no matter how often it is read, until a new mailgram arrives for that mailbox. The new mailgram replaces the old one on arrival, whether or not the old mailgram has ever been read.
- (5) The mailbox writer decides when to replace the mailgram. This may be performed independent of external information or may be influenced by "flow control" factors.
(Section II.2.)
- (6) In general, only one process is authorized to write into a mailbox at a time. In the case where more than one process may write into a specific mailbox, implicit or explicit "flow control" must be implemented.
(Section II.2.)
- (7) Any number of reader processes can retrieve the current mailgram in a mailbox.
- (8) Any reader process can pick up the same mailgram several times if the writer does not change it in the interim. Likewise, any reader process may miss several mailgrams if the writer changes the mailgram more often than the reader picks it up. When it is important to assure that a particular recipient has read the mailgram before a new one is issued, the writer and reader must agree to a "flow control" protocol. (Section II.2.)

The mailbox management mechanism guarantees that a new mailgram will be distinct from its predecessors. However, the mailbox management mechanism does not guarantee that any particular receiver will have picked up a mailgram before it is replaced. If it is necessary to assure that a particular receiver has read the mailgram before it is replaced, the sender and that receiver must agree to a protocol by which the sender refrains from replacing the mailgram until it has an indication that the receiver has read it.

- (9) Every mailbox has a fixed size which is defined when the mailbox is created. There is no predefined maximum on mailbox size. There may be a maximum mailbox size for individual systems as a consequence of hardware or

software limitations. Any given mailbox must be large enough to contain the largest mailgram agreed upon between the sender and receiver(s).

(10) Mailgrams can be of variable length. Each mailgram has associated with it information on how long it is. A mailgram may never be longer than the mailbox in which it is placed. If necessary, the mailgram will be truncated before it is transferred into the destination mailbox by the common memory service routines.

1.2. Common Memory Manager

Depending on the capabilities of the host operating system within any computer (or single-board computer in a Multibus chassis), it is possible to have an active common memory manager agent. However, an active common memory manager is not absolutely required. A passive manager is also possible.

1.2.1. Active Common Memory Manager

An active common manager is one that actively controls mailbox access for each common memory client using either of the following methods.

(1) Pass a token among the clients. The token might be passed sequentially from client to client. Alternately, the token could be held by the common memory manager; when a client requests common memory access, he is given the token. The token is returned to the manager when the client has completed common memory access.

It is possible for the token to get lost. That is, if a client that has been issued the token never returns it (due to a program abort or unexpected endless loop), all other clients are prohibited from further access. In order to minimize the consequences of such a deadlock situation, a token regeneration scheme must then be developed to recognize when the token has been unequivocally lost and a new token must be generated.

(2) Present a single interface to all common memory clients: the manager's access interface. The manager then coordinates, internally, the individual accesses of the clients. This is somewhat similar to token passing, where the token only passes between the manager and a client, and never between two clients. However, it does give the manager the flexibility of coordinating more than one concurrent access.

1.2.2. Passive Common Memory Manager

In the case of a passive common memory manager, there is no token or single interface. Instead, there is an access convention that is adhered to by all common memory clients. For example, the robot controller referenced previously (Section I.1.1.) uses a scheme involving repetitive common memory access cycles [8]. In this instance, there are three distinct divisions of the cycle:

- (1) READ division: During the READ division all processes (each of which exists on its own single-board computer) compete for bus access in order to READ from common memory mailboxes. No WRITE actions are performed during this time period. The duration of this division is fixed. All READ accesses that are not completed during this time period are postponed until the next READ division.
- (2) PROCESS division: During this division, all data reduction, equipment control and data acquisition occurs. No READ or WRITE access to the mailboxes is performed. Its duration is fixed.
- (3) WRITE division: During the WRITE division, all processes once again compete for bus access in order to WRITE to common memory mailboxes. NO READ actions are performed during this time period. The duration of this division is fixed. All WRITE accesses that are not completed during this time period are postponed until the next WRITE division.

2. COORDINATING COMMON MEMORY ACCESS

Common memory environments can be susceptible to several problems related to coordinating access to these areas. The following subsections identify the potential problems and some possible solutions. Section III of this document specifies the solutions implemented in the PC common memory architecture.

2.1. Read While Write is Active

A process may be attempting to read information from a common memory mailbox at the same time that a second process is attempting to update that mailbox (or vice-versa). Consequently, the reading process may get inconsistent information (e.g., the current value of field A and the former value of field B).

To avoid this, one could:

- (1) use a semaphore for each common memory buffer area (a mechanism that supports single-process access to the buffer). Some processors provide atomic "test and set" operations which can be used as hardware semaphores. Unfortunately, the PC does not. Software semaphores, using Dekker's algorithm [4], for example, can be extended to provide mutual exclusion between any number of processes.
- (2) define a regular, recurring real-time interval and divide it into a write-only period and a read-only period. Any process not prepared to perform a write operation during the write-only period would have to wait for the next write-only period. The same restriction holds for read-only periods.
- (3) pass a token among participating processes. The process that has the token can perform any read or write operation it wants. Fixed length or varying length time quanta can be employed. Token passing has an unfortunate drawback: if the process with the token halts (or appears to do so), passing of the token becomes impossible and all access to common memory is barred. In a fixed length time quantum implementation, the token can be reissued by some governing process after the expiration of the time quantum (plus some extra "safety margin"); in a varying length time quantum implementation, the recovery algorithm is much less obvious.

(4) utilize a hardware architecture that does not support interrupt processing. Once a processor has control of the bus (and consequent access to common memory), no other processor can interrupt. This assures that overlapped access does not occur.

2.2. Update Frequency Exceeds Read Frequency

A process may update the common memory area more often than a reading process is able to retrieve the information.

This may only be an "application-specific" problem. That is, if the reader process only wants the "current" information (as from a temperature sensor, for example), then the fact that any amount of older information may have been missed is a moot point. However, if it becomes important that the reader process have access to each information set before it gets updated, then some form of "flow-control" must be used.

For example, if the information set in a common memory mailbox includes a unique identifier (a time stamp or sequence number), then flow control could be implemented by defining a second mailbox in the common memory area into which the reader process could echo the unique identifier. When the writer of the original information sees the echoed identifier in this second common memory mailbox, it knows that it can proceed with the next update.

This method of flow control is feasible when the reader process can consume the information as fast as it is produced. However, if the reader process is too slow, it can have negative ramifications for the information writer. For example, if a temperature sensor at a nuclear power plant is attempting to report rapidly rising temperatures but is prohibited from reporting the current temperature because the mailgram reader has not acknowledged the previous temperature, undesirable side-effects can result.

2.3. Read Frequency Exceeds Update Frequency

A process may read the data in the common memory area more often than a writer process updates it. This can result in "old" information unintentionally being considered "new" information.

In the case where the information happens to be a command such as "hit nail on head with hammer", an undesirable number of duplicate executions could be performed.

A possible solution is to identify new information whenever it is placed into the common memory buffer by implementing a flag field

within the mailbox. This flag field could take the form of a sequence number that gets incremented with each update of the mailbox or a time stamp that identifies when the information was placed into the mailbox. In each case, the reader process is looking for a change in the flag field to indicate that mailbox contents have been updated.

An alternative method is for the common memory manager to maintain a list containing the names of mailboxes that have been updated. A separate list is maintained for each common memory application. The application can then be given the update list upon request. By reading only the changed mailboxes, the application can minimize unproductive time spent examining unchanged mailboxes. However, this method is only available for systems with an active common memory manager.

2.4. Multiple Readers of a Common Mailbox

In the case where a single mailbox is being accessed by multiple readers, if it is important that each of the readers have the opportunity to retrieve the mailgram before it is overwritten, then a more elaborate form of flow control must be implemented.

One solution is to share a single "flow control" mailbox between all the readers. Each reader sets a specific "flag" in the mailbox indicating he has retrieved the message. When all flags have been set, the shared-read mailbox contents can be overwritten. This "solution" immediately introduces another problem: multiple writers to a single mailbox. (See Section II.2.5.)

A simpler, more reliable solution is to assign each reader process its own flow control mailbox.

2.5. Multiple Writers to a Common Mailbox

Unpredictable results can occur when more than one process is permitted to write into a single common memory buffer:

- (1) predicting the sequence in which information is written to the common memory buffer may be impossible,
- (2) guaranteeing that all reader clients have seen the contents of the common memory buffer before it is updated may be impossible, and
- (3) identifying the intended reader client audience for any particular memory buffer update may be impossible.

A simple solution is to stipulate that any common memory buffer is permitted to have only a single process writing data into it, although it can have any number of reader clients.

More complex solutions that support the use of a single common memory buffer by more than one writing process are possible. In general, these solutions require the implementation of enhanced flow control and flag field techniques.

3. MAILBOX MANAGEMENT

Mailbox management, as discussed in the following subsections, applies only to implementations with an active common memory manager. There are four fundamental common memory functions used for mailbox management. They are: DECLARE, UNDECLARE, READ and WRITE. Other AMRF extensions exist and are discussed in [7].

Passive common memory management involves the static designation of specific memory areas for each mailbox. Although this might be considered a DECLARE action, there is no equivalent UNDECLARE action. Likewise, READ and WRITE access is uniquely different from passively managed common memories (Section II.1.2.2.).

3.1. DECLARE Access Requirements for Each Mailbox

Before an application can utilize a common memory mailbox for a read or write operation, it must DECLARE the mailbox to the common memory manager. Without this declaration, the common memory manager will not allow access to the mailbox.

A declaration must be issued for each mailbox that is to be accessed and must specify if the mailbox declaration is for a READ function or for a WRITE function. If the mailbox does not already exist in common memory, it is created and space is allocated dynamically. The user application can declare the same mailbox more than once.

There is no logical limit to the number of mailboxes that any application may declare or access. Available memory and/or memory addressing constraints impose the only limitation on the number of common memory participants (applications), number of mailboxes, and mailbox size.

3.2. Perform the READ or WRITE Action

Using the common memory READ and WRITE functions, an application can access the mailbox as often as desired. However, the application must previously have declared that mailbox for the respective access. The common memory manager will return a fatal error status indication if the application is not a client of the mailbox for the requested access.

3.3. UNDECLARE Mailbox Access

After an application has completed all desired accesses to a mailbox and before the application terminates, it should undeclare all previously-declared mailboxes.

When a user application undeclares a mailbox, the common memory manager removes that application from the client list of the mailbox. If the mailbox has other clients, no further action is taken. However, if the mailbox has no other clients, it is removed from common memory and the space it occupied is freed.

4. MAILBOX ACCESS METHODS

There are two possible mailbox access methods: implicit and explicit.

4.1. Explicit Mailbox Access

The word "explicit" is used to imply that there is no common memory manager present. Predesignated, static memory areas are assigned mailbox functions and are accessed directly by more than one application for the purpose of exchanging information.

Since there is no memory manager agent, mailboxes and their starting address and size are static and designated manually by a human agent. The mailbox specifications are loaded (or coded) into each participating application. Coordinating mailbox access (Section II.2.) is not a problem as long as interrupts are disabled when any application accesses a mailbox.

This method of common memory access can result if:

- (1) the participating processes are actually different states within a single program,
- (2) the host operating system does not enforce a memory protection scheme whereby a process is prohibited from accessing memory allocated to a second process,
- (3) the host operating system supports the declaration of "common" memory regions that can be shared by multiple applications, or
- (4) a Multibus implementation with multiple single-board computers (SBC) is used together with a separate memory board that maps into the address space of each SBC.

The mailgram format for explicit common memory would tend to approximate Figure II-2 in order to clearly and easily identify new information and coordinate mailbox access.

4.2. Implicit Mailbox Access

When the implicit method is used, each process associates an internal "logical unit number" (or numeric handle) with a common memory mailbox. This logical unit number is supplied to the process when it creates the mailbox through the respective common memory service. The process then references the logical unit number when it performs READ or WRITE operations in order to exchange mailgrams with the common memory.

The memory location of the mailbox is never accessed directly by any of the participating processes. Instead, the common memory manager has the responsibility of transferring data between the mailbox and the user data area. This activity is performed whenever the application requests a READ or WRITE.

4.3. Conversion of Implicit to Explicit Access and Explicit to Implicit Access

Programs designed for explicit access can be moved to an implicit access environment by inserting the necessary CM_READs, CM_WRITEs, and CM_DECLAREs to create the mailbox. CM_UNDECLAREs are recommended to discontinue mailboxes after their usefulness has expired.

Likewise, programs designed for an implicit mailbox access environment can be moved to an explicit environment by removing the CM_DECLARE, CM_READ, CM_WRITE and CM_UNDECLARE sections and replacing them with the necessary code to identify and access target mailbox memory areas.

5. EXISTING COMMON MEMORY IMPLEMENTATIONS

5.1. Common Memory Maps into the Process's Address Space

In order to reduce the time needed to access areas of common memory, the most desirable implementation is one where the common memory occupies memory in the addressable range of the process. Additional processes within the same "computer" can have access to this same common memory area as long as they have a means of directly accessing that same address space.

The word "computer" is placed in quotes in the preceding paragraph because the reference can be to a single computer such as the Digital Equipment Corporation VAX. It can also refer to a collection of single-board computers resident in a single Multibus chassis.

Within the VAX, the common memory areas can be included in memory space of multiple processes, all active concurrently, by linking to them as a shared READ or WRITE memory area.

For Multibus systems, the local common memory maps into the address space of each of the single-board computers sharing the same bus. Each process within its respective single-board computer sees that memory as its own, and is able to access it directly for READ or WRITE purposes.

5.2. Common Memory as a Separate Process

Some multitasking computer systems used within the AMRF are not immediately amenable to sharing memory space with other active processes. By altering the operating system, it is possible to make them amenable. However, it is actually easier (and safer) to create a separate common memory task.

The common memory information is then transferred between a subset of information maintained by the (user) application process and the actual common memory maintained by this separate task. The interprocess communication is performed using Transmission Control Protocol (TCP) routines [9].

5.3. The Global Common Memory

Information is exchanged with other common memory systems implemented on remote computer hosts via another process resident on the local host, called the network interface process (NIP) [7]. The NIPs have access to all of their local common memory. Connected to each other over a network, NIPs are able to transfer information between common memories. (Figure I-2)

Except for the time delays associated with the transfer of the information across the network, the processes accessing the common memory have no knowledge of what is actually happening to the information that they provide or access.

6. SUMMARY

6.1. Common Memory - An Application Interface

The result of this implemented architecture is that common memory is an application interface to communications with any other processes, both local and remote. It provides a uniform and portable interface for every application. If the application is later moved to a new location, no code changes need to be made for any of its correspondents in order to continue data exchanges. Some changes to the moved application may be necessary if the new location provides a different hardware or operating system architecture. The location changes are reported to the network service and the network service adjusts the mailgram delivery paths [7].

Providing a single application interface allows the application developers to concentrate on the application and frees them from the dependencies of the host-dependent interprocess communication, including network communication.

Further benefits of the common memory interface are listed in the following subsection. This is followed by a short discussion of some perceived drawbacks to common memory as an application interface.

6.1.1. Benefits of Common Memory

The benefits that common memory provides are listed below. Only benefit 3 relies on the availability of network services. However, all benefits are enhanced by the availability and use of network services.

- (1) Asynchronous communications occur between processes. The application process is not interrupted by communications from other processes. It accesses the desired information whenever necessary (i.e., whenever it is ready for it).
- (2) Information can be shared with additional processes with a minimum of effort. Additional processes can read from the same areas of common memory without any action on the part of the initial information provider to deliver it.
- (3) Communication is independent of the location of related processes. The application process does not need to know the location of any other process with which it communicates. If the second process is within the same processor, it is directly connected to the same common

memory. If the second process is located remotely, it has its own common memory with which it communicates. Network services provide the connectivity between common memories.

- (4) It supports coordinated activity between independent processes. Two processes can coordinate their activities by using common memory for command/status information, independent of their respective locations.
- (5) It supports independent evolution of individual processes. With the structured interface between processes that common memory provides, individual processes may evolve in response to changing requirements without mandating equivalent changes in other processes interconnected through the common memory (including the network interface process.)
- (6) It provides a consistent communications methodology for a diverse collection of computers and operating systems. Application processes are freed from machine-dependent communication primitives (e.g., subroutine calls) for both interprocess communications and network communications.

6.1.2. Drawbacks to Common Memory

No paper has yet been published discussing the drawbacks of common memory although one is in preparation [10]. The following are a few thoughts based upon personal experience and do not represent any consensus of opinion.

- (1) The reader of a mailgram does not know the state of the writer of that mailgram:
 - (a) Is the writer on a local or remote host?
 - (b) Is the writer still active or has it terminated or been aborted?
 - (c) Is the logical network connection linking the applications, if applicable, still established?
- (2) The writer of a mailgram does not know the state of the reader of that mailgram:
 - (a) Is the reader on a local or remote host?
 - (b) Is the reader still active or has it terminated or been aborted?
 - (c) Is the logical network connection linking the writer with the reader, if applicable, still established?
 - (d) Has the reader retrieved the mailgram yet?

(3) Why use common memory at all if the correspondents are all known to each other and the number of correspondents and their location will never change? Other communications techniques, such as network message passing, would be more efficient.

With the exception of item (3), these concerns are associated with the delivery and receipt of mailgrams. To some extent, they are answerable with the implementation of a mailbox to acknowledge the receipt of message similar to the flow control mailboxes discussed in Section II.2. However, they actually extend beyond that level of concern.

For example, if a control process is waiting for a data report that originates from a sensory process at irregular time intervals, it is important for the control process to know whether the sensory process is ever going to deliver the next data report. Simply knowing that the sensory process has not halted or been aborted may not be enough: it may be stuck in an unintentional/undesirable endless loop! Perhaps it is necessary for the sensory process to provide a "heartbeat" status in a mailbox?

On the other hand, a sensory process that reports status into a common memory mailbox may not have been programmed to care whether the deposited information is ever read.

This list of drawbacks identifies part of the next logical evolution (or extension) of the common memory architecture. Each drawback is only a minor obstacle that can easily be overcome by having the necessary information provided in common memory, either by the common memory manager or by one or more of the participating processes.

Item (3) raises a good question. The value of common memory as an application interface in a manufacturing or production environment where configuration changes are extremely infrequent has yet to be determined. However, in a research environment where applications evolve and shift from one host system to another and from one architecture to another, the flexibility of common memory has proven invaluable, for all the reasons listed in Section II.6.1.1. Some alternatives to common memory are identified in the following section.

6.2. Alternatives to Common Memory

Simple alternatives that encompass a single computer and networking architecture are easily identified and implemented. Significant difficulty arises when dissimilar architectures comprise the applications environment. In those environments,

the user process is responsible for providing for all communications mechanisms, including gateway routing to devices connected to dissimilar networks.

The variety of hardware and software systems implemented in the AMRF preclude the use of any simple common solution to provide the same service as common memory. Communications among processes within a single computer system would have to use facilities provided by the operating system (if any), or new capabilities similar to those provided by common memory would have to be developed.

Attempts to use services provided by commercially-available network solutions would also be difficult. Commercial network solutions are not available for all computer systems in use in the AMRF. Although more applications (solutions) using the TCP/IP protocol are becoming available all the time, the current migration for networks in manufacturing environments is towards the Manufacturing Automation Protocol (MAP) and the Technical and Office Protocol (TOP). With network companies concentrating on major computer systems at this time, MAP and TOP products are not available for all computer systems.

Any alternative to common memory that must provide service to multiple applications located on several computer systems distributed across multiple network topologies is most likely a connection-oriented message passing solution. For example, although TCP/IP is connectionless, MAP and TOP are connection oriented.

In a connection-oriented network, application process 'A' establishes and maintains a connection to application process 'B' for the bi-directional exchange of messages. Any new application introduced into this environment will have to be accommodated through code changes in those applications with which it must communicate to provide for the additional connection and messages.

III. PC COMMON MEMORY ARCHITECTURE DESCRIPTION

This section describes the common memory architecture as it is implemented for the personal computer using the DOS operating system. Hardware and operating system limitations were the most influential factors affecting the architecture's development. Extensive effort was made to avoid operating system and hardware dependencies in order to maximize portability to other computer architectures and operating systems.

The presentation of specifications in the following sections assumes that the reader has previously become familiar with Section II and the general common memory architecture.

1. MAILBOXES AND MAILGRAMS

1.1 Coordinating Common Memory Access

This topic was originally discussed in Section II.2. The following subsections identify the solutions specifically adopted for the PC common memory without further reference to the problems or alternative solutions.

1.1.1. Read While Write is Active

Coordinating multiple accesses to a common memory mailbox in the PC architecture approximates token passing. DOS is a single-user operating system, so the common memory code is incorporated as "in-line" code to the user application. Likewise, if the user wishes to include other capabilities, such as a network interface program, they too must be included as inline code. (An example of this is shown in Appendix D.) This results in a single program or application when viewed from the DOS perspective.

This single application will not interrupt itself to access common memory. As incorporated in the preceding paragraph, neither will the network interface process interrupt either itself or the user application. Effectively, when any section of the program accesses the common memory, it can be considered to "have the token". The token can only be lost as a consequence of a program crash; appropriate recovery and application debugging steps must be taken following abnormal program terminations.

Other methods (i.e., local software development as well as possible commercial offerings) of implementing multitasking within DOS were considered. They were avoided because they were operating system specific and significantly compromised the portability of the common memory application to other hardware or operating systems.

1.1.2. Update Frequency Exceeds Read Frequency

The PC common memory takes a totally passive role in implementing flow control. The decision to avoid flow control by the common memory manager is intentional: common memory should function just like computer memory. An application can access a mailbox for read or write purposes as often as it desires. The READ operation will retrieve whatever was deposited last, all prior contents having been overwritten.

If flow control is important to processes exchanging information, then it is the responsibility of the respective application processes to provide for it. (Refer to Sections II.2.2 and II.2.5. for further discussion about flow control.)

1.1.3. Read Frequency Exceeds Update Frequency

The PC common memory manager maintains a separate linked list of updated mailboxes for each application process that is a client of common memory. The application process can then be given the update list upon request, thereby providing an indication of mailgram update without the overhead of a mailgram transfer. By reading only the changed mailboxes, the application processes minimize unproductive time spent examining unchanged mailboxes.

1.1.4. Multiple Readers of a Common Mailbox

As stated in Section III.1.1.2, the PC common memory takes a totally passive role regarding flow control. It is the responsibility of the application processes to coordinate the mailgram update procedure if it is desirable that each application have the opportunity to retrieve a mailgram before it is updated.

1.1.5. Multiple Writers to a Common Mailbox

The PC common memory allows more than one application to write to a mailbox. If only one application is to have access to a mailbox for WRITE access, then that application must request exclusive WRITE access when declaring the mailbox.

One potential future extension to PC common memory is to provide for an "access list": the application originally declaring the mailbox is considered to "own" it and may grant other applications READ/WRITE access to its mailbox.

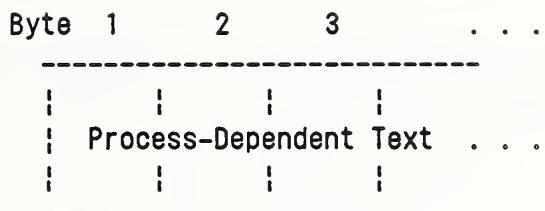


Figure III-1. Generic PC Mailbox Structure

Coordinating writer access in order to avoid a mailgram deposited by one writer from being overwritten by another writer before it has been retrieved by any reader is the responsibility of the user application. The PC common memory takes a totally passive role regarding mailbox flow control.

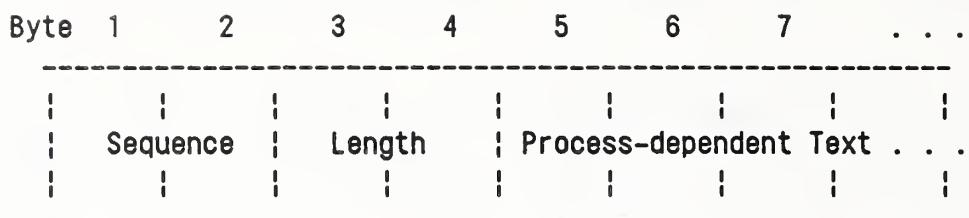
1.2. Mailgram Format

The PC common memory architecture assumes a mailbox structure as shown in Figure III-1.

However, since PC application processes may need to communicate with other applications resident on systems without a common memory manager, the assumed mailbox structure is as shown in Figure III-2. Since the PC common memory manager does not know which mailgram format is in use, it does not manipulate or monitor any of the fields of this alternate mailgram format.

It is the responsibility of the user application reading from the mailbox to be knowledgeable about the mailgram format and the mailgram contents. Furthermore, it is the responsibility of the user application to update and/or provide the information in the SEQUENCE and LENGTH fields in accordance with AMRF conventions. That is, the sequence number must change each time the mailgram is intended to be considered "new". Writing a mailgram with an unchanged sequence number but different process-dependent text may result in the mailgram not being read by an application on a host system that does not have an active common memory manager.

The LENGTH field specifies the number of bytes in the process-dependent section of this mailgram format. A process performing a READ request with PC common memory is returned the total length of the mailgram, including the SEQUENCE and LENGTH field sizes. This number is expected to be different from the value of the LENGTH field in the mailgram.



Sequence is a sequence number attached to the mailgram in the particular mailbox. Every time the text of the mailgram is changed, the sequence number is incremented, so that the change can be detected by examining only the sequence field. The numeric representation is unsigned binary integer, and "wraps" back to zero when the maximum integer representation is incremented.

Length is the length of the process-dependent text in the mailgram in bytes. The numeric representation is unsigned binary integer.

Text is the information portion of the mailgram, defined entirely by the communicating processes.

Figure III-2. Underlying PC Mailgram Format

2. COMMON MEMORY ACCESS METHOD

The PC common memory is implemented using the implicit common memory access method. All mailbox access is performed indirectly through calls to common memory manager routines.

3. MAILBOX MANAGEMENT

The following subsections describe the software interfaces that support explicit common memory, as implemented on the PC. There are four fundamental common memory functions: DECLARE, UNDECLARE, READ and WRITE. Two additional functions, CKMAIL and DISC, have been provided in the PC version of common memory in order to expedite memory access and management, and are also discussed.

3.1. DECLARE Access Requirements for Each Mailbox

Before an application can utilize a common memory mailbox for a READ or WRITE operation, it must DECLARE the mailbox to the common memory manager. Without this declaration, the common memory manager will not allow access to the mailbox.

A declaration must be issued for each mailbox that is to be accessed. However, the application may issue multiple types of access for the declared mailbox within the same mailbox declaration (Appendix B).

If the mailbox does not already exist in common memory, it is created and space is allocated dynamically. The common memory manager maintains two separate lists of clients for each mailbox: a list of reader clients and a list of writer clients. If access is granted, the user application is registered on the appropriate list in accordance with the requested access.

The user application can declare a mailbox more than once. Subsequent declarations for the same mailbox are assumed to be requests to change the application's access rights. That is, these requests can ADD an access that was not previously requested, or CHANGE an existing access. A particular access can only be removed by undeclaring it.

For example, a "XREAD | WRITE" declaration for a mailbox followed by an "READ" declaration results in an access of "READ | WRITE". The character "|" is used to represent a bit-wise OR of the bit representations associated with the specific access function. (The result assumes that a non-fatal error status is returned from function `cm_declare`).

There is no logical limit to the number of mailboxes that any application may declare or access. Available memory and/or memory addressing constraints impose the only limitation on the number of common memory participants (applications) and mailboxes.

3.2. Perform the READ or WRITE Action

Using the common memory READ and WRITE functions, an application can access the mailbox as often as desired. However, the application must previously have declared that mailbox for the respective access. The common memory manager will return a fatal error status if the application is not on the client list of the mailbox or has not been granted the requested (READ or WRITE) access.

3.3. UNDECLARE Mailbox Access

After an application has completed all desired accesses to a mailbox and before the application terminates, it should undeclare all previously-declared mailboxes. This can be done individually for each mailbox, or all mailboxes can be undeclared through a DISConnect request.

When a user application undeclares a mailbox, the common memory manager removes that application from the client list of the mailbox. If the mailbox has other clients, no further action is taken. However, if the mailbox has no other clients, it is removed from common memory and the space it occupied is freed.

3.4. Mailbox Management Extensions

The PC common memory architecture provides two notable extensions to the standard mailbox management services. These are provided in order to meet mailbox access coordination criteria and to simplify the mailbox management process.

3.4.1. DISConnect the Application from Common Memory

If an application intends to terminate its involvement in the common memory, it should undeclare all of its mailboxes. The DISConnect function provided by the common memory manager reduces this to a single request. When an application issues a disconnect request, the common memory manager undeclares all mailboxes for that application.

This function is provided in order to simplify and expedite the departure of an application from the common memory environment. Well-behaved applications can be expected to undeclare their mailboxes either individually or via the DISConnect function. However, in some future multitasking environment, it is conceivable that some task participating in common memory may abruptly terminate. It will be the responsibility of the common memory manager to detect the abnormal termination of the process

("how" will depend on the specific operating system) and undeclare all of that process's mailboxes, thereby purging the common memory of unnecessary mailboxes (i.e., mailboxes to which no other process has declared access).

3.4.2. Check for Mail

Since common memory mailboxes serve as a communications interface, it is important to know when (or whether) new information has been deposited in any mailbox. As mentioned previously (Section III.1.1.3.), a mechanism exists within the PC common memory architecture to identify new information.

Rather than force every application participating in the common memory to examine each of its mailboxes to discover new data, the PC common memory manager maintains a list of mailboxes that have been updated. The application can simply check its mail and is presented with a list of mailboxes that have changed. The common memory manager then starts a new list. The application can now limit its READ requests to those mailboxes whose contents have changed, and thereby minimize the amount of CPU time spent looking at common memory mailboxes.

If the application performs a READ of a mailbox for which there is an entry on the update list (i.e., without checking for mail first), the update list entry for that mailbox is removed by the common memory manager.

IV. PROGRAMMER REFERENCE

The common memory library was developed using the C programming language, as specified by Kernighan and Ritchie [1]. The C programming language was selected in order to maximize portability to other computer architectures.

1. IMPLEMENTATION ISSUES

1.1. The Compiler

The Turbo C (version 1.0) small memory model compiler was used [2,3]. Users who wish to use a different compiler model may need to recompile the common memory library in order to have a consistent data structure definition for variables such as pointers. Defaults are used in the common memory library in order to facilitate such redefinition.

1.2. Memory Allocation and Usage

The common memory manager uses linked lists extensively to maintain information about its clients and their associated mailboxes. The memory space for these data structures is allocated and freed dynamically using functions MALLOC and FREE, respectively. The MALLOC function allocates memory space directly from the user static data area. In the small model, the user application is given 64 Kbytes of static data space.

The 64 Kbyte limit is an artificial one, however, since data space outside of that 64 kbyte can be dynamically allocated and freed using functions FARMALLOC and FARFREE. However, these functions are specific to MS-DOS, and it is desirable to avoid such dependencies.

The amount of data space available to the application process varies with the compiler memory model used. The compact and large models, for example, provide for up to 1 Mbyte of static data space. If the user prefers to use the small memory model (perhaps because compiling is faster), it is preferable that the user application perform the FARMALLOC and FARFREE in order to continue the current level of common memory library independence.

1.3. Interfaces to Languages Other Than C

The common memory library was developed using the C programming language. This does not preclude software developers, who wish to use another programming language compiler for their application, from linking with the library.

For example, the Turbo C user's guide [2] discusses how to link Turbo Pascal and Prolog programs with C object modules. Likewise, it may be feasible to link this C version of the common memory library to other languages, such as Lisp and Ada. This will depend on the compiler implementation. The interested reader must research the respective language reference manual(s).

1.4. Non-portable Common Memory Functions

Only one function referenced in the library is not immediately portable to other C language compilers. This is function "eprintf". It is located in file `cm_utils.c`, and is used by the common memory functions to print debugging statements. Its specificity is to Turbo C and is based on its use of a variable-length argument list. However, other (but not necessarily all) C compilers are known to support variable-length argument lists, so recoding this function to comply with another compiler should not be too difficult.

2. THE COMMON MEMORY INTERFACE LIBRARY

The common memory library provides three service categories: (1) functions that provide common memory access, (2) functions that provide information about common memory usage, and (3) convenience functions. These service categories are described in detail below.

The library distribution kit consists of two INCLUDE (source) files and two object files. The user application must include the source files during the application program compilation period (described in Section IV.5.2). The object files are linked during the application program linking period (described in Section IV.5.3). Section IV.5.1. lists files in the distribution kit.

The source files and the object files have an embedded variable that identifies the common memory library version number. It is critical that both version numbers be identical in order to insure the proper performance of the common memory interface. Section IV.4.3. details how to determine the respective version numbers.

The common memory interface descriptions shown below incorporate the C language convention that a function returns a value through a RETURN statement. These returned values are of type SHORT INT (2 bytes) and report a completion status for each routine. Appendix A lists all possible status codes and describes their significance.

The library routines utilize various data structures to contain and convey specific information to the user application. In the following pages, the data structure for each argument of the functions is identified. This identification is prefaced by the word "TYPE". The data structures are detailed in Appendix C.

2.1. Introduction to Common Memory Manipulations

The actual functions that provide access to common memory are presented in subsequent pages. Before using these functions, it is necessary to have an understanding of the relationship they have to each other and the sequence in which they must be accessed. Section III.3 describes this relationship.

2.2. Functions That Provide Common Memory Access

This section describes the functions that provide common memory access. Only four of them (`cm_declare`, `cm_undeclare`, `cm_write`, and `cm_read`) are actually necessary. The others provide expeditious extensions to these basic functions. (For example, `cm_disc` will undeclare all mailboxes previously declared by the user process, thereby relieving the user from having to submit a series of `cm_undeclare`'s.)

The argument list for each function is described in detail. Each function returns an integer status value that correlates with what the function is to perform (hence the 'int' before each function name). The list of all potential status values that this family of functions can return and their significance is provided in Appendix A.

2.2.1. Function cm_declare

This function creates the necessary application, mailbox, and mailbox client structures within common memory to support future mailbox manipulations by the declaring application.

If an application calls cm_declare for an existing mailbox for READ or XREAD access, the common memory manager will NOT place an entry into its "update list" for that mailbox. The purpose of the update list is to indicate that the mailbox contents have been written SINCE the time of the cm declare or cm_read. It is assumed that the user application will perform an initial cm_read as a matter of course.

int <u>cm_declare</u> (fsm	- INPUT TYPE char *fsm user application name string. String must be null-terminated and must be less than or equal to 32 characters in length (excluding the trailing NULL).
mbxname	- INPUT TYPE char *mbxname mailbox name string. String must be null-terminated and must be less than or equal to 32 characters in length (excluding the trailing NULL).
mbxsize	- INPUT TYPE int mbxsize max size of mailbox to be created.
mbxaccess	- INPUT TYPE int mbxaccess Potential values are READ WRITE XREAD XWRITE but not both of the same kind in the same declaration. The associated constants are listed in <u>cm_const.h</u> .
mbxhandle	- OUTPUT TYPE int *mbxhandle value returned in the integer variable is used as a shorthand reference for mailbox for calls to all other cm routines.
)	

RETURNS: status, as identified in cm_const.h

2.2.2. Function cm_undeclare

Function cm_undeclare is used to remove a user application from a particular mailbox's client list for the specified access. More than one access type may be specified at each call, subject to the access rules identified in the cm_declare section.

If the undeclare action results in a mailbox without any clients, that mailbox is deleted and the space returned to the operating system. Likewise, if the action results in a user application that has no other mailbox's declared, that user application is removed as a common memory client.

If the deleted mailbox is referenced on the undeclaring user application's update list, that specific update entry will also be deleted.

```

int cm_undeclare(fsm
                  - INPUT
                  TYPE char *fsm
                  user application name string.
                  String must be null-terminated
                  and must be less than or equal to
                  32 characters in length
                  (excluding the trailing NULL).

                  mbxaccess - INPUT
                  TYPE int mbxaccess
                  Potential values are
                  READ | WRITE | XREAD | XWRITE,
                  but not both of the same kind in
                  the same declaration. The
                  access constants are
                  listed in cm_const.h.

                  mbxhandle - INPUT
                  TYPE int *mbxhandle
                  Variable value was initially set
                  by cm_declare and is used as a
                  fast way to reference a specific
                  mailbox. Although this does
                  not need to be a pointer, it is
                  specified as such for
                  compatibility with cm_declare
                  (which requires it) and other
                  common memory routines.

)
RETURNS: status, as identified in cm_const.h

```

2.2.3. Function cm_write

Function cm_write is used to transfer a specified number of bytes from the user data area to the common memory mailbox. As a consequence of the write operation, all user application's that have declared READ (or XREAD) access to this mailbox will have an entry made on their update list.

```

int cm_write (fsm
              mbxhandle
              usr_data
              nr_bytes
)
RETURNS: status, as identified in cm_const.h

```

- INPUT
TYPE char *fsm
user application name string.
String must be null-terminated
and must be less than or equal to
32 characters in length
(excluding the trailing NULL).
- INPUT
TYPE int *mbxhandle
Variable value was initially set
cm_declare and is used as a fast
way to reference a specific
mailbox. Although this does not
need to be a pointer, it is
specified as such for
compatibility with cm_declare
(which requires it) and other
common memory routines.
- INPUT
TYPE byte *usr_data
points to user-data area from
which bytes are transferred.
- INPUT
TYPE int *nr_bytes
int variable contains nr of bytes
to be transferred from user data
area to common memory. Although
this does not need to be a
pointer, it is specified as such
for compatibility with cm_read,
which requires it.

2.2.4. Function `cm_read`

Function `cm_read` is used to transfer a specified number of bytes to the user data area from the common memory mailbox. It is recommended that `cm_ckmail` be used together with `cm_read` to minimize unnecessary common memory accesses.

If an entry for this mailbox exists on the update list of this user application, it is removed at completion of the `cm_read` operation.

int <code>cm_read</code> (fsm	- INPUT TYPE <code>char *fsm</code> user application name string. String must be null-terminated, and must be less than or equal to 32 characters in length (excluding the trailing NULL).
mbxhandle	- INPUT TYPE <code>int *mbxhandle</code> Variable value was initially set <code>cm_declare</code> and is used as a fast way to reference a specific mailbox. Although this does not need to be a pointer, it is specified as such for compatibility with <code>cm_declare</code> (which requires it) and other common memory routines.
usr_data	- INPUT TYPE <code>byte *usr_data</code> points to user data area from which bytes are to be transferred.
nr_bytes	- INPUT/OUTPUT TYPE <code>int *nr_bytes</code> When <code>cm read</code> is called, if : (1) the int variable = 0, then all data bytes are transferred from the mailbox to the user's data area. (2) the int variable is not equal to 0, the nr of bytes transferred will be the minimum of (<code>nr_bytes</code> , <code>nr_bytes_in_mbx</code>).

Upon return, the variable pointed to by `nr_bytes` will contain the actual number of bytes transferred. If fewer bytes are transferred to the user data area than are available in the mailbox, an "information-only" status of `I_CM_MOREDATA` is returned to alert the user who may have inadvertently called `cm_read` without clearing the variable pointed to by `nr_bytes`.

It is the user's responsibility to make sure that the data area is large enough to contain the mailgram.

RETURNS: `status`, as identified in `cm_const.h`

2.2.5. Function cm_ckmail

For each user application that is a READER client of common memory, the common memory manager creates and maintains a list of those mailboxes that have changed since the last time the user application read them (i.e., an update list). Whenever a user application writes to a common memory mailbox, the common memory manager posts an entry on this "update" list. If an entry already exists for a changed mailbox, no additional entry is made. The list is maintained in first-in-first-out (FIFO) order.

Entries are removed from this list when a user application calls cm_read for the respective mailbox or when the application calls cm_ckmail. If an update list exists, cm_ckmail returns a pointer to the top of the list and releases the list to the application. If no update list exists, cm_ckmail returns NULL. (If a list is passed to the user application, the common memory manager will start a new list when the next mailbox update arrives.)

Once the list is released to the user application, it is the responsibility of the user application to FREE the memory allocated for the list.

Using the update list, the user application can now perform sequential cm_read operations and only access those mailboxes that have changed since the last read operation.

int cm_ckmail (fsm

- INPUT
 TYPE char *fsm
 user application name string.
 String must be null-terminated
 and must be less than or equal to
 32 characters in length
 (excluding the trailing NULL).

list_ptr

- INPUT
 TYPE is
 struct update_list **list_ptr;

If an update list exists for this user application, cm_ckmail will return a ptr to the top of the update list in this location. If none exists, cm_ckmail will return NULL.

(continued on next page)

nr_entries - INPUT
TYPE int *nr_entries;
If an update list exists, the int
variable will contain the number
of entries in the update list;
else, it will contain ZERO.
)

RETURNS: status, as identified in cm_const.h

2.2.6. Function cm_disc

Function cm_disc provides a shortcut method for a user application to undeclare all of its mailboxes at one time. All data structures within common memory that are associated with that user application are freed. The user application must issue a cm_declare before it can again access common memory variables.

```
int cm_disc (fsm
```

- INPUT

TYPE char *fsm

user application name string.

String must be null-terminated and must be less than or equal to 32 characters in length (excluding the trailing NULL).

```
)
```

RETURNS: status, as identified in cm_const.h

2.3. Functions That Provide Information About Common Memory Usage

The following functions provide a "window" into the common memory environment. They are available to any application. In fact, applications that are not participating in common memory (i.e., applications that do not have any mailboxes declared) can use these calls to determine common memory activity.

These functions were intended for use during the common memory development process. During that time, it was determined that they would be useful for reporting local common memory status to some supervisory and/or monitoring agent (located either across the network or on the local host).

2.3.1. Function `cm_get_fsm_list`

This function allows any user application to determine what user applications are currently active in common memory. Using the `mbxname` and `list_type` function arguments appropriately, the caller can retrieve the list of all user applications known to the common memory manager or only the list of clients (read or write) for a specific mailbox. The information that is returned on the list can be used to solicit other user application (and, indirectly, mailbox) statistics.

`int cm_get_fsm_list(
 mbxname`

- INPUT

TYPE `char *mbxname`
 If `NULL`, this routine will return, through `fsm_list_ptr`, the list of all user application names known to the common memory manager. Argument "list_type" has no effect. If not `NULL`, it must point to a mailbox name. This routine will return a list of all user applications that are a client of that specific mailbox. The argument "list_type" is used to qualify whether the caller wants the list of READER clients or the list of WRITER clients.

`list_type`

- INPUT

TYPE `char`
 May only have the values 'R' (READ) and 'W' (WRITE) when `*mbxname` is non-`NULL`. If `*mbxname` is `NULL`, `list_type` is ignored.

`fsm_list_ptr`

- OUTPUT

TYPE `struct fsm_list_type`
`**fsm_list_ptr`
 This routine will create a linked list of user application names and return a ptr to the top of the list if any user applications exists or `NULL` if none exist. It is the user's responsibility to free this list when it is no longer needed.

int_ptr - OUTPUT
TYPE int *int_ptr
Upon return, the int variable
will equal the number of entries
in the list.
)

RETURNS: status, as identified in cm_const.h

2.3.2. Function `cm_get_mbx_list`

This function allows any user application to determine what mailboxes (mbx) are currently active in common memory. The information that is returned on the list can be used to solicit other common memory user application (and, indirectly, mailbox) statistics. Using the fsm and list_type function arguments appropriately, the caller can retrieve the list of all mailboxes in common memory, or only the list of mailboxes (read or write) for a specific user application.

`int cm_get_mbx_list(`

`fsm`

- INPUT

TYPE `char *fsm`

user application name string. String must be null-terminated and must be less than or equal to 32 characters in length (excluding the trailing NULL). If NULL, this routine will return, through `mbx_list_ptr`, the list of all mailbox names known to the common memory manager. Argument "list_type" has no effect. If not NULL, it must point to a user application name. This routine will return a list of all mailboxes that are a declared by that user application. The argument "list_type" is used to qualify whether the caller wants the list of READER mailboxes or the list of WRITER mailboxes.

`list_type`

- INPUT

TYPE `char`

May only have the values 'R' (READ) and 'W' (WRITE) when `*fsm` is non-NULL. If `*fsm` is NULL, `list_type` is ignored.

`mbx_list_ptr`

- OUTPUT

TYPE `struct mbx_list_type`
`**mbx_list_ptr`

This routine will create a linked list of mailbox names and return

a ptr to the top of the list if any mailboxes exist, or NULL if none exist. It is the user's responsibility to free this list when it is no longer needed.

int_ptr

- INPUT

TYPE int *int_ptr

Upon return, the int variable will equal the number of entries in the list.

)

RETURNS: status, as identified in cm_const.h

2.3.3. Function `cm_get_cm_stats`

This function provides common memory operating statistics. Since the size of the statistics areas is static, the user application must provide a pointer to space in the user data area into which the statistics will be copied. This avoids the overhead associated with dynamic memory allocation in case the user application calls this routine multiple times.

```
int cm_get_cm_stats(
    activity_ptr      - INPUT
                        TYPE cm_activity_stats
                        *activity_ptr
    Points to user-allocated data
    area of appropriate size. This
    routine will copy the activity
    statistics into that data area.
    It has been implemented in this
    fashion to minimize malloc and
    free operations, since it is
    assumed the user will want this
    information more than once.

    client_ptr      - INPUT
                        TYPE cm_client_stats
                        *client_ptr;
    Points to user-allocated data
    area of appropriate size. This
    routine will copy the client
    statistics into that data area.
    It has been implemented in this
    fashion to minimize malloc and
    free operations, since it is
    assumed the user will want this
    information more than once.

)
```

RETURNS: status, as identified in `cm_const.h`

2.3.4. Function `cm_get_fsm_stats`

This function returns the common memory statistics for the specified user application, as identified in structure `cm_fsm_stats_rec` (Appendix C). Since the size of the statistics area is static, the user application must provide a pointer to space in the user data area into which the statistics will be copied. This avoids the overhead associated with dynamic memory allocation in case the user application calls this routine multiple times.

```
int cm_get_fsm_stats(  
    fsm  
    )  
    ptr  
)  
RETURNS: status, as identified in cm_const.h
```

- INPUT
 TYPE `char *fsm`
 user application name string.
 String must be null-terminated
 and must be less than or equal to
 32 characters in length
 (excluding the trailing NULL).
- INPUT
 TYPE `cm_fsm_stats_rec *ptr`

2.3.5. Function `cm_get_mbx_stats`

This function returns the common memory statistics for the specified mailbox, as identified in structure `cm_mbx_stats_rec` (Appendix C). Since the size of the statistics area is static, the user application must provide a pointer to space in the user data area into which the statistics will be copied. This avoids the overhead associated with dynamic memory allocation in case the user application calls this routine multiple times.

```
int cm_get_mbx_stats(  
    mbx                                - INPUT  
                                         TYPE char *mbx  
                                         mailbox name string. String must  
                                         be null-terminated and must be  
                                         less than or equal to 32  
                                         characters in length (excluding  
                                         the trailing NULL).  
    ptr                                - INPUT  
                                         TYPE cm_mbx_stats_rec *ptr  
)  
RETURNS: status, as identified in cm_const.h
```

3. CONVENIENCE FUNCTIONS

During the development of the PC common memory library, several utility functions were developed that might be useful to an application programmer developing a program that uses PC common memory. These functions provide information or functionality directly related to PC common memory usage and are documented below. (Other functions that might be more generally useful are documented in the source listing of file "sfuncs.c". However, they are not specific only to the use of PC common memory. The interested reader is encouraged to examine the "sfuncs.c" file.)

3.1. Function cm_free_update_list

This function frees the memory allocated to an update list. The update list is passed to the user by function cm_ckmail. It is the user's responsibility to free the memory allocated to that linked list. If the user does not free the memory associated with that linked list, it is conceivable that the application may run out of dynamically allocable memory and the common memory manager will be unable to function.

The user may free the blocks of the linked list with calls to the C function "free" or may use this function.

```
void cm_free_update_list(
    list
)
RETURNS: nothing
```

- INPUT
 TYPE struct update_list *list
 This is a pointer to the top of
 the update list. The function
 will free each block of the
 structure, advancing to the next
 block, until NULL is reached.

3.2. Function `cm_get_statusname`

This function converts the status code from its numeric representation into a string containing the corresponding status mnemonic. The conversion is based on the status codes maintained in file `cm_const.h` and listed in Appendix A. All unrecognizable status code is converted into the string "unknown code".

```
char *cm_get_statusname (
    code
    - INPUT
    TYPE int code
    This is presumed to be the status
    code returned by one of the
    common memory interface routines.
)
```

RETURNS: pointer to a string containing the status code associated with the code passed to it. If the code is not recognizable, the function returns the string "unknown code".

3.3. Function `cm_ini`

This function can be used to (1) perform common memory initialization at a known point in the user application program and (2) display the version of the common memory object library.

The user application does not have to call this function. Each common memory interface function, when it is called, checks to see if common memory has been initialized. If common memory has not been initialized, the interface function makes a call to `cm_ini`. Conversely, if common memory has been initialized, no call to `cm_ini` is made.

Function `cm_ini` is responsible for establishing the data structures for the use and management of common memory. During `cm_ini` execution, the value of global variable `CM_DEBUG_LEVEL` is examined. If the value of this variable is greater than zero, `cm_ini` will display the common memory object library version number and will leave the value of `CM_DEBUG_LEVEL` unchanged. If it is less than zero, `cm_ini` sets `CM_DEBUG_LEVEL` to zero to turn off common memory debugging statements.

The interested reader is referred to Section IV.4 for further discussions about global common memory variables `CM_DEBUG_LEVEL`, `CM_GET_STATS`, and `CM_VERSION`.

`void cm_ini()`

RETURNS: nothing

4. GLOBAL COMMON MEMORY VARIABLES

The variables listed in the following subsections are declared by the common memory manager and made available to the user application program. By manipulating these variables, the application program can affect the operation of the local common memory manager. CM_VERSION, as the exception to the previous statement, has no control function. It is only used to provide information about the common memory library.

4.1. CM_DEBUG_LEVEL - Control the Amount of Common Memory Debugging Information Displayed

This variable is declared and initialized in the common memory object library and is located in file cm_globa.h. It is an integer variable whose assigned value determines the amount of diagnostic and/or debugging information the common memory manager will display at the user console. The debugging levels "build" upon each other. That is, selecting a debug level also selects those levels below it (i.e., those with a lower debug level number are also included). The current CM_DEBUG_LEVEL values and their effect are:

- 0 - no debugging data is displayed.
- 1 - display the common memory library version. (This value has been compiled into the object library and cannot be easily accessed or changed by the user application program. Section IV.4.3 provides more information.)
- 2 - 4 <reserved for future use>
- 5 - display identifying messages whenever a mailbox is written or read.
- 6 - display identifying messages whenever an application or mailbox entry is added or deleted to the list maintained by the common memory manager or whenever a cm_declare fails.
- 7 - display identifying messages whenever a mailbox client is added or deleted.
- 8 - display identifying messages whenever a mailbox update notification is posted or removed from the list maintained by the common memory manager.
- 9 - everything (includes 0-8 and more)

The value of CM_DEBUG_LEVEL is initially set to -1 in file cm_globa.h. If the application program changes the value to be greater than or equal to 0, then the common memory manager will not change it. Function cm_ini is responsible for initializing the value of CM_DEBUG_LEVEL to zero if it has a negative value when common memory is initialized.

The only time this variable is examined and (potentially) changed is in function `cm_ini`. This function is only called once by the common memory manager although the user application can call it as frequently as desired. Consequently, the user application can manipulate the value of `CM_DEBUG_VALUE` in order to change or halt the amount of common memory debugging information provided by the common memory manager.

If the user application intends to manipulate or monitor the value of `CM_DEBUG_LEVEL`, it must include the following line in the application program:

```
extern int CM_DEBUG_LEVEL;
```

4.2. CM_GET_STATS - Control the Acquisition of Common Memory Statistics

The PC common memory manager includes the capability of gathering common memory usage statistics. These statistics are available to any application that has access to the common memory manager.

Gathering these statistics takes CPU time away from the application process. If the host processor is slow and cannot provide the level of response necessary for the application, it may be necessary to analyze where the bottleneck is located. If it occurs at the interface between the application and common memory, turning OFF statistics gathering is one way to improve responsiveness. If it occurs within the user application, then changing the value associated with `CM_GET_STATS` will have no affect.

This variable is declared and initialized in the common memory object library, in file `cm_globa.h`, and is of type "boolean". Its initial value is set to TRUE to enable the gathering of statistics.

If the user application intends to manipulate or monitor the value of `CM_GET_STATS`, it must include the following line in the application program:

```
extern boolean CM_DEBUG_LEVEL;
```

In order to assure the proper definition of structure "boolean", this line should appear in the user application source code after the line containing:

```
INCLUDE "cm_types.h"
```

4.3. CM VERSION - Determine Version Numbers of the Common Memory Distribution Components

The PC common memory distribution kit contents are listed in Section IV.5.1. Without some sort of tag or label, it becomes a nearly impossible task to make sure that the components of the distribution kit are all at the same version number.

No matter how much care is exerted, it is always feasible that the common memory version of the two INCLUDE files can get out of synchronization with the version of the common memory object library or even with each other. Therefore, it is the application developer's responsibility to verify that distribution components are at the same version number. If this verification is not done, unintentional, undesirable and unexplainable errors may appear during the execution of the application program. These errors may be a direct result of differences in constant definitions or data structures that may have been introduced in subsequent versions of the common memory distribution.

CM_VERSION is defined and initialized in file cm_const.h via:

```
#define CM_VERSION      "1.0a"
```

However, two logical variables called CM VERSION actually exist: one that is easily accessible to the application program and one that is not so easily accessible. It is neither intended nor desirable for the user application to change the common memory version identifiers.

To determine and verify that all common memory distribution components are at the same version, perform the following steps:

- (1) Manually inspect the source listings of files cm_const.h and cm_types.h. Both files will have an indication of their common memory version number. File cm_const.h will have it as part of a #DEFINE statement and file cm_types.h will have it as part of a comment area near the beginning of the file. Optionally, you can place a "printf" at the beginning of your program to remind you what the INCLUDE file versions are once you have verified that both cm_const.h and cm_types.h are at the same version level. For example:

```
printf("file cm_const.h is at version %s\n",CM_VERSION);
```

- (2) Display the common memory version number in file cm_funcs.obj by compiling and linking a (simple) program that sets the CM_DEBUG_LEVEL equal to an integer value of 1 (or greater) and then calls function cm_ini. Compare this version number with the value determined using step (1).

If the common memory components do not have the same version number, it will be necessary to locate the (most recent) matching set of files.

5. APPLICATION PROGRAM DEVELOPMENT

5.1. The PC Common Memory Distribution Kit

The PC common memory distribution kit consists of this documentation set and the following files:

- `cm_const.h` - contains the definition of constants.
- `cm_types.h` - contains the data structure definitions.
- `cm_funcs.obj` - contains the common memory interface functions. This is distributed in its object file form in order to minimize the potential for user-initiated changes that may later result in inexplicable common memory behavior and to control the possible divergence of common memory interfaces from that identified in this documentation.
- `sfuncs.obj` - contains utilities used by the common memory library, some of which may be useful to the application process, too. It must be included during application program linking.

5.2. Compiling Programs That Use the Common Memory Library

Two of the common memory source files must be included during the user application compilation. They are files `cm_types.h` and `cm_const.h`. File `cm_types.h` references variables defined in `cm_const.h`, so `cm_const.h` must be included before `cm_types.h`. For the C language compiler, the directive is

```
#include "cm_const.h"
#include "cm_types.h"
```

If the user application intends to access or manipulate variables `CM_DEBUG_LEVEL` or `CM_GET_STATS`, the respective "extern" statement should be inserted following the above two lines, as:

```
extern int      CM_DEBUG_LEVEL;
extern boolean CM_GET_STATS;
```

File `cm_types.h` defines all the common memory data structures.

File `cm_const.h` defines all the status return constants. This file is necessary if the user application will be testing the status value returned by each common memory function.

Applications that are written in programming languages other than C and that cannot import these declarations must provide equivalent data structures and constant declarations.

5.3. Linking Programs That Use the Common Memory Library

Two object files must be included in the link process. They are files `cm_funcs.obj` and `sfuncs.obj`.

File `cm_funcs.obj` contains all the common memory interface routines.

File `sfuncs.obj` contains some general support utilities required by the common memory interface routines. They are maintained in a separate file because they are useful to applications that do not need to use the common memory library.

APPENDIX A

 STATUS REPORT CODES
 FOR THE
 COMMON MEMORY INTERFACE FUNCTIONS

The following list identifies the status report codes that the common memory interface routines can return. They are found in file CM_CONST.H.

Although each status is associated with a numeric value, it is strongly recommended that the user application use only the status name (e.g., I_CM_OK) when testing completion status. It is conceivable that the values associated with the status may change in future versions of the common memory interface, whereas status variable names will not change.

The status values are divided into two groups: informational (i.e., non-fatal) and fatal. The boundary is established at variable E_CM_FATALERR. Status values less than E_CM_FATALERR are informational. Status values greater than or equal to E_CM_FATALERR report fatal errors.

Informational status values are used to indicate that the call to the common memory interface function was successfully completed while concurrently providing some additional information affecting that call. Fatal status returns are used to indicate that the call to the common memory interface function was aborted and the cause of the abort.

CONSTANT	HEX VALUE	DESCRIPTION
I_CM_OK	0x0	normal success indicator
I_CM_MBXACTV	0x2	mailbox successfully undeclared, but other applications are still connected. In the case of cm_disc, this refers to the status of one or more mailboxes.
I_CM_DUPMBX	0x4	mailbox was previously declared by this application and for similar access (READ/XREAD or WRITE/XWRITE). If DECLARE was used to change access from exclusive to non-exclusive (or vice-versa), then the change was made. Additional client entries are not made in the client list of the respective mailbox.

CONSTANT	HEX VALUE	DESCRIPTION
I_CM_MOREDATA	0x5	cm_read was successfully performed. However the int variable pointed to by "nr_bytes" was non-zero when the call was made and specified a number of bytes that was less than the number of bytes actually available in the mailbox. Only the number of bytes specified in the int variable pointed to by "nr bytes" was transferred to the user data area. More data is actually available.

Errors greater than 0fx are FATAL errors. This means that the requested action was NOT performed.

E_CM_FATALERR	0x10	defines the start of FATAL ERROR range
E_CM_INSUFFMEM	0x10	insufficient memory. malloc failed
E_CM_MBXERR	0x20	base value for mbx errors
E_CM_MBXNOREAD	0x21	can't have READ - another has XREAD
E_CM_MBXNOXREAD	0x22	can't have XREAD - another fsm has either READ or XREAD
E_CM_MBXNOWRITE	0x24	can't have WRITE - another has XWRITE
E_CM_MBXNOXWRITE	0x28	can't have XWRITE - another fsm has either WRITE or XWRITE
E_CM_MBXSIZEx	0x29	invalid size, returned if: (1) negative size in cm_declare, or (2) byte count < 1 for cm_write, or (3) size doesn't match previously-declared size (for cm_declare), or (4) attempt to write nr_bytes greater than declared size
E_CM_MBXACCESS	0x2a	invalid mbx access, returned if: (1) unrecognizable mbx access code supplied for a cm_declare or cm_undeclare, or (2) invalid list type supplied to cm_get_fsm_list

<u>CONSTANT</u>	<u>HEX</u>	<u>DESCRIPTION</u>
	<u>VALUE</u>	
<u>E_CM_MBXACCBOTH</u>	<u>0x2b</u>	can't declare both READ and XREAD or WRITE and XWRITE in the same mbx declare. However, you can later declare a mbx to be READ after having previously declared it XREAD ... this changes your access and lets other fsm's have access.
<u>E_CM_MBXNAME</u>	<u>0x2c</u>	invalid mbx name - too long or none given
<u>E_CM_MBXNOTDECL</u>	<u>0x2d</u>	returned if : (1) fsm attempts to undeclare a mbx for which it is not a client for the respective access, or (2) fsm attempts to perform read or write actions w/ a mbx for which it is not a client for the respective access, or (3) fsm requests any cm action without being a client of cm, or (4) attempt to get info for a mbx that is not in common memory
<u>E_CM_FSMERR</u>	<u>0x30</u>	base value for fsm errors
<u>E_CM_FSMNAME</u>	<u>0x31</u>	invalid name - too long or none given
<u>E_CM_FSMNOTINCM</u>	<u>0x32</u>	fsm not a common memory client

APPENDIX B

COMMON MEMORY MAILBOX ACCESS CODES

The following list identifies the mailbox access codes that the common memory interface routines will accept. They are found in file CM_CONST.H.

Although each access type is associated with a numeric value, it is strongly recommended that the user application use only the access name (e.g., CM_READ_ACCESS). It is conceivable that the values associated with the access names may change in future versions of the common memory interface, whereas access names will not change.

<u>CONSTANT</u>	<u>HEX</u>	<u>DESCRIPTION</u>
<u>VALUE</u>		
CM_READ_ACCESS	0x1	The declaring application is requesting shared mailbox read access. This will be granted as long as no other application has previously declared exclusive read access. Other applications may declare shared read access for the same mailbox.
CM_XREAD_ACCESS	0x2	The declaring application is requesting exclusive mailbox read access. This will be granted as long as no other application has read or exclusive read access. Use this option with caution, since it also precludes the network from accessing this mailbox. A future extension will allow the mailbox declarer to generate a mailbox access list.
CM_WRITE_ACCESS	0x4	The declaring application is requesting shared mailbox write access. This will be granted as long as no other application has previously declared exclusive write access. Other applications may declare shared write access for the same mailbox.
CM_XWRITE_ACCESS	0x8	The declaring application is requesting exclusive mailbox write access. This will be granted as long as no other application has write or exclusive write access. Use this option with caution, since it also

precludes the network from accessing this mbx. A future extension will allow mbx declarer to generate a mbx access list.

Mailbox access can be specified by using the appropriate access constant individually or by specifying the bit-wise OR of two or more access constants. For example, using the C language syntax, both READ and WRITE access can be specified via

CM_READ_ACCESS | CM_WRITE_ACCESS

The common memory manager checks the validity of the access request. Any access code combination can be submitted except one where the user process is requesting both exclusive and non-exclusive access for the same purpose (e.g., READ or WRITE). The following combinations are illegal:

<u>CM_READ_ACCESS</u>		<u>CM_XREAD_ACCESS</u>		<anything else>
<u>CM_WRITE_ACCESS</u>		<u>CM_XWRITE_ACCESS</u>		<anything else>

APPENDIX C

DATA STRUCTURES
USED IN THE
COMMON MEMORY INTERFACE FUNCTIONS

The data structures specific to the common memory functions available to the user application are detailed below. They are a subset of the complete set of data structures used by the common memory library. They can be found in file `cm_types.h`.

In some instances, constants are referenced in the type definitions. When referenced, they are displayed in uppercase characters. The values for these constants are found in file `cm_const.h`.

C.1. STANDARDIZED DEFINITIONS

The following are some standardized definitions used in subsequent type declarations.

```
typedef unsigned char     byte;
typedef long int          timestamp; /* nr of seconds since
                                         01 Jan 1970 */
```

C.2. UPDATE LIST STRUCTURE

When a mailbox is written to, the common memory manager checks to see if an entry identifying this mailbox already exists on the update list maintained by each reader client of that mailbox. If an entry already exists, no further action is taken. If an entry does not exist, an entry identifying this mailbox is appended to the update list. The list is maintained in FIFO (first-in-first-out) order.

```
typedef struct update_list {
    int mbxhandle;           /* identifies the mbx that has
                                been changed */
    struct update_list *next; /* points to next list entry */
};
```

C.3. COMMON MEMORY STATISTICS

The following structures show how the common memory manager returns common memory utilization statistics.

C.3.1. Function Call Statistics

When a user application requests common memory statistics for this category, the common memory manager places a copy of the statistics into the user-specified data area. The data area must be large enough to contain the data or unpredictable and undesirable side effects may result.

The following structure is used to return usage statistics for each of the user-callable functions of common memory. Its primary use is expected to be as a diagnostic tool.

```

typedef struct {
    char fsm[CM_MAXFSMNAMELENGTH];           /* name of fsm */
    char mbx[CM_MAXMBXNAMELENGTH];           /* name of mbx */
    timestamp when;
    unsigned int nr_times;                   /* nr times this
                                                service called */
} base_stats;

typedef struct {
    base_stats success;                     /* for successes */
    base_stats failure;                   /* for failures */
} group_stats;

typedef struct {
    group_stats cm_declare,                /* id the various */
                cm_undeclare,             /* routines */
                cm_write,
                cm_read,
                cm_ckmail,
                cm_disc,
                cm_get_mbx_list,
                cm_get_fsm_list,
                cm_get_cm_stats,
                cm_get_fsm_stats,
                cm_get_mbx_stats;
} cm_activity_stats;                      /* for cm calls */

typedef struct {
    unsigned int
        mbx_ttl,                         /* total nr of mbx's declared */
        mbx_active,                      /* nr of mbx's currently active */
        fsm_ttl,                         /* total nr fsm's declared */
        fsm_active;                     /* nr fsm's currently active */
} cm_client_stats;

```

C.3.2. Mailbox and Client Lists

Two functions (cm_get_mbx_list and cm_get_fsm_list) return a pointer to a linked list containing the requested information in one of the arguments of the function call. In this case, it is the user's responsibility to FREE the linked list after the list's usefulness has been completed. As mentioned elsewhere (Section IV.1.2), UNIX-portable functions MALLOC and FREE are used within the common memory manager so the user must use function FREE to release the space allocated for this list.

The following structure is used to return a copy of the mailbox list to the common memory client:

```
typedef struct mbx_list_type {
    mbxname[CM_MAXMBXNAMELENGTH];      /* contains name of mbx */
    struct mbx_list_type *next;         /* next block, or NULL */
};
```

The following structure is used to return a copy of the common memory client list to the user application:

```
typedef struct fsm_list_type {
    fsmname[CM_MAXFSMNAMELENGTH];      /* contains name of fsm */
    struct fsm_list_type *next;         /* next block, or NULL */
};
```

C.3.3. Mailbox and Client Statistics

When a user application requests common memory statistics for this category, the common memory manager places a copy of the statistics into the user-specified data area. The data area must be large enough to contain the data or unpredictable and undesirable side effects may result.

The following is returned whenever common memory client information is requested via function cm_get_fsm_stats:

```

typedef struct {
    int nr_read_mbx;          /* nr of read mbx declared */
    unsigned int nr_reads;    /* nr of mbx reads performed */
    timestamp read_time;     /* time of last read */
    char mbx_read[CM_MAXMBXNAMELENGTH]; /* name of last mbx
                                         read, null-terminated */
    int nr_write_mbx;         /* nr of write mbx declared */
    unsigned int nr_writes;   /* nr of writes performed */
    timestamp write_time;    /* time of last write */
    char mbx_write[CM_MAXMBXNAMELENGTH]; /* name of last mbx
                                         written, null-terminated */
    int nr_updates;          /* nr entries on update_list */
}cm_fsm_stats_rec;

```

The following is returned whenever mailbox information is requested via function `cm_get_mbx_stats`:

```

typedef struct {
    int handle;                /* mbxhandle for this mbx */
    int declaredlength;        /* nr bytes declared */
    int msglength;             /* nr bytes currently stored */
    int read_fsms;             /* nr of READ subscribers */
    unsigned int read_accesses; /* nr of READ accesses */
    timestamp read_time;       /* time of last read */
    char reader[CM_MAXFSMNAMELENGTH]; /* name of last reader
                                         fsm, null-terminated */
    int write_fsms;            /* nr of WRITE subscribers */
    unsigned int write_accesses; /* nr of WRITE accesses */
    timestamp write_time;       /* time of last write */
    char writer[CM_MAXFSMNAMELENGTH]; /* name of last writer
                                         fsm, null-terminated */
    char xreader[CM_MAXFSMNAMELENGTH]; /* name of exclusive
                                         reader, null-terminated */
    char xwriter[CM_MAXFSMNAMELENGTH]; /* name of exclusive
                                         writer, null-terminated */
}cm_mbx_stats_rec;

```

APPENDIX D

SAMPLE PROGRAM DEMONSTRATING A COMMON MEMORY MAILBOX INTERACTION BETWEEN TWO LOGICALLY SEPARATE APPLICATIONS

D.1. PURPOSE OF THE PROGRAM

This program serves as a coding example to programmers wishing to develop a PC common memory application. It incorporates all available common memory function calls and is heavily commented in order to document what the program is attempting to do.

The reader is referred to the program listing for further comments about the program.

D.2. PROGRAM LISTING

The program name is "cm_sampl.c" and begins on the following page.

```
/* cm_sample.c - sample program to show usage of the common memory routines.  
   Some procedural and explanatory notes are listed below.
```

- 1) When compiling: Two common memory files must be available during the compile process, cm_types.h and cm_const.h, in order to bring in the common memory data structure definitions. These files do not reference any others, so they may be placed into a library directory, if desired.

File cm_const.h must be included before cm_types.h.

- 2) When linking: Include files cm_funcs.obj and sfuncs.obj in the linking process.

In this sample program, the following is expected to happen:

- 3) Since the common memory version of the two INCLUDE files can get out of sync with the version of the common memory object library, it is a good idea to check the respective versions. To avoid cryptic errors, these version numbers MUST be identical. If they do not match, it will be necessary to locate the (most recent) matching set of files.

Check the respective versions via the following (demonstrated below):

- 1) "printf" the CM_VERSION from the common memory INCLUDE files, or check the include file's source code manually, and
- 2) set CM_DEBUG_LEVEL = 1 to generate a "printf" from the common memory initialization routine that is internal to the (object) library. The initialization routine is called the very first time you attempt to perform a DECLARE, UNDECLARE, READ, WRITE, CKMAIL, or DISCONNECT action. Set CM_DEBUG_LEVEL to zero for those applications where you don't need to monitor the common memory library version level.

- 4) A common memory interaction occurs between two application processes:

- (1) process_1 declares a mailbox for READ and WRITE access and writes something into it.
- (2) process_2 declares the same mailbox for READ access and checks if any of its READ mailboxes have been updated. Since it has only one mailbox, and it was just declared, no updates have yet been posted. process_1 does not perform any read function until

after the next time the mailbox is written to. Thus process_1 misses the first mailgram. The next time process_1 writes into the mailbox, process_2 expects to see an update record (since they are called sequentially from the main program). process_2 will return FALSE until it receives its first mailgram. Thereafter it will always return TRUE.

- (3) process_1 now enters a state where it writes a new mailgram to the mailbox every delta_time seconds. If delta_time seconds have not elapsed since the last time it was called, process_1 will return FALSE. Only after process_1 has written the specified nr_times to the mailbox will it return TRUE.
- (4) Since process_2 is called immediately after each call to process_1, process_2 should display the contents of the newly-written mailbox (almost) immediately after it was written by process_1.
- (5) only when both process_1 and process_2 return TRUE does the main program exit the "when" loop.

5) Get and display the list of all applications active in common memory.

6) Get and display the names of all mailboxes currently in common memory.

7) Get and display the list of READ/XREAD clients for mailbox proc_1_mbx.

8) Get and display the list of mailboxes declared for READ/XREAD access by process_1.

9) Get and display the available common memory statistics for an application that is not in common memory to see the error generated. Use "process_3".

10) Get and display the available common memory statistics for an application that is in common memory. Use "process_1".

11) Get and display the available common memory statistics for a mailbox that is not in common memory to see the error generated. Use "proc_3_mbx".

12) Get and display the available common memory statistics for a mailbox that is in common memory. Use "proc_1_mbx".

13) Undeclare the process_2 mailbox connections by having it call cm_undeclare. Once all of it's mailboxes are undeclared, the common memory manager will remove it from it's client list. Demonstrate this by displaying the list of applications still in common memory. Only 1 should be left.

14) Undeclare the process_1 mailbox connections by having it call cm_disc
 Once all of it's mailboxes are undeclared, the common memory manager will
 remove it from it's client list. Demonstrate this by displaying the list
 of applications still in common memory. The common memory client list
 should be empty.

15) Retrieve and display the common memory statistics.

*/

```
#include <stdio.h>
#include "cm_const.h"                                /* Include before cm_types.h */
#include "cm_types.h"
extern int CM_DEBUG_LEVEL;                          /* declared in the cm library */

enum states {INIT, NORMAL, SHUTDOWN};

/* =====
 |  main  |
 =====
 */
main ()
{
  boolean DONE1 = FALSE;
  boolean DONE2 = FALSE;
  char *mbxname, *fsmname;
  struct mbx_list_type *mbx_list;
  struct fsm_list_type *fsm_list, *tptr;
  cm_fsm_stats_rec fsm_stats;
  cm_mbx_stats_rec mbx_stats;
  cm_activity_stats cm_activity;
  cm_client_stats cm_clients;
  Int status, nr;
  enum states p1_state = INIT,
  p2_state = INIT;

  printf("\n\n\t|-----|\\n"
         "\t| section corresponding to |\\n"
         "\t| |\\n"
         "\t| NOTE # 3 |\\n"
         "\t|-----|\\n"
         "\t| Display the common memory |\\n"
         "\t| library version numbers for both |\\n"
         "\t| the INCLUDE files and the |\\n"
         "\t| linked object file. |\\n"
         "\t|-----|\\n");
  /* the first common memory version number comes from the common memory
   INCLUDE file, cm_const.h */
  printf("\nThe common memory include file is version %s\\n", CM_VERSION);
}
```

```

CM_DEBUG_LEVEL = 1;
/* a common memory debug level of '1' means: display the common memory
object library version nr. The version nr will be displayed by function
cm_ini (located in file cm_utils.c, and thus part of the object file
cm_funcs.obj) the first time a common memory interface function is
called. This can be forced to occur NOW by inserting a call to
function cm_ini here, in order to perform common memory initialization
NOW. The CM_DEBUG_LEVEL must be set to 1 (or greater) before you make
the call to cm_ini in order to display the version number. */
cm_ini();

printf("\n\n\t\t-----|\n"
       "\t\t| section corresponding to |\n"
       "\t\t| |\n"
       "\t\t| NOTE # 4 |\n"
       "\t\t| |\n"
       "\t\t| Two processes access a common |\n"
       "\t\t| mailbox, transferring information |\n"
       "\t\t| from one to the other. |\n"
       "\t\t| |\n"
       "\t\t-----|\n");
while ((! DONE1) || (! DONE2)) { /* continue until both are done */
    DONE1 = process_1(&p1_state);
    DONE2 = process_2(&p2_state);
}

printf("\n\n\t\t-----|\n"
       "\t\t| section corresponding to |\n"
       "\t\t| |\n"
       "\t\t| NOTE # 5 |\n"
       "\t\t| |\n"
       "\t\t| Get and display the list of ALL |\n"
       "\t\t| applications participating in |\n"
       "\t\t| common memory. |\n"
       "\t\t| |\n"
       "\t\t-----|\n");
printf("\n\nmain: Get and display the list of ALL applications currently"
      "\n      participating in common memory\n");
status = cm_get_fsm_list(NULL, '', &fsm_list, &nr);
if (status >= E_CM_FATALERR) /* was there a fatal error */
    printf("main: FATAL ERROR %s returned from cm_get_fsm_list\n",
          cm_get_statusname(status));

```

```

else
  if (fsm_list) {
    printf("main: following %d applications are active in common memory...\n", nr);
    while (fsm_list) { /* display the application names */
      printf("\t%s\n", fsm_list->fsmname);
      tptr = fsm_list->next;
      free(fsm_list); /* free the blocks when no longer needed */
      fsm_list = tptr;
    }
  }else printf("main: no applications active in common memory\n");

printf("\n\n\t\t-----\n\t\t| section corresponding to |\n\t\t| NOTE # 6 |\n\t\t| Get and display the list of ALL |\n\t\t| mailboxes in common memory. |\n\t\t-----\n");
printf("\n\nmain: Get and display the list of ALL mailboxes in common memory\n");
status = cm_get_mbx_list(NULL, '', &mbx_list, &nr);
if (status >= E_CM_FATALERR) /* was there a fatal error */
  printf("main: FATAL ERROR %s returned from cm_get_mbx_list\n",
         cm_get_statusname(status));
else
  if (mbx_list) {
    printf("main: following %d mailboxes are declared in common memory...\n", nr);
    while (mbx_list) { /* display the application names */
      printf("\t%s\n", mbx_list->mbxname);
      tptr = mbx_list->next;
      free(mbx_list); /* free the blocks when no longer needed */
      mbx_list = tptr;
    }
  }else printf("main: no mailboxes declared in common memory\n");

printf("\n\n\t\t-----\n\t\t| section corresponding to |\n\t\t| NOTE # 7 |\n\t\t| Get and display the list of |\n\t\t| READ/XREAD clients for mailbox |\n\t\t| 'proc_1_mbx'. |\n\t\t-----\n");
printf("\n\nmain: Get and display the list of READ/XREAD clients for mailbox"
      "\n      proc_1_mbx.\n");

```

```

mbxname = "proc_1_mbx";
status = cm_get_fsm_llist(mbxname,'r',&fsm_llist,&nr);
if (status >= E_CM_FATALERR)           /* was there a fatal error */
    printf("main: FATAL ERROR %s returned from cm_get_fsm_llist\n",
           cm_get_statusname(status));
else
    if (fsm_llist) {
        printf("main: %s has following %d READ/XREAD clients ...\n", mbxname,nr);
        while (fsm_llist) {           /* display the application names */
            printf("\t%s\n",fsm_llist->fsmname);
            tptr = fsm_llist->next;
            free(fsm_llist);          /* free the blocks when no longer needed */
            fsm_llist = tptr;
        }
    } else printf("main: no READ/XREAD applications declared mailbox %s\n",mbxname);

printf("\n\n\t\t-----|\n"
      "\t\t | section corresponding to |\n"
      "\t\t | |\n"
      "\t\t | NOTE # 8 |\n"
      "\t\t | |\n"
      "\t\t | Get and display the list of |\n"
      "\t\t | mailboxes declared for READ/ |\n"
      "\t\t | XREAD access by 'process_1'. |\n"
      "\t\t | |\n"
      "\t\t-----|\n");
printf("\n\nmain: Get and display the list of mailboxes declared for READ/XREAD"
      "\n      by process_1.\n");
fsmname = "process_1";
status = cm_get_mbx_llist(fsmname,'r',&mbx_llist,&nr);
if (status >= E_CM_FATALERR)           /* was there a fatal error */
    printf("main: FATAL ERROR %s returned from cm_get_mbx_llist\n",
           cm_get_statusname(status));
else
    if (mbx_llist) {
        printf("main: %s has following %d READ mailboxes ...\n", fsmname,nr);
        while (mbx_llist) {           /* display the application names */
            printf("\t%s\n",mbx_llist->mbxname);
            tptr = mbx_llist->next;
            free(mbx_llist);          /* free the blocks when no longer needed */
            mbx_llist = tptr;
        }
    } else printf("main: no READ/XREAD mailboxes declared by %s\n",fsmname);

```

```

printf("\n\n\t\t-----\n\t\t| section corresponding to\n\t\t| NOTE # 9\n\t\t| Get and display the available\n\t\t| common memory statistics for an\n\t\t| application that is not in\n\t\t| common memory. Error will be\n\t\t| generated. Use 'process_3'.\n\t\t|-----\n\t\t-----\n\nmain: Get and display the available common memory statistics for\n      an application that doesn't exist. Expect an error report.\n");
fsmname = "process_3";
status = cm_get_fsm_stats(fsmname,&fsm_stats);
if (status >= E_CM_FATALERR)           /* was there a fatal error */
    printf("main: FATAL ERROR %s returned from cm_get_fsm_stats\n",
           cm_get_statusname(status));
else {
    printf("\t%d read mbx's\n\t%u reads performed\n\tlast one at ",
           fsm_stats.nr_read_mbx,
           fsm_stats.nr_reads);
    printf("%s\tmbx: %s\n\t%d write mbx's\n\t%u writes performed\n\t"
           "last one at ",
           (fsm_stats.read_time) ? ctime(&fsm_stats.read_time) : "<none>\n",
           fsm_stats.mbx_read,
           fsm_stats.nr_write_mbx,
           fsm_stats.nr_writes);
    printf("%s\tmbx: %s\n\t%d updates on the update list\n",
           (fsm_stats.write_time) ? ctime(&fsm_stats.write_time) : "<none>\n",
           fsm_stats.mbx_write,
           fsm_stats.nr_updates);
}

```

```

printf("\n\n\t\t-----\n\t\t| section corresponding to\n\t\t| NOTE # 10\n\t\t| Get and display the available\n\t\t| common memory statistics for\n\t\t| process_1.\n\t\t|-----\n\t\t-----\n\nmain: Get and display the available common memory statistics for\n      an application that does exist.\n");
fsmname = "process_1";
status = cm_get_fsm_stats(fsmname,&fsm_stats);

```

```

if (status >= E_CM_FATALERR) /* was there a fatal error */
    printf("main: FATAL ERROR %s returned from cm_get_fsm_stats\n",
           cm_get_statusname(status));
else {
    printf("\t%d read mbx's\n\t%u reads performed\n\tlast one at ",
           fsm_stats.nr_read_mbx,
           fsm_stats.nr_reads);
    printf("%s\tmbx: %s\n\t%d write mbx's\n\t%u writes performed\n\t"
           "last one at ",
           (fsm_stats.read_time) ? ctime(&fsm_stats.read_time) : "<none>\n",
           fsm_stats.mbx_read,
           fsm_stats.nr_write_mbx,
           fsm_stats.nr_writes);
    printf("%s\tmbx: %s\n\t%d updates on the update list\n",
           (fsm_stats.write_time) ? ctime(&fsm_stats.write_time) : "<none>\n",
           fsm_stats.mbx_write,
           fsm_stats.nr_updates);
}

printf("\n\n\t-----\n\t\t| section corresponding to\n\t\t| NOTE # 11\n\t\t| Get and display the available\n\t\t| common memory statistics for a\n\t\t| nonexistent mailbox.\n\t\t|-----\n\t\t|\n");
printf("\n\nmain: Get and display the available common memory statistics for"
      "\n      a mailbox that doesn't exist.\n");
mbxname = "proc_3_mbx";
status = cm_get_mbx_stats(mbxname, &mbx_stats);
if (status >= E_CM_FATALERR) /* was there a fatal error */
    printf("main: FATAL ERROR %s returned from cm_get_mbx_stats\n",
           cm_get_statusname(status));
else {
    printf("\tmbxhandle %d, declared length %d, msg length %d\n\t"
           "%d readers w/ %u read actions\n\t  last by %s @",
           mbx_stats.handle, mbx_stats.declaredlength, mbx_stats.msglength,
           mbx_stats.read_fsms, mbx_stats.read_accesses,
           mbx_stats.reader);
    printf("\t%s\t%d writers w/ %u write actions\n\t  last by %s @",
           (mbx_stats.read_time) ? ctime(&mbx_stats.read_time) : "<none>\n",
           mbx_stats.write_fsms, mbx_stats.write_accesses,
           mbx_stats.writer);
}

```

```

printf(" %s\txreader: %s\n\txwriter: %s\n",
      (mbx_stats.write_time) ? ctime(&mbx_stats.write_time) : "<none>\n",
      mbx_stats.xreader,
      mbx_stats.xwriter);
}

printf("\n\n\t\t-----|\n\t\t| section corresponding to |\n\t\t| |\n\t\t| NOTE # 12 |\n\t\t| |\n\t\t| Get and display the available |\n\t\t| common memory statistics for |\n\t\t| 'proc_1_mbx'. |\n\t\t| |\n\t\t-----|\n");
printf("\n\nmain: Get and display the available common memory statistics for"
      "\n      a mailbox that does exist.\n");
mbxname = "proc_1_mbx";
status = cm_get_mbx_stats(mbxname, &mbx_stats);
if (status >= E_CM_FATALERR)           /* was there a fatal error */
  printf("main: FATAL ERROR %s returned from cm_get_mbx_stats\n",
         cm_get_statusname(status));
else {
  printf("\tmbxhandle %d, declared length %d, msg length %d\n\t"
         "%d readers w/ %u read actions\n\t last by %s @",
         mbx_stats.handle, mbx_stats.declaredlength, mbx_stats.msglength,
         mbx_stats.read_fsms, mbx_stats.read_accesses,
         mbx_stats.reader);
  printf(" %s\t%d writers w/ %u write actions\n\t last by %s @",
         (mbx_stats.read_time) ? ctime(&mbx_stats.read_time) : "<none>\n",
         mbx_stats.write_fsms, mbx_stats.write_accesses,
         mbx_stats.writer);
  printf(" %s\txreader: %s\n\txwriter: %s\n",
         (mbx_stats.write_time) ? ctime(&mbx_stats.write_time) : "<none>\n",
         mbx_stats.xreader,
         mbx_stats.xwriter);
}

```

```

printf("\n\n\t\t-----|\n"
      "\t\t| section corresponding to |\n"
      "\t\t| |\n"
      "\t\t| NOTE # 13 |\n"
      "\t\t| |\n"
      "\t\t| Cause process_2 to undeclare all |\n"
      "\t\t| Its common memory mailbox |\n"
      "\t\t| connections. Display the list |\n"
      "\t\t| or remaining common memory |\n"
      "\t\t| clients. Only one client should |\n"
      "\t\t| be left. |\n"
      "\t\t| |\n"
      "\t\t| Process_2 will use cm_undeclare |\n"
      "\t\t| to disconnect from the mailboxes |\n"
      "\t\t| |\n"
      "\t\t-----|\n");
printf("\n\nmain: Cause process_2 to undeclare its common memory mailbox connections."
      "\n      Then display the common memory client list. Expect only one client.\n");
p2_state = SHUTDOWN;
DONE2 = process_2(&p2_state);
status = cm_get_fsm_list(NULL, '', &fsm_list, &nr);
if (status >= E_CM_FATALERR)           /* was there a fatal error */
  printf("main: FATAL ERROR %s returned from cm_get_fsm_list\n",
         cm_get_statusname(status));
else
  if (fsm_list) {
    printf("main: following %d applications are active in common memory...\n", nr);
    while (fsm_list) {           /* display the application names */
      printf("\t%s\n", fsm_list->fsmname);
      tptr = fsm_list->next;
      free(fsm_list);           /* free the blocks when no longer needed */
      fsm_list = tptr;
    }
  } else printf("main: no applications active in common memory\n");

```

```

printf("\n\n\t\t-----\n\t\t| section corresponding to |\n\t\t| NOTE # 14 |\n\t\t| Cause process_1 to undeclare all |\n\t\t| its common memory mailbox |\n\t\t| connections. Display the list |\n\t\t| or remaining common memory |\n\t\t| clients. Only one client should |\n\t\t| be left. |\n\t\t|\n\t\t| Process_1 will use cm_disc to |\n\t\t| disconnect from the mailboxes. |\n\t\t|\n\t\t-----\n\t\t-----\n\nmain: Cause process_1 to undeclare its common memory mailbox connections.\n\t\tThen display the common memory client list. Expect zero clients.\n\np1_state = SHUTDOWN;\nDONE1 = process_1(&p1_state);\nstatus = cm_get_fsm_list(NULL, '', &fsm_list, &nr);\nif (status >= E_CM_FATALERR) /* was there a fatal error */\nprintf("main: FATAL ERROR %s returned from cm_get_fsm_list\n",
cm_get_statusname(status));\nelse\nif (fsm_list) {\nprintf("main: following %d applications are active in common memory...\n", nr);\nwhile (fsm_list) /* display the application names */\nprintf("\t%s\n", fsm_list->fsmname);\nptr = fsm_list->next;\nfree(fsm_list); /* free the blocks when no longer needed */\nfsm_list = ptr;\n}\n} else printf("main: no applications active in common memory\n");

printf("\n\n\t\t-----\n\t\t| section corresponding to |\n\t\t| NOTE # 15 |\n\t\t| Retrieve and display the common |\n\t\t| memory usage statistics. |\n\t\t|\n\t\t-----\n\t\t-----\n\nstatus = cm_get_cm_stats(&cm_activity, &cm_clients);\nif (status >= E_CM_FATALERR) /* was there a fatal error */\nprintf("main: FATAL ERROR %s returned from cm_get_cm_stats\n",
cm_get_statusname(status));\n

```

```

else {
    printf("A total of %u mbx's were created. %u still active.\n"
          "A total of %u fsm's were clients of cm. %u still active.\n",
          cm_clients.mbx_ttl, cm_clients.mbx_active,
          cm_clients.fsm_ttl, cm_clients.fsm_active);
    dmp_base_stats(cm_activity.cm_declare.failure, "cm_declare failures");
    dmp_base_stats(cm_activity.cm_declare.success, "cm_declare successes");
    dmp_base_stats(cm_activity.cm_undeclare.failure, "cm_undeclare failures");
    dmp_base_stats(cm_activity.cm_undeclare.success, "cm_undeclare successes");
    dmp_base_stats(cm_activity.cm_write.failure, "cm_write failures");
    dmp_base_stats(cm_activity.cm_write.success, "cm_write successes");
    dmp_base_stats(cm_activity.cm_read.failure, "cm_read failures");
    dmp_base_stats(cm_activity.cm_read.success, "cm_read successes");
    dmp_base_stats(cm_activity.cm_ckmail.failure, "cm_ckmail failures");
    dmp_base_stats(cm_activity.cm_ckmail.success, "cm_ckmail successes");
    dmp_base_stats(cm_activity.cm_disc.failure, "cm_disc failures");
    dmp_base_stats(cm_activity.cm_disc.success, "cm_disc successes");
    dmp_base_stats(cm_activity.cm_get_mbx_list.failure, "cm_get_mbx_list failures");
    dmp_base_stats(cm_activity.cm_get_mbx_list.success, "cm_get_mbx_list successes");
    dmp_base_stats(cm_activity.cm_get_fsm_list.failure, "cm_get_fsm_list failures");
    dmp_base_stats(cm_activity.cm_get_fsm_list.success, "cm_get_fsm_list successes");
    dmp_base_stats(cm_activity.cm_get_cm_stats.failure, "cm_get_cm_stats failures");
    dmp_base_stats(cm_activity.cm_get_cm_stats.success, "cm_get_cm_stats successes");
    dmp_base_stats(cm_activity.cm_get_fsm_stats.failure, "cm_get_fsm_stats failures");
    dmp_base_stats(cm_activity.cm_get_fsm_stats.success, "cm_get_fsm_stats successes");
    dmp_base_stats(cm_activity.cm_get_mbx_stats.failure, "cm_get_mbx_stats failures");
    dmp_base_stats(cm_activity.cm_get_mbx_stats.success, "cm_get_mbx_stats successes");
}
}

/*
 * | dmp_base_stats |
 */
int dmp_base_stats(func,msg)
base_stats func;
char *msg;
{
    printf("%s: %u\n\tlast by fsm: %s\n\tfor mbx: %s\n\tat: %s\n",
           msg,func.nr_times,
           (func.fsm) ? func.fsm : " ",
           (func.mbx) ? func.mbx : " ",
           (func.when) ? ctime(&func.when) : "\n");
}

```

```

/* ||||||||||||||| This process initializes itself the first time it is called.
| process_1 | Next, it writes into a mailbox, records the time of write,
||||||||||||||| and decrements the number of times that it is supposed to
| write into the mailbox. If the result is zero, it
| returns TRUE, else it returns FALSE.

The next time it is called, it checks if delta_time has
elapsed. If not, it returns FALSE. If so, it once again
writes into the mailbox, records the time of the write,
decrements the counter, and returns either TRUE or FALSE
according to the algorithm described before.

At shutdown, the mailboxes are disconnected via cm_disc.

*/
boolean process_1(state)
enum states *state;
{ char data[200];
  int mbxsize, mbxaccess;
  timestamp time_now;
  struct update_llist *update_ptr, *tptr;
  int nrbytes, status;

  static timestamp last_write,           /* timestamp for last cm_write */
            delta_time = 10;          /* nr seconds between cm_write */
  static int nr_writes = 5;             /* total nr of writes to perform */
  static int mbxhandle;
  static char *fsm = "process_1";
  static char *mbxname = "proc_1_mbx";

  switch (*state) {
    case INIT : {
      *state = NORMAL;           /* for next call */
      time(&last_write);       /* fake a last_write time */
      last_write -= 10;          /* take off 10 seconds */
      mbxsize = 150;
      mbxaccess = CM_READ_ACCESS | CM_WRITE_ACCESS;
      status = cm_declare(fsm, mbxname, mbxsize, mbxaccess, &mbxhandle);
      if (status >= E_CM_FATALERR)
        printf("process_1: FATAL ERROR %s returned from cm_declare\n",
               cm_get_statusname(status));
      break;
    }
  }
}

```

```

case NORMAL :
    time(&time_now);
    if ((time_now - last_write) >= delta_time) {
        last_write = time_now;
        sprintf(data,"process_1 wrote mbx (%d writes remaining) on %s",
                nr_writes-1, ctime(&time_now));
        nrbytes = strlen(data);
        status = cm_write(fsm, &mbxhandle, data, &nrbytes);
        if (status >= E_CM_FATALERR)
            printf("process_1: FATAL ERROR %s returned from cm_write\n",
                   cm_get_statusname(status));
        return ((--nr_writes) == 0);
    }else return(FALSE);           /* insufficient time elapsed */
    break;
case SHUTDOWN :
    status = cm_disc(fsm);
    if (status >= E_CM_FATALERR) {
        printf("process_1: FATAL ERROR %s returned from cm_disc\n",
               cm_get_statusname(status));
        return (TRUE);           /* give MAIN the option for clean getaway */
    }
    break;
}
}

/*
| process_2 | This process initializes itself the first time it is called.
|           | Next, and every time it is called thereafter, it checks its
|           | mail to see if any of the mailboxes that it has declared as
|           | READ or XREAD access have been updated. If none, it returns
|           | a boolean value indicating whether it has read any of its
|           | declared mailboxes since they were declared. If updated
|           | mailboxes exist, they are sequentially displayed, the
|           | update list freed, and a boolean value is (set and) returned
|           | to indicate that a mailbox has been read.

At shutdown, the mailboxes are disconnected via cm_undeclare.
*/
boolean process_2(state)
enum states *state;
{ char data[200];
  int mbxsize, mbxaccess;
  timestamp time_now;
  struct update_list *update_ptr, *tptr;
  int nrbytes, nr_updates, status;

```

```

static boolean read_one = FALSE;
static int mbxhandle;
static char *fsm = "process_2";
static char *mbxname = "proc_1_mbx";

switch (*state) {
    case INIT :
        mbxsize = 150;
        mbxaccess = CM_READ_ACCESS;
        status = cm_declare(fsm, mbxname, mbxsize, mbxaccess, &mbxhandle);
        if (status >= E_CM_FATALERR) {
            printf("process_2: FATAL ERROR %s returned from cm_declare\n",
                   cm_get_statusname(status));
            return (TRUE);
        }
        *state = NORMAL;
        break;
    case NORMAL :
        status = cm_ckmail(fsm, &update_ptr, &nr_updates);
        if (status >= E_CM_FATALERR) {
            printf("process_2: FATAL ERROR %s returned from cm_ckmail\n",
                   cm_get_statusname(status));
            return (TRUE);           /* try to give MAIN a graceful exit */
        }
        if (update_ptr) {           /* read updated mailboxes */
            read_one = TRUE;       /* remember that you read at least one */
            tptr = update_ptr;     /* remember start of list for later FREE */
            time(&time_now);
            printf("\nprocess_2: It is now %s", ctime(&time_now));
            printf("process_2: %d mailbox update notifications\n", nr_updates);
            while (update_ptr) {
                nrbytes = 0;          /* to insure we get all of the mailgram */
                status = cm_read(fsm, &update_ptr->mbxhandle, data, &nrbytes);
                if (status >= E_CM_FATALERR) {
                    printf("process_2: FATAL ERROR %s returned from cm_read\n",
                           cm_get_statusname(status));
                    exit();
                }
                data[nrbytes] = 0;      /* trailing NULL so we can treat as string w/ printf */
                printf("\tmbxhandle: %d, msg is %d bytes long\n"
                       "\tstatus returned is %s, or %x (hex)\n\tCONTENTS: %s",
                       update_ptr->mbxhandle, nrbytes, cm_get_statusname(status), status, data);
                update_ptr = update_ptr->next;
            }
            cm_free_update_list(tptr);      /* FREE the update list */
            return (read_one);
        }
        break;
}

```

```
case SHUTDOWN :
    mbxaccess = CM_READ_ACCESS;
    status = cm_undeclare(fsm, mbxaccess, &mbxhandle);
    If (status >= E_CM_FATALERR) {
        printf("process_1: FATAL ERROR %s returned from cm_undeclare\n",
               cm_get_statusname(status));
        return (TRUE);           /* give MAIN the option for clean getaway */
    }
    break;
}
return (read_one);
}
```

D.3. SAMPLE PROGRAM OUTPUT

The output generated by the sample program is shown below. The correlation between sections of the output and notations made in the program listing are clearly identified by in-line comments. In some cases, the sample program generates a line that is wider than can be represented on these typewritten pages. In these cases, the line is broken into two parts and the information is displayed on two sequential lines.

The output begins on the following page.

| section corresponding to |
|

NOTE # 3

| Display the common memory
| library version numbers for both
| the INCLUDE files and the
| linked object file.
|

The common memory include file is version 1.0a
cm_inl: using common memory library version 1.0a

| section corresponding to |
|

NOTE # 4

| Two processes access a common
| mailbox, transferring information
| from one to the other.
|

process_2: It is now Sat Feb 27 19:50:55 1988
process_2: 1 mailbox update notifications
mbxhandle: 1, msg is 69 bytes long
status returned is I_CM_OK, or 0 (hex)
CONTENTS: process_1 wrote mbx (4 writes remaining)
on Sat Feb 27 19:50:55 1988

process_2: It is now Sat Feb 27 19:51:05 1988
process_2: 1 mailbox update notifications
mbxhandle: 1, msg is 69 bytes long
status returned is I_CM_OK, or 0 (hex)
CONTENTS: process_1 wrote mbx (3 writes remaining)
on Sat Feb 27 19:51:05 1988

```
process_2: It is now Sat Feb 27 19:51:15 1988
process_2: 1 mailbox update notifications
  mbxhandle: 1, msg is 69 bytes long
  status returned is I_CM_OK, or 0 (hex)
  CONTENTS: process_1 wrote mbx (2 writes remaining)
            on Sat Feb 27 19:51:15 1988
```

```
process_2: It is now Sat Feb 27 19:51:25 1988
process_2: 1 mailbox update notifications
  mbxhandle: 1, msg is 69 bytes long
  status returned is I_CM_OK, or 0 (hex)
  CONTENTS: process_1 wrote mbx (1 writes remaining)
            on Sat Feb 27 19:51:25 1988
```

```
process_2: It is now Sat Feb 27 19:51:35 1988
process_2: 1 mailbox update notifications
  mbxhandle: 1, msg is 69 bytes long
  status returned is I_CM_OK, or 0 (hex)
  CONTENTS: process_1 wrote mbx (0 writes remaining)
            on Sat Feb 27 19:51:35 1988
```

```
=====
| section corresponding to
|                               |
|       NOTE # 5
|                               |
|       Get and display the list of ALL
|       applications participating in
|       common memory.
|                               |
=====
```

```
main: Get and display the list of ALL applications currently
      participating in common memory
main: following 2 applications are active in common memory...
  process_1
  process_2
```

```
=====
| section corresponding to
|                               |
|       NOTE # 6
|                               |
|       Get and display the list of ALL
|       mailboxes in common memory.
|                               |
=====
```

```
main: Get and display the list of ALL mailboxes in common memory
main: following 1 mailboxes are declared in common memory...
proc_1_mbx
```

```
-----+
| section corresponding to |
| |
| NOTE * 7 |
| |
| Get and display the list of |
| READ/XREAD clients for mailbox |
| 'proc_1_mbx'. |
| |
-----+
```

```
main: Get and display the list of READ/XREAD clients for mailbox
proc_1_mbx.
main: proc_1_mbx has following 2 READ/XREAD clients ...
process_1
process_2
```

```
-----+
| section corresponding to |
| |
| NOTE * 8 |
| |
| Get and display the list of |
| mailboxes declared for READ/ |
| XREAD access by 'process_1'. |
| |
-----+
```

```
main: Get and display the list of mailboxes declared for
      READ/XREAD by process_1.
main: process_1 has following 1 READ mailboxes ...
proc_1_mbx
```

```
-----  
| section corresponding to  
|  
| NOTE # 9  
|  
| Get and display the available  
| common memory statistics for an  
| application that is not in  
| common memory. Error will be  
| generated. Use 'process_3'.  
|  
-----
```

```
main: Get and display the available common memory statistics for  
an application that doesn't exist. Expect an error report.  
main: FATAL ERROR E_CM_FSMNOTINCM returned from cm_get_fsm_stats
```

```
-----  
| section corresponding to  
|  
| NOTE # 10  
|  
| Get and display the available  
| common memory statistics for  
| process_1.  
|  
-----
```

```
main: Get and display the available common memory statistics for  
an application that does exist.  
1 read mbx's  
0 reads performed  
last one at <none>  
mbx: <none>  
1 write mbx's  
5 writes performed  
last one at Sat Feb 27 19:51:35 1988  
mbx: proc_1_mbxx  
1 updates on the update list
```

```
-----  
| section corresponding to  
|  
| NOTE # 11  
|  
| Get and display the available  
| common memory statistics for a  
| nonexistent mailbox.  
-----
```

```
main: Get and display the available common memory statistics for  
a mailbox that doesn't exist.  
main: FATAL ERROR E_CM_MBXNOTDECL returned from cm_get_mbx_stats
```

```
-----  
| section corresponding to  
|  
| NOTE # 12  
|  
| Get and display the available  
| common memory statistics for  
| 'proc_1_mbx'.  
-----
```

```
main: Get and display the available common memory statistics for  
a mailbox that does exist.  
mbxhandle 1, declared length 150, msg length 69  
2 readers w/ 5 read actions  
    last by process_2 @ Sat Feb 27 19:51:35 1988  
1 writers w/ 5 write actions  
    last by process_1 @ Sat Feb 27 19:51:35 1988  
xreader: <none>  
xwriter: <none>
```

```
-----  
| section corresponding to  
|  
| NOTE # 13  
|  
| Cause process_2 to undeclare all  
| its common memory mailbox  
| connections. Display the list  
| or remaining common memory  
| clients. Only one client should  
| be left.  
|  
| Process_2 will use cm_undeclare  
| to disconnect from the mailboxes  
|  
-----
```

main: Cause process_2 to undeclare its common memory mailbox
connections. Then display the common memory client list.
Expect only one client.
main: following 1 applications are active in common memory...
process_1

```
-----  
| section corresponding to  
|  
| NOTE # 14  
|  
| Cause process_1 to undeclare all  
| its common memory mailbox  
| connections. Display the list  
| or remaining common memory  
| clients. Only one client should  
| be left.  
|  
| Process_1 will use cm_disc to  
| disconnect from the mailboxes.  
|  
-----
```

main: Cause process_1 to undeclare its common memory mailbox
connections. Then display the common memory client list.
Expect zero clients.
main: no applications active in common memory

```
-----  
| section corresponding to  
|  
| NOTE # 15  
|  
| Retrieve and display the common  
| memory usage statistics.  
|  
-----
```

A total of 1 mbx's were created. 0 still active.
A total of 2 fsm's were clients of cm. 0 still active.

cm_declare failures: 0

 last by fsm:
 for mbx:
 at:

cm_declare successes: 2
 last by fsm: process_2
 for mbx: proc_1_mbx
 at: Sat Feb 27 19:50:55 1988

cm_undeclare failures: 0
 last by fsm:
 for mbx:
 at:

cm_undeclare successes: 1
 last by fsm: process_2
 for mbx: proc_1_mbx
 at: Sat Feb 27 19:51:35 1988

cm_write failures: 0
 last by fsm:
 for mbx:
 at:

cm_write successes: 5
 last by fsm: process_1
 for mbx: proc_1_mbx
 at: Sat Feb 27 19:51:35 1988

cm_read failures: 0
 last by fsm:
 for mbx:
 at:

```
cm_read successes: 5
  last by fsm: process_2
  for mbx: proc_1_mbx
  at: Sat Feb 27 19:51:35 1988

cm_ckmall failures: 0
  last by fsm:
  for mbx:
  at:

cm_ckmall successes: 9002
  last by fsm: process_2
  for mbx:
  at: Sat Feb 27 19:51:35 1988

cm_disc failures: 0
  last by fsm:
  for mbx:
  at:

cm_disc successes: 1
  last by fsm: process_1
  for mbx:
  at: Sat Feb 27 19:51:36 1988

cm_get_mbx_list failures: 0
  last by fsm:
  for mbx:
  at:

cm_get_mbx_list successes: 2
  last by fsm: process_1
  for mbx:
  at: Sat Feb 27 19:51:35 1988

cm_get_fsm_list failures: 0
  last by fsm:
  for mbx:
  at:

cm_get_fsm_list successes: 3
  last by fsm:
  for mbx:
  at: Sat Feb 27 19:51:35 1988

cm_get_cm_stats failures: 0
  last by fsm:
  for mbx:
  at:
```

```
cm_get_cm_stats successes: 0
  last by fsm:
  for mbx:
  at:

cm_get_fsm_stats failures: 1
  last by fsm: process_3
  for mbx:
  at: Sat Feb 27 19:51:35 1988

cm_get_fsm_stats successes: 1
  last by fsm: process_1
  for mbx:
  at: Sat Feb 27 19:51:35 1988

cm_get_mbx_stats failures: 1
  last by fsm: proc_3_mbx
  for mbx:
  at: Sat Feb 27 19:51:35 1988

cm_get_mbx_stats successes: 1
  last by fsm:
  for mbx: proc_1_mbx
  at: Sat Feb 27 19:51:35 1988
```


APPENDIX E

SOURCE CODE LISTINGS OF THE COMMON MEMORY PROGRAMS

The source code listings of the individual program files comprising the PC common memory have been placed at this appendix in the order shown below.

- E.1. CM_CONST.H
- E.2. CM_GLOBALS.H
- E.3. CM_TYPES.H
- E.4. CM_FUNCS.C
- E.5. CM_UTILS.C
- E.6. SFUNCS.C

Their pages are formatted and numbered differently from the rest of this document in order to provide a reference base for discussions of content.

52 transferred to the user data are. More
53 data is actually available. */

55 /* Errors greater than 0Fx are FATAL errors. This means that the requested
56 action was NOT performed. */

59 /* 0x10 - 0x1F = general errors (or errors common to multiple categories */
60 #define E_CM_FATALERR 0x10 /* defines the start of FATAL ERROR range */
61 #define E_CM_INSUFFMEM 0x10 /* insufficient memory. malloc failed */

64 /* 0x20 - 0x2F = mbx errors */
65 #define E_CM_MBXERR 0x20 /* base value for mbx errors */
66 /* the following 4 error codes must always be related according to */
67 /* <err_code> = E_CM_MBXERR + CM_<type>_ACCESS */
68 #define E_CM_MBXNOREAD 0x21 /* can't have READ - another has XREAD */
69 #define E_CM_MBXNOXREAD 0x22 /* can't have XREAD - another fsm has
70 either READ or XREAD */
71 #define E_CM_MBXNOWRITE 0x24 /* can't have WRITE - another has XWRITE */
72 #define E_CM_MBXNOXWRITE 0x28 /* can't have XWRITE - another fsm has
73 either WRITE or XWRITE */
74 #define E_CM_MBXSIZE 0x29 /* Invalid size, returned if:
75 (1) negative size in cm_declare, or
76 (2) byte count < 1 for cm_write, or
77 (3) size doesn't match previously-
78 declared size (for cm_declare), or
79 (4) attempt to write nr_bytes
80 greater than declared size */
81 #define E_CM_MBXACCESS 0x2a /* Invalid mbx access, returned if:
82 (1) unrecognizable mbx access code
83 supplied for a cm_declare or
84 cm_undeclare, or
85 (2) invalid list_type supplied to
86 cm_get_fsm_list */
87 #define E_CM_MBXACC BOTH 0x2b /* can't declare both READ and XREAD or
88 WRITE and XWRITE in the same mbx declare.
89 However, you can later declare a mbx
90 to be READ after having previously
91 declared it XREAD ... this changes your
92 access and lets other fsm's have
93 access. */
94 #define E_CM_MBXNAME 0x2c /* Invalid mbx name - too long or none given */
95 #define E_CM_MBXNOTDECL 0x2d /* returned if:
96 (1) fsm attempts to undeclare a mbx
97 for which it is not a client for
98 the respective access, or
99 (2) fsm attempts to perform read or
100 write actions w/ a mbx for which
101 it is not a client for the
102 respective access, or

```
103 (3) fsm requests any cm action
104 without being a client of cm, or
105 (4) attempt to get info for a mbx
106 that is not in common memory */

108 /* 0x30 - 0x3F = fsm errors */
109 #define E_CM_FSMERR 0x30 /* base value for fsm errors */
110 #define E_CM_FSMNAME 0x31 /* Invalid name - too long or none given */
111 #define E_CM_FSMNOTINCM 0x32 /* fsm not a common memory client */
```



```
1 /* cm_globals.h - contains all global variable declarations */

3 #include <stdio.h>
4 #include "cm_const.h"           /* Include cm_const.h before cm_types.h */
5 #include "cm_types.h"

7 int cm_mbxhandle = 0;          /* used by CMM in assigning mbxhandles */

9 fsm_rec *cm_fsm_list = NULL;    /* list of active fsm's maintained by CMM */
10 mbx_rec *cm_mbx_list = NULL;    /* list of declared mbx's maintained by CMM */

13 /* The following are initialized in function cm_ini found in cm_utils.h.
14   The first 2 are EXTERN so they can be adjusted by the user application pgm.
15   For testing and development purposes, they are currently initialized in
16   routine cm_ini. */

18 int CM_DEBUG_LEVEL = -1;        /* causes debugging statements to be displayed */
19 boolean CM_GET_STATS = TRUE;    /* tells cmm to gather some performance stats */
20 cm_activity_stats *cm_activity = NULL; /* tracks useage of cm routines */
21 cm_client_stats *cm_clients = NULL; /* keeps count of mbx's and fsm's */
```



```
1 /* cm_types.h - definitions of types used in the common memory and
2      interface routines.

4      The source program that includes this file must also
5      include file "cm_const.h".

7      #####
8      | library version 1.0a | <-- chg here and in cm_const.h
9      #####

11 */
12 #ifndef boolean
13     typedef int boolean;
14 #endif

18 /* some standardized definitions */
19 typedef unsigned char byte;
20 typedef long int timestamp;           /* nr of seconds since 01 Jan '70 */

24 /* this is how mailbox variables are stored */
25 typedef struct mbx_stats {
26     int nr_fsms;                      /* nr of subscribers to this mbx */
27     unsigned int nr_accesses;          /* nr of times mbx accessed */
28     timestamp when;                 /* time of last access */
29     struct fsm_rec_type *who;        /* pts to last accessing fsm */
30 };
31 typedef struct client_chain {
32     struct fsm_rec_type *who;        /* ptr to subscribing fsm record */
33     boolean exclusive;              /* TRUE if fsm has XREAD or XWRITE */
34     struct client_chain *next;       /* ptr to next fsm or NULL */
35 };
36 typedef struct mbx_rec_type {
37     char mbxname [CM_MAXMBXNAMELENGTH];
38     int handle;                     /* short (numeric) name for mbx */
39     int declaredlength;             /* nr bytes declared */
40     int msglength;                 /* nr bytes currently stored */
41     byte *data;                   /* ptr to actual data area */
42     struct mbx_stats readers;      /* reader statistics */
43     struct mbx_stats writers;      /* writer statistics */
44     struct client_chain *readerlist; /* linked list of fsm_rec ptrs */
45     struct client_chain *writerlist; /* linked list of fsm_rec ptrs */
46     struct mbx_rec_type *next;      /* ptr to next mbx list entry */
47 }mbx_rec;

49 /* this is how fsm variables are stored */
50 typedef struct mbx_decl_chain {
51     struct mbx_rec_type *mbx;        /* ptr to subscribed mbx record */
```

```
52     struct mbx_decl_chain *next;           /* ptr to next mbx or NULL */
53 };
54 typedef struct fsm_stats {
55     int nr_mbxes;                         /* nr of mbxes declared for "this" access */
56     unsigned int nr_accesses;              /* ttl nr of accesses */
57     timestamp when;                      /* time of last access */
58     int mbxhandle;                       /* last mbx accessed */
59     struct mbx_decl_chain *mbx_llist;    /* chain of declared mailboxes for "this" access */
60 };
61 typedef struct update_llist {
62     int mbxhandle;                         /* Id's mbx that has been changed */
63     struct update_llist *next;             /* pts to next llist entry */
64 };
65 typedef struct fsm_rec_type {
66     char fsmname[CM_MAXFSMNAMELENGTH];   /* name, null terminated */
67     struct fsm_stats read;               /* READ stats */
68     struct fsm_stats write;              /* WRITE stats */
69     struct update_llist *update_top;     /* FIFO llist of subscribed READ mbxes
70                                     that have been updated. CMM will FREE them after fsm reads
71                                     the mbxes. Or, fsm can be handed this llist w/ CM_SYNC call,
72                                     and fsm must FREE the llist when done. If a subsequent update
73                                     is made to a mbx that is already on the llist, no additional
74                                     llist entry will be submitted. */
75     struct update_llist *update_bot;     /* position of the update llist where
76                                     new entries are made to maintain
77                                     the FIFO order. */
78     int nr_updates;                      /* nr of update entries on the llist */
79     struct fsm_rec_type *next;           /* ptr to next fsm record or NULL */
80 } fsm_rec;

84 /* this is how the cmm stores statistics */
85 typedef struct {
86     char fsm[CM_MAXFSMNAMELENGTH];       /* name of fsm */
87     char mbx[CM_MAXMBXNAMELENGTH];       /* name of mbx, if appropriate */
88     timestamp when;                     /* nr times this service called */
89     unsigned int nr_times;
90 } base_stats;

92 typedef struct {
93     base_stats success;                /* for when the call succeeded */
94     base_stats failure;               /* for when it failed */
95 } group_stats;

97 typedef struct {
98     group_stats cm_declare,
99                 cm_undeclare,
100                cm_write,
101                cm_read,
102                cm_ckmail,
```



```
154     int read_fsms;          /* nr of READ subscribers */  
155     unsigned int read_accesses; /* nr of READ accesses */  
156     timestamp read_time;    /* time of last read */  
157     char reader[CM_MAXFSMNAMELENGTH]; /* name of last reader fsm */  
158     int write_fsms;          /* nr of WRITE subscribers */  
159     unsigned int write_accesses; /* nr of WRITE accesses */  
160     timestamp write_time;    /* time of last write */  
161     char writer[CM_MAXFSMNAMELENGTH]; /* name of last writer fsm */  
162     char xreader[CM_MAXFSMNAMELENGTH]; /* name of xreader fsm */  
163     char xwriter[CM_MAXFSMNAMELENGTH]; /* name of xwriter fsm */  
164 }cm_mbx_stats_rec;
```

1 /* cm_funcs.c - contains the common memory interface routines
2 used by client processes.

5 The following section describes the purpose of each interface function.
6 The argument list for each function is described in detail. Each
7 function returns an integer status value that correlates with what the
8 function is to perform (hence the 'int' before each function name).

10 The list of all potential status values that this family of functions
11 can return, and their significance, is provided in file `cm const.h`

PURPOSE: Routine creates the necessary fsm, mbx, and mbx client structures within common memory to support future mailbox manipulations by the declaring fsm.

43 NOTE: When you cm_declare to an existing mbx for READ or XREAD,
44 the cmm will NOT place an entry into your "update list" for
45 that mbx. The purpose of the update list is to indicate that
46 the mbx contents have been written SINCE the time of your
47 cm_declare or cm_read. It is assumed that you will perform an
48 initial cm read as a matter of course (if desired).

50 RETURNS: status, as id'd in cm const.h

```
52 int cm_undeclare(fsm
53
54
55
56
57     mbxaccess - INPUT
58             TYPE char *fsm
59             fsm name string. String must
60             be null-terminated, and must be <= 32
61             characters in length (excluding NULL).
62
63     mbxhandle - INPUT
64             TYPE int *mbxhandle
65             Variable value was initially set by
66             cm_declare and is used as a fast way to
67             reference a specific mbx. Although this does
68             not need to be a pointer, it is specified as
69             such for compatibility with cm_declare
70             (which requires it) and other common memory
71             routines.
72
73 PURPOSE: Routine is used to remove an fsm from a particular mbx's
74 client list for the specified access. More than
75 one access type may be specified at each call, subject to
76 the access rules identified in the cm_declare section.
77 If the action results in a mbx without any clients, the
78 mailbox is deleted and the space returned to the operating
79 system. Likewise, if the action results in an fsm that has
80 no other mbx's declared, that fsm is removed as a cm client.
81 If the mbx was identified on the update list, then the
82 update list entry will be purged.

84 RETURNS: status, as defined in cm_const.h
```

```
87 int cm_write (fsm
88
89
90
91
92     mbxhandle - INPUT
93             TYPE int *mbxhandle
94             Variable value was initially set by
95             cm_declare and is used as a fast way to
96             reference a specific mbx. Although this does
97             not need to be a pointer, it is specified as
98             such for compatibility with cm_declare
99             (which requires it) and other common memory
100            routines.
101
102     usr_data - INPUT
103             TYPE byte *usr_data
```

```
121 Int cm_read (fsm
122
123
124
125
126     mbxhandle
127
128
129
130
131
132
133
134
135     usr_data
136
137
138
139     nr_bytes
140
141
142
143
144
145
146
147
148
149
150
151
152
153
- INPUT
  TYPE char *fsm
  fsm name string. String must
  be null-terminated, and must be <= 32
  characters in length (excluding NULL).
- INPUT
  TYPE int *mbxhandle
  Variable value was initially set by
  cm_declare and is used as a fast way to
  reference a specific mbx. Although this does
  not need to be a pointer, it is specified as
  such for compatibility with cm_declare
  (which requires it) and other common memory
  routines.
- INPUT
  TYPE byte *usr_data
  points to user data area to which
  bytes are to be transferred.
- INPUT/OUTPUT
  TYPE Int *nrbytes
  When cm_read is called, if :
    (1) the Int variable = 0, then all data
        bytes are transferred from the mailbox
        to the user's data area.
    (2) the Int variable is not equal to 0,
        the nr of bytes transferred will be
        min(nr_bytes, nr_bytes_in_mbx).
  Upon return, the variable pointed to by
  nr_bytes will contain the actual number of
  bytes transferred. If fewer bytes are
  transferred to the user data area than are
  available in the mailbox, an "information-
  only" status of 1 CM MOREDATA is returned
```

154 to alert the user who may have inadvertently
155 called `cm_read` without clearing the variable
156 pointed to by `nr_bytes`.

158 It is the user's responsibility to make
159 sure that the data area is large enough
160 to contain the mailgram.

161)

162 PURPOSE: Routine is used to transfer the specified number of bytes
163 to the user data area from the common memory mailbox.
164 It is recommended that `cm_ckmail` be used together with
165 `cm_read` to minimize common memory accesses.

167 If an entry for this mbx exists on the update list of this fsm,
168 it is removed at completion of the `cm_read` operation.

170 RETURNS: status, as defined in `cm_const.h`

174 int `cm_ckmail`(fsm
175 - INPUT
176 TYPE char *fsm
177 fsm name string. String must
178 be null-terminated, and must be <= 32
179 characters in length (excluding NULL).
180 list_ptr - INPUT
181 TYPE struct update_list **list_ptr;
182 If an update list exists for this fsm, `cm_ckmail`
183 will return a ptr to the top of the
184 update list in this location. If none
185 exists, `cm_ckmail` will return NULL.
186 nr_entries - INPUT
187 TYPE int *nr_entries;
188 If an update list exists, the int variable
189 will contain the number of entries in the
190 update list; else, it will contain ZERO.
191)

192 PURPOSE: For each fsm that is a READER client of common memory, the
193 common memory manager creates and maintains a list of those
194 mailboxes that have changed since the last time the fsm read
195 them (ie, an update list). Whenever an fsm writes to a
196 common memory mbx, the common memory manager posts an entry
197 on this "update" list. If an entry already exists for a
198 changed mbx, no additional entry is made. The list is
199 maintained in first-in-first-out (FIFO) order.

201 Entries are removed from this list whenever an fsm calls
202 `cm_read` for the respective mailbox. Alternately, an fsm may
203 call `cm_ckmail`. If an update list exists, `cm_ckmail` returns a
204 pointer to the top of the list and releases the list to the

205 fsm. If no update list exists, cm_ckmail returns NULL. (The
206 common memory manager will start a new list when the next
207 mbx update arrives.)

209 Once the list is released to the fsm, it is the responsibility
210 of the fsm to FREE the memory allocated for the list. The
211 fsm may do so itself, or it may call cm_free_update_list,
212 passing it a pointer to the top of the list.

214 Using the update list, the fsm can now perform sequential
215 cm_read operations and only access those mailboxes that have
216 changed since the last read operation.

218 RETURNS: status, as id'd in cm_const.h

221 int cm_disc (fsm
222 - INPUT
223 TYPE char *fsm
224 fsm name string. String must
225 be null-terminated, and must be <= 32
226 characters in length (excluding NULL).
)

228 PURPOSE: Provides a short-cut method for a process to undeclare all of
229 its mailboxes at one time. All data structures within common
230 memory that are associated with that fsm are freed. The fsm
231 must issue a cm_declare before it can again access common
232 memory variables.

234 RETURNS: status, as id'd in cm_const.h

238 int cm_get_fsm_list(
239 mbxname - INPUT
240 TYPE char *mbxname
241 if NULL, this routine will return, through
242 fsm_list_ptr, the list of all fsm names
243 known to the common memory manager.
244 Arg "list_type" has no effect.
245 If not NULL, it must point to a mbx name.
246 This routine will return a list of all fsm's
247 that are a client of that specific mbx. The
248 arg "list_type" is used to qualify whether
249 the caller wants the list of READER clients
250 or the list of WRITER clients.
251 list_type - INPUT
252 TYPE char
253 May only have the values 'R' (READ) and
254 'W' (WRITE) when *mbxname is non-NULL. If
255 *mbxname is NULL, list_type is ignored.

256 fsm_llist_ptr - OUTPUT
257 TYPE struct fsm_llist_type **fsm_llist_ptr
258 This routine will create a linked list of fsm
259 names and return a ptr to the top of the
260 list if any fsm's exists, or NULL if none
261 exist. It is the user's responsibility
262 to free this list when it is no longer needed.
263 int_ptr - OUTPUT
264 TYPE int *int_ptr
265 Upon return, the int variable will equal the
266 number of entries in the list.
267)

269 PURPOSE: This routine allows any fsm to determine what fsm's are
270 currently active in common memory. Using the mbxname and
271 list_type appropriately, the caller can retrieve the list of
272 all fsm's, or only the list of clients (read or write) for a
273 specific mbx. The information that is returned on the list
274 can be used to solicit other fsm (and, indirectly, mbx)
275 statistics.

277 RETURNS: status, as id'd in cm_const.h

281 int cm_get_mbx_llist(
282 fsmname - INPUT
283 TYPE char *fsmname
284 fsm name string. String must
285 be null-terminated, and must be <= 32
286 characters in length (excluding NULL).
287 If NULL, this routine will return, through
288 mbx_llist_ptr, the list of all mbx names
289 known to the common memory manager.
290 Arg "list_type" has no effect.
291 If not NULL, it must point to a fsm name.
292 This routine will return a list of all mbx's
293 that are declared by that fsm. The
294 arg "list_type" is used to qualify whether
295 the caller wants the list of READER mbx's
296 or the list of WRITER mbx's.
297 list_type - INPUT
298 TYPE char
299 May only have the values 'R' (READ) and
300 'W' (WRITE) when *fsmname is non-NULL. If
301 *fsmname is NULL, list_type is ignored.
302 mbx_llist_ptr - OUTPUT
303 TYPE struct mbx_llist_type **mbx_llist_ptr
304 This routine will create a linked list of mbx
305 names and return a ptr to the top of the
306 list if any mbx's exists, or NULL if none

307 exist. It is the user's responsibility
308 to free this list when it is no longer needed.
309 int_ptr - INPUT
310 TYPE int *int_ptr
311 Upon return, the int variable will equal the
312 number of entries in the list.
313)

315 PURPOSE: This routine allows any fsm to determine what mbx's are
316 currently active in common memory. The information that is
317 returned on the list can be used to solicit other cm fsm
318 (and, indirectly, mbx) statistics. Using the fsmname and
319 list_type appropriately, the caller can retrieve the list of
320 all mbx's in common memory, or only the list of mbx's (read
321 or write) for a specific fsm.

323 RETURNS: status, as id'd in cm_const.h

327 int cm_get_cm_stats(
328 activity_ptr - INPUT
329 TYPE cm_activity_stats *activity_ptr
330 Points to user-allocated data area of
331 appropriate size. This routine will copy
332 the activity statistics into that data area.
333 It has been implemented in this fashion to
334 minimize malloc and free operations, since it
335 is assumed the user will want this information
336 more than once.

338 client_ptr - INPUT
339 TYPE cm_client_stats *client_ptr;
340 Points to user-allocated data area of
341 appropriate size. This routine will copy
342 the client statistics into that data area.
343 It has been implemented in this fashion to
344 minimize malloc and free operations, since it
345 is assumed the user will want this information
346 more than once.
347)

349 PURPOSE: This routine provides cm operating statistics.
350 Since the size of the statistics areas is static,
351 the user must provide pointers to space in the
352 data area into which the statistics will be
353 copied. This avoids malloc overhead in case the
354 user wishes to call this routine multiple times.

356 RETURNS: status, as id'd in cm_const.h

```
360 int cm_get_fsm_stats(  
361         fsm           - INPUT  
362                     TYPE char *fsm  
363                     fsm name string. String must  
364                     be null-terminated, and must be <= 32  
365                     characters in length (excluding NULL).  
366         ptr           - INPUT  
367                     TYPE cm_fsm_stats_rec *ptr  
368     )
```

370 PURPOSE: This routine returns the common memory statistics
371 for the specified fsm, as identified in cm_fsm_stats_rec.
372 Since the size of the statistics area is static,
373 the user must provide a pointer to space in the
374 user data area into which the statistics will be
375 copied. This avoids malloc overhead in case the
376 user wishes to call this routine multiple times.

378 RETURNS: status, as id'd in cm_const.h

```
382 int cm_get_mbx_stats(  
383         mbx           - INPUT  
384                     TYPE char *mbx  
385                     mbx name string. String must  
386                     be null-terminated, and must be <= 32  
387                     characters in length (excluding NULL).  
388         ptr           - INPUT  
389                     TYPE cm_mbx_stats_rec *ptr  
390     )
```

392 PURPOSE: This routine returns the common memory statistics
393 for the specified mbx, as identified in cm_mbx_stats_rec.
394 Since the size of the statistics area is static,
395 the user must provide a pointer to space in the
396 user data area into which the statistics will be
397 copied. This avoids malloc overhead in case the
398 user wishes to call this routine multiple times.

400 RETURNS: status, as id'd in cm_const.h

403 */

405 #include "cm_utils.c"

408 /* ***** Declare a mailbox in common memory. This

```
409 * | cm_declare | can result in a connection to an already
410 * |::::::::::::::::::::::::::| existing mbx, or creation of a new one.
411 */
412 int cm_declare (fsm, mbxname, mbxsize, mbxaccess, mbxhandle)
413 char *fsm, *mbxname;
414 int mbxsize, mbxaccess, *mbxhandle;
415 { int return_status;
416 byte *data_ptr;
417 fsm_rec *tmp_fsm_rec;
418 mbx_rec *tmp_mbx_rec;
419 struct client_chain *tmp_client_rec;
420 boolean new_fsm, new_mbx;
421
422 if (lcm_activity) cm_init(); /* init cm structures */
423 new_fsm = FALSE;
424 new_mbx = FALSE;
425
426 if (return_status = cm_validate_fsm(fsm)) {
427 log_status(&cm_activity->cm_declare.failure, NULL, NULL);
428 return(return_status); /* fsm name error */
429 }
430 if (return_status = cm_validate_mbx(mbxname, mbxsize)) {
431 log_status(&cm_activity->cm_declare.failure, fsm, NULL);
432 return(return_status); /* mbx error */
433 }
434 if (return_status = cm_validate_access(mbxaccess)) {
435 log_status(&cm_activity->cm_declare.failure, fsm, mbxname);
436 return(return_status); /* invalid access */
437 }
438
439 if (tmp_fsm_rec = cm_fsm_find(fsm)) {
440 eprintf(9, "cm_declare: fsm already known to common memory\n");
441 } else { /* init new rec for this fsm */
442 eprintf(6, "cm_declare: add '%s' as cm client\n", fsm);
443 tmp_fsm_rec = (fsm_rec *) malloc (sizeof(fsm_rec));
444 if (tmp_fsm_rec == NULL) {
445 log_status(&cm_activity->cm_declare.failure, fsm, mbxname);
446 return(E_CM_INSUFFMEM);
447 }
448 bclr(tmp_fsm_rec, sizeof(fsm_rec));
449 strcpy(tmp_fsm_rec->fsmname, fsm); /* save fsm name */
450 new_fsm = TRUE;
451 }
452
453 if (tmp_mbx_rec = cm_mbx_find(mbxname)){ /* ck if mbx already exists */
454 eprintf(9, "cm_declare: This mbx is already known to common memory\n");
455 *mbxhandle = tmp_mbx_rec->handle;
456 } else { /* init new rec for this mbx */
457 eprintf(6, "cm_declare: add '%s' as new cm mbx\n", mbxname);
458 tmp_mbx_rec = (mbx_rec *) malloc (sizeof(mbx_rec));
459 if (tmp_mbx_rec == NULL) {
```

```
460     log_status(&cm_activity->cm_declare.failure,fsm,mbxname);
461     return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec,E_CM_INSUFFMEM));
462 }
463 bcir(tmp_mbx_rec,sizeof(mbx_rec));
464 strcpy(tmp_mbx_rec->mbxname,mbxname); /* save fsm name */
465 *mbxhandle = ++cm_mbxhandle;
466 tmp_mbx_rec->handle = cm_mbxhandle;
467 tmp_mbx_rec->declaredlength = mbxsize;
468 new_mbx = TRUE;
469 data_ptr = (byte *) malloc (mbxsize); /* spc available for data mbx ? */
470 if (data_ptr == NULL) {
471     log_status(&cm_activity->cm_declare.failure,fsm,mbxname);
472     return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec,E_CM_INSUFFMEM));
473 }
474 bcir(data_ptr,mbxsize);
475 tmp_mbx_rec->data = data_ptr;
476 }

477 if (mbxaccess & CM_XREAD_ACCESS) { /* ck if another fsm has READ or XREAD */
478     if ((tmp_mbx_rec->readers.nr_fsms > 1) /* DIE If more than 1 reader */
479         || ((tmp_mbx_rec->readerlist) /* DIE If I'm not the only reader */
480             && (find_mbx_client(tmp_mbx_rec->readerlist, tmp_fsm_rec) == NULL))) {
481         log_status(&cm_activity->cm_declare.failure,fsm,mbxname);
482         return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec,E_CM_MBXNOREAD));
483     }else
484         if (tmp_client_rec = find_mbx_client(tmp_mbx_rec->readerlist, tmp_fsm_rec)){
485             tmp_client_rec->exclusive = TRUE; /* reset from READ to XREAD */
486             return_status = I_CM_DUPMBX; /* let user know he was already a client for
487                                         * this mbx w/ either READ/XREAD access */
488         }else {
489             ++(tmp_mbx_rec->readers.nr_fsms); /* NEW, so Increment nr of readers */
490             ++(tmp_fsm_rec->read.nr_mbxes); /* incr nr of mbxes this fsm reads */
491             return_status = add_mbx_client (tmp_fsm_rec, &tmp_mbx_rec->readerlist, TRUE);
492             if (return_status >= E_CM_FATALERR)
493                 return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec,return_status));
494             return_status = add_fsm_mbx (&tmp_fsm_rec->read.mbx_list, tmp_mbx_rec); /* add mbx to fsm's list */
495             if (return_status >= E_CM_FATALERR) {
496                 log_status(&cm_activity->cm_declare.failure,fsm,mbxname);
497                 return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec,return_status));
498             }
499         }
500     }
501 }else
502     if (mbxaccess & CM_READ_ACCESS) /* ck if another fsm has XREAD */
503         if ((tmp_mbx_rec->readers.nr_fsms == 1) /* DIE If someone has XREAD */
504             && ((tmp_mbx_rec->readerlist->exclusive) /* and It Isn't ME */
505                 && (tmp_mbx_rec->readerlist->who != tmp_fsm_rec))) {
506             log_status(&cm_activity->cm_declare.failure,fsm,mbxname);
507             return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec,E_CM_MBXNOREAD));
508         }else
509             if (tmp_client_rec = find_mbx_client(tmp_mbx_rec->readerlist, tmp_fsm_rec)){
510                 tmp_client_rec->exclusive = FALSE; /* reset from XREAD to READ */
```

```
511     return_status = I_CM_DUPMBX; /* let user know he was already a client for
512                                /* this mbx w/ either READ/XREAD access */
513
514     }else {
515         ++(tmp_mbx_rec->writers.nr_fsms); /* NEW, so increment nr of writers */
516         ++(tmp_fsm_rec->read.nr_mbxes); /* incr nr of mbxes this fsm reads */
517         return_status = add_mbx_client (tmp_fsm_rec, &tmp_mbx_rec->writerlist, FALSE);
518         if (return_status >= E_CM_FATALERR) {
519             log_status(&cm_activity->cm_declare.failure, fsm, mbxname);
520             return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec, return_status));
521         }
522         return_status = add_fsm_mbx (&tmp_fsm_rec->read.mbx_list, tmp_mbx_rec); /* add mbx to fsm's list */
523         if (return_status >= E_CM_FATALERR) {
524             log_status(&cm_activity->cm_declare.failure, fsm, mbxname);
525             return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec, return_status));
526         }
527     }
528
529     if (mbxaccess & CM_XWRITE_ACCESS) /* ck if another fsm has WRITE or XWRITE */
530     if ((tmp_mbx_rec->writers.nr_fsms > 1) /* DIE if more than 1 writer */
531         || ((tmp_mbx_rec->writerlist) /* DIE if I'm not the only writer */
532             && (find_mbx_client(tmp_mbx_rec->writerlist, tmp_fsm_rec) == NULL))) {
533         log_status(&cm_activity->cm_declare.failure, fsm, mbxname);
534         return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec, E_CM_MBXNOXWRITE));
535     }else /* allowed to have XWRITE */
536     if (tmp_client_rec = find_mbx_client(tmp_mbx_rec->writerlist, tmp_fsm_rec)) {
537         tmp_client_rec->exclusive = TRUE; /* reset from WRITE to XWRITE */
538         return_status = I_CM_DUPMBX; /* let user know he was already a client for
539                                * this mbx w/ either WRITE/XWRITE access */
540     }else {
541         ++(tmp_mbx_rec->writers.nr_fsms); /* NEW, so increment nr of writers */
542         ++(tmp_fsm_rec->write.nr_mbxes); /* incr nr of mbxes this fsm writes to */
543         return_status = add_mbx_client (tmp_fsm_rec, &tmp_mbx_rec->writerlist, TRUE);
544         if (return_status >= E_CM_FATALERR) {
545             log_status(&cm_activity->cm_declare.failure, fsm, mbxname);
546             return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec, return_status));
547         }
548         return_status = add_fsm_mbx (&tmp_fsm_rec->write.mbx_list, tmp_mbx_rec); /* add mbx to fsm's list */
549         if (return_status >= E_CM_FATALERR) {
550             log_status(&cm_activity->cm_declare.failure, fsm, mbxname);
551             return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec, return_status));
552         }
553     }
554 }else
555     if (mbxaccess & CM_WRITE_ACCESS) /* ck if other fsm has XWRITE */
556     if ((tmp_mbx_rec->writers.nr_fsms == 1) /* DIE if someone has XWRITE */
557         && ((tmp_mbx_rec->writerlist->exclusive) /* and it isn't ME */
558             && (tmp_mbx_rec->writerlist->who != tmp_fsm_rec))) {
559         log_status(&cm_activity->cm_declare.failure, fsm, mbxname);
560         return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec, E_CM_MBXNOWRITE));
561     }else /* only add client if not already on list */
562
563 }
```

```
562     if (tmp_client_rec = find_mbx_client(tmp_mbx_rec->writerlist, tmp_fsm_rec)){
563         tmp_client_rec->exclusive = FALSE; /* reset from XWRITE to WRITE */
564         return_status = I_CM_DUPMBX;          /* let user know he was already a client for
565                                         * this mbx w/ either WRITE/XWRITE access */
566     }else {
567         ++(tmp_mbx_rec->writers.nr_fsms); /* NEW, so increment nr of writers */
568         ++(tmp_fsm_rec->write.nr_mbxes); /* Incr nr of mbxes this fsm writes to */
569         return_status = add_mbx_client (tmp_fsm_rec, &tmp_mbx_rec->writerlist, FALSE);
570         if (return_status >= E_CM_FATALERR) {
571             log_status(&cm_activity->cm_declare.failure, fsm, mbxname);
572             return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec, return_status));
573         }
574         return_status = add_fsm_mbx (&tmp_fsm_rec->write.mbx_list, tmp_mbx_rec); /* add mbx to fsm's list
575         if (return_status >= E_CM_FATALERR) {
576             log_status(&cm_activity->cm_declare.failure, fsm, mbxname);
577             return(clear_declare(new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec, return_status));
578         }
579     }
580
582     if (new_fsm)                                /* Insert new fsm onto list */
583         if (cm_fsm_list == NULL)
584             cm_fsm_list = tmp_fsm_rec;
585         else {
586             tmp_fsm_rec->next = cm_fsm_list;      /* Insert at front of list */
587             cm_fsm_list = tmp_fsm_rec;
588         }
589
590     if (new_mbx)                                /* Insert new mbx onto list */
591         if (cm_mbx_list == NULL)
592             cm_mbx_list = tmp_mbx_rec;
593         else {
594             tmp_mbx_rec->next = cm_mbx_list;      /* Insert at front of list */
595             cm_mbx_list = tmp_mbx_rec;
596         }
597     log_status(&cm_activity->cm_declare.success, fsm, mbxname);
598     if (new_mbx) {
599         ++(cm_clients->mbx_ttl);
600         ++(cm_clients->mbx_active);
601     }
602     if (new_fsm) {
603         ++(cm_clients->fsm_ttl);
604         ++(cm_clients->fsm_active);
605     }
606     return (return_status);
607 }
610 /* ////////////////////////////// Remove a client from a mbx client list, but
611 * | cm_undeclare | only for the specified access.
612 * //////////////////////////////
```

```
613 */
614 int cm_undeclare (fsm, mbxaccess, mbxhandle)
615 char *fsm;
616 int mbxaccess, *mbxhandle;
617 { int l, return_status;
618 fsm_rec *tmp_fsm_rec;
619 mbx_rec *tmp_mbx_rec;
620 struct client_chain *tmp_client_rec;

622 eprintf(9,"cm_undeclare: fsm %s, access %d, mbxhandle %d\n",
623 fsm, mbxaccess, *mbxhandle);
624 if (!cm_activity) cm_init(); /* init cm structures */
625 if ((!(tmp_fsm_rec = cm_fsm_find(fsm))) /* fsm active in cm ? */
626 || (!(tmp_mbx_rec = cm_mbxhandle_find(*mbxhandle)))) /* mbxhandle in cm ? */
627 log_status(&cm_activity->cm_undeclare.failure,fsm,NULL,NULL);
628 return(E_CM_MBXNOTDECL);
629 }
630 if (return_status = cm_validate_access(mbxaccess)) {
631 log_status(&cm_activity->cm_undeclare.failure,fsm,NULL);
632 return(return_status); /* invalid access */
633 }

635 eprintf(9,"cm_undeclare: passed validation tests\n");
636 if (mbxaccess & (CM_READ_ACCESS | CM_XREAD_ACCESS)) /* is fsm a READ client ? */
637 if (tmp_client_rec = find_mbx_client(tmp_mbx_rec->readerlist,tmp_fsm_rec)) {
638 eprintf(9,"cm_undeclare: found client in readerlist\n");
639 if ((mbxaccess & CM_READ_ACCESS) && (tmp_client_rec->exclusive)) {
640 log_status(&cm_activity->cm_undeclare.failure,fsm,tmp_mbx_rec->mbxname);
641 return(E_CM_MBXNOTDECL); /* DIE if undeclaring READ but have XREAD */
642 }
643 if ((mbxaccess & CM_XREAD_ACCESS) && (!tmp_client_rec->exclusive)) {
644 log_status(&cm_activity->cm_undeclare.failure,fsm,tmp_mbx_rec->mbxname);
645 return(E_CM_MBXNOTDECL); /* DIE if undeclaring XREAD but have READ */
646 }
647 del_update_rec(tmp_fsm_rec,tmp_mbx_rec->handle); /* remove update notification, if on list */
648 del_mbx_client(tmp_client_rec,&(tmp_mbx_rec->readerlist),&(tmp_mbx_rec->readers.nr_fsms));
649 del_fsm_mbx_entry(&tmp_fsm_rec->read.mbx_list,tmp_mbx_rec);
650 -(tmp_fsm_rec->read.nr_mbxes); /* decrease nr of mbxes this fsm reads */
651 eprintf(7,"cm_undeclare: deleted READ/XREAD entry\n");
652 }else {
653 log_status(&cm_activity->cm_undeclare.failure,fsm,tmp_mbx_rec->mbxname);
654 return(E_CM_MBXNOTDECL);
655 }

657 eprintf(9,"cm_undeclare: finished READ/XREAD checks\n");
658 if (mbxaccess & (CM_WRITE_ACCESS | CM_XWRITE_ACCESS)) /* is fsm a WRITE client ? */
659 if (tmp_client_rec = find_mbx_client(tmp_mbx_rec->writerlist,tmp_fsm_rec)) {
660 eprintf(9,"cm_undeclare: found client in writerlist\n");
661 if ((mbxaccess & CM_WRITE_ACCESS) && (tmp_client_rec->exclusive)) {
662 log_status(&cm_activity->cm_undeclare.failure,fsm,tmp_mbx_rec->mbxname);
663 return(E_CM_MBXNOTDECL); /* DIE if undeclaring WRITE but have XWRITE */
664 }
```

```

664     }
665     if ((mbxaccess & CM_XWRITE_ACCESS) && (tmp_client_rec->exclusive)) {
666         log_status(&cm_activity->cm_undeclare.failure, fsm, tmp_mbx_rec->mbxname);
667         return(E_CM_MBXNOTDECL);           /* DIE if undeclaring XWRITE but have WRITE */
668     }
669     del_mbx_client(tmp_client_rec, &(tmp_mbx_rec->writerlist), &(tmp_mbx_rec->writers.nr_fsms));
670     del_fsm_mbx_entry(&tmp_fsm_rec->write.mbx_list, tmp_mbx_rec);
671     --(tmp_fsm_rec->write.nr_mbxes);      /* decrease nr of mbxes this fsm writes */
672     eprintf(7, "cm_undeclare: deleted WRITE/XWRITE entry\n");
673 }else {
674     log_status(&cm_activity->cm_undeclare.failure, fsm, tmp_mbx_rec->mbxname);
675     return(E_CM_MBXNOTDECL);
676 }

678 /* If mbx has no other clients, remove it from cm_mbx_list */
679 if ((tmp_mbx_rec->readerlist) && (tmp_mbx_rec->writerlist)) {
680     --(cm_clients->mbx_active);
681     cm_mbx_del(tmp_mbx_rec);
682 }

684 /* If the fsm has no other mailboxes declared, remove it from cm_fsm_list */
685 if ((tmp_fsm_rec->read.mbx_list) && (tmp_fsm_rec->write.mbx_list)) {
686     --(cm_clients->fsm_active);
687     cm_fsm_del(tmp_fsm_rec);
688 }
689 log_status(&cm_activity->cm_undeclare.success, fsm, tmp_mbx_rec->mbxname);
690 return(I_CM_OK);
691 }

694 /* ===== Transfer the specified number of bytes from the
695 * | cm_write | the user data area to the common memory mbx.
696 * ===== */
697 */
698 int cm_write (fsm, mbxhandle, usr_data, nrbytes)
699 char *fsm;
700 byte *usr_data;
701 int *mbxhandle, *nrbytes;
702 { fsm_rec *tmp_fsm_rec;
703     mbx_rec *tmp_mbx_rec;
704     timestamp when;

706     eprintf(5, "cm_write: at entry, fsm = %s, mbxhandle = %d, nrbytes = %d\n",
707             fsm, *mbxhandle, *nrbytes);

709     if (!cm_activity) cm_init();           /* Init cm structures */
710     if ((!(tmp_fsm_rec = cm_fsm_find(fsm))) /* fsm active in cm ? */
711         || (!(tmp_mbx_rec = cm_mbxhandle_find(*mbxhandle)))) /* mbxhandle in cm ? */
712         log_status(&cm_activity->cm_write.failure, NULL, NULL);
713     return(E_CM_MBXNOTDECL);
714 }

```

```
715 if (!find_mbx_client(tmp_mbx_rec->writerlist,tmp_fsm_rec)) {
716     log_status(&cm_activity->cm_write.failure,fsm,tmp_mbx_rec->mbxname);
717     return(E_CM_MBXNOTDECL);           /* not a WRITE/XWRITE client of this mbx */
718 }
719 if (*nrbytes > tmp_mbx_rec->declaredlength) {
720     log_status(&cm_activity->cm_write.failure,fsm,tmp_mbx_rec->mbxname);
721     return(E_CM_MBXSIZE);           /* attempt to write more than mbx space allows */
722 }
723 eprintf(5,"cm_write: passed validation tests\n");
724 if (*nrbytes)                  /* transfer data to mbx */
725     bcpys(usr_data,tmp_mbx_rec->data,*nrbytes);
726     tmp_mbx_rec->msglength = *nrbytes;           /* nr bytes in msg */
727     ++(tmp_mbx_rec->writers.nr_accesses);
728     tmp_mbx_rec->writers.who = tmp_fsm_rec;
729     time(&when);           /* timestamp */
730     tmp_mbx_rec->writers.when = when;
731     tmp_fsm_rec->write.when = when;
732     ++(tmp_fsm_rec->write.nr_accesses);
733     tmp_fsm_rec->write.mbxhandle = *mbxhandle;

735 /* notify all READ clients of this mbx that a change has occurred in the mbx */
736 notify_clients(tmp_mbx_rec->readerlist,*mbxhandle);
737 log_status(&cm_activity->cm_write.success,fsm,tmp_mbx_rec->mbxname);
738 return(I_CM_OK);
739 }
```

```
743 /* ********************| Transfer the specified number of bytes from the
744 * |      cm_read      | the common memory mbx to the user data area.
745 * | ********************|
746 */
747 Int cm_read (fsm, mbxhandle, usr_data, nrbytes)
748 char *fsm;
749 byte *usr_data;
750 Int *mbxhandle, *nrbytes;
751 { fsm_rec *tmp_fsm_rec;
752     mbx_rec *tmp_mbx_rec;
753     timestamp when;
754     Int return_status;

756 eprintf(5,"cm_read: at entry, fsm = %s, mbxhandle = %d, nrbytes = %d\n",
757         fsm, *mbxhandle, *nrbytes);
758 if (!cm_activity) cm_init();           /* Initialize cmm structures */
759 return_status = I_CM_OK;           /* assume clean exit */
760 if ((!(tmp_fsm_rec = cm_fsm_find(fsm))) /* fsm active in cm ? */
761     || (!(tmp_mbx_rec = cm_mbx_rec_find(*mbxhandle)))) { /* mbxhandle in cm ? */
762     log_status(&cm_activity->cm_read.failure,NULL,NULL);
763     return(E_CM_MBXNOTDECL);
764 }
765 if (!find_mbx_client(tmp_mbx_rec->readerlist,tmp_fsm_rec)) {
```

```
766     log_status(&cm_activity->cm_read.failure,fsm,tmp_mbx_rec->mbxname);
767     return(E_CM_MBXNOTDECL);           /* not a READ/XREAD client of this mbx */
768 }
769 if ((*nrbytes < tmp_mbx_rec->msglength) && (*nrbytes)) {
770     log_status(&cm_activity->cm_read.failure,fsm,tmp_mbx_rec->mbxname);
771     return_status = I_CM_MOREDATA;      /* user wants to read less than available */
772 }
773 else *nrbytes = tmp_mbx_rec->msglength;
774 eprintf(5,"cm_read: passed validation tests\n");
775 if (*nrbytes)                      /* transfer data to mbx */
776     bcpys(tmp_mbx_rec->data,usr_data,*nrbytes);
777     ++(tmp_mbx_rec->readers.nr_accesses);
778     tmp_mbx_rec->readers.who = tmp_fsm_rec;
779     time(&when);                   /* timestamp */
780     tmp_mbx_rec->readers.when = when;
781     tmp_fsm_rec->read.when = when;
782     ++(tmp_fsm_rec->read.nr_accesses);
783     tmp_fsm_rec->read.mbxhandle = *mbxhandle;

785 /* Check if fsm has an update record for this mailbox in his update list. */
786 /* If found, remove it. */
787 del_update_rec(tmp_fsm_rec,*mbxhandle);
788 log_status(&cm_activity->cm_read.success,fsm,tmp_mbx_rec->mbxname);
789 return(return_status);
790 }

793 /* ===== Pass the fsm the update list if one exists, or
794 * | cm_ckmail | pass it NULL if none exists.
795 * =====
796 */
797 Int cm_ckmail(fsm, llist_ptr, nr_entries)
798 char *fsm;
799 struct update_llist **llist_ptr;
800 Int *nr_entries;
801 { fsm_rec *tmp_fsm_rec;

803 eprintf(9,"cm_ckmail: at entry, fsm = %s\n", fsm);
804 if (!cm_activity) cm_init();           /* Initialize cmm structures */
805 if (!tmp_fsm_rec = cm_fsm_find(fsm)) { /* fsm active in cm ? */
806     log_status(&cm_activity->cm_ckmail.success,NULL,NULL);
807     return(E_CM_MBXNOTDECL);
808 }
809 if (tmp_fsm_rec->update_top) {        /* if update llist exists, send it */
810     *llist_ptr = tmp_fsm_rec->update_top;
811     *nr_entries = tmp_fsm_rec->nr_updates;
812     tmp_fsm_rec->update_top = NULL;      /* reset */
813     tmp_fsm_rec->update_bot = NULL;
814     tmp_fsm_rec->nr_updates = 0;
815 }else {
816     *llist_ptr = NULL;
```

```
817     *nr_entries = 0;
818 }
819 log_status(&cm_activity->cm_ckmail.success,fsm,NULL);
820 return(I_CM_OK);
821 }

824 /* ===== Pass the fsm the update list if one exists, or
825 * | cm_disc | pass it NULL if none exists.
826 * ===== */
827 */
828 Int cm_disc(fsm)
829 char *fsm;
830 { Int return_status;
831 byte *data_ptr;
832 fsm_rec *tmp_fsm_rec;
833 mbx_rec *tmp_mbx_rec;
834 struct client_chain *tmp_client_rec;
835 struct mbx_decl_chain *tptr;

837 eprintf(9,"cm_disc: at entry, fsm = %s\n", fsm);
838 if (!cm_activity) cm_init(); /* Initialize cmn structures */
839 if (!tmp_fsm_rec = cm_fsm_find(fsm)) { /* fsm active in cm ? */
840     log_status(&cm_activity->cm_disc.success,NULL,NULL);
841     return(E_CM_MBXNOTDECL);
842 }
843 cm_free_update_list(tmp_fsm_rec->update_top); /* free the update list */

845 while (tmp_fsm_rec->read.mbx_list) { /* process the READ/XREAD list */
846     tmp_mbx_rec = tmp_fsm_rec->read.mbx_list->mbx;
847     eprintf(7,"cm_disc: processing read mbx %s\n",tmp_mbx_rec->mbxname);
848     if (tmp_client_rec = find_mbx_client(tmp_mbx_rec->readerlist,tmp_fsm_rec)) {
849         del_mbx_client(tmp_client_rec, &(tmp_mbx_rec->readerlist), &(tmp_mbx_rec->readers.nr_fsms));
850         if ((!tmp_mbx_rec->readerlist) && (!tmp_mbx_rec->writerlist)) {
851             eprintf(7,"cm_disc: %s has no more clients... freed\n",tmp_mbx_rec->mbxname);
852             --(cm_clients->mbx_active);
853             cm_mbx_del(tmp_mbx_rec); /* delete mbx if it has no more clients */
854         }
855     }else {
856         printf("cm_disc: FATAL ERROR: attempt to undeclare mbx for which fsm isn't client.\n"
857               "Contact the network team.\n");
858         exit(0x40);
859     }
860     tptr = tmp_fsm_rec->read.mbx_list;
861     tmp_fsm_rec->read.mbx_list = tmp_fsm_rec->read.mbx_list->next;
862     free(tptr);
863 }

865 while (tmp_fsm_rec->write.mbx_list) { /* process the WRITE/XWRITE list */
866     tmp_mbx_rec = tmp_fsm_rec->write.mbx_list->mbx;
867     eprintf(7,"cm_disc: processing write mbx %s\n",tmp_mbx_rec->mbxname);
```

```
868     if (tmp_client_rec = find_mbx_client(tmp_mbx_rec->writerlist,tmp_fsm_rec)) {
869         del_mbx_client(tmp_client_rec, &(tmp_mbx_rec->writerlist), &(tmp_mbx_rec->readers.nr_fsms));
870         if ((tmp_mbx_rec->readerlist) && (tmp_mbx_rec->writerlist)) {
871             --(cm_clients->mbx_active);
872             eprintf(7,"cm_disc: %s has no more clients... freed\n",tmp_mbx_rec->mbxname);
873             cm_mbx_del(tmp_mbx_rec);           /* delete mbx if it has no more clients */
874         }
875     }else {
876         printf("cm_disc: FATAL ERROR: attempt to undeclare mbx for which fsm isn't client.\n"
877               "Contact the network team.\n");
878         exit(0x40);
879     }
880     tptr = tmp_fsm_rec->write.mbx_list;
881     tmp_fsm_rec->write.mbx_list = tmp_fsm_rec->write.mbx_list->next;
882     free(tptr);
883 }

885 /* fsm should have no other mailboxes declared, so remove it from cm_fsm_list */
886 if ((tmp_fsm_rec->read.mbx_list) && (tmp_fsm_rec->write.mbx_list)) {
887     --(cm_clients->fsm_active);
888     cm_fsm_del(tmp_fsm_rec);
889 }else{
890     printf("cm_disc: Didn't work... still have read/write clients for this fsm.\n"
891           "Contact the network team.\n");
892     exit(0x40);
893 }
894 log_status(&cm_activity->cm_disc.success,fsm,NULL);
895 return(I_CM_OK);
896 }

899 /* ===== Construct a linked list of mbx names and return
900  * | cm_get_mbx_list | a pointer to it to the caller.  If fsmname == NULL
901  * | ===== | return a list of all mbx's in cm; else return the
902  * | ===== | list of mbx's of the specified fsm and for the
903  * | ===== | specified list_type access.
904 */
905 Int cm_get_mbx_list(fsmname, list_type, mbx_list_ptr, int_ptr)
906 char *fsmname;
907 struct mbx_list_type **mbx_list_ptr;
908 Int *int_ptr;
909 { struct mbx_list_type *tptr;
910   fsm_rec *tmp_fsm_rec;
911   mbx_rec *tmp_mbx_rec;
912   struct mbx_decl_chain *tmp_client_rec;
913   Int return_status;

915   if (!cm_activity) cm_init();           /* initialize cmm structures */
916   *mbx_list_ptr = NULL;                 /* initialize ptr */
917   *int_ptr = 0;                        /* initialize nr entries on list */
918   if (!tmp_mbx_rec = cm_mbx_list)
```

```
919     return(I_CM_OK);                      /* no mbx's in cm */
920     tmp_client_rec = NULL;
921     if (fsmname) {                         /* Is a mbx name provided ? */
922         if (return_status = cm_validate_fsm(fsmname)) {
923             eprintf(9,"cm_get_mbxx_llist: Invalid fsm name '%s'\n",fsmname);
924             log_status(&cm_activity->cm_get_mbxx_llist.failure,NULL,NULL);
925             return(return_status);           /* fsm error */
926         }
927         if (!(tmp_fsm_rec = cm_fsm_find(fsmname))){ /* ck if fsm exists */
928             eprintf(9,"cm_get_mbxx_llist: fsm '%s' doesn't exist in cm\n",fsmname);
929             log_status(&cm_activity->cm_get_mbxx_llist.failure,fsmname,NULL);
930             return(E_CM_FSMNOTINCM);
931         }
932         switch (llist_type) {
933             case 'R':
934             case 'r': tmp_client_rec = tmp_fsm_rec->read.mbx_llist;
935                 eprintf(9,"cm_get_mbxx_llist: for fsm '%s', read llist\n",fsmname);
936                 break;
937             case 'W':
938             case 'w': tmp_client_rec = tmp_fsm_rec->write.mbx_llist;
939                 eprintf(9,"cm_get_mbxx_llist: for fsm '%s', write llist\n",fsmname);
940                 break;
941             default : eprintf(9,"cm_get_mbxx_llist: Invalid llist_type '%c' for fsm '%s'\n",llist_type,fsmname);
942                 log_status(&cm_activity->cm_get_mbxx_llist.failure,fsmname,NULL);
943                 return(E_CM_MBXXACCESS);
944         }
945         tmp_mbx_rec = NULL;                  /* so only tmp_client_rec is used */
946     }else                         /* no fsm name... get all mbx's */
947         eprintf(9,"cm_get_mbxx_llist: no fsm name specified\n");

949     /* build the llist of fsm names */
950     while ((tmp_mbx_rec) || (tmp_client_rec)) {
951         tptr = (struct mbxx_llist_type *) malloc (sizeof(struct mbxx_llist_type));
952         if (tptr) {
953             bclr(tptr,sizeof(struct mbxx_llist_type));
954             if (fsmname)
955                 strcpy(tptr->mbxname,tmp_client_rec->mbx->mbxname);
956             else strcpy(tptr->mbxname,tmp_mbx_rec->mbx->mbxname);
957             tptr->next = *mbxx_llist_ptr;
958             *mbxx_llist_ptr = tptr;
959             if (fsmname)
960                 tmp_client_rec = tmp_client_rec->next;
961             else tmp_mbx_rec = tmp_mbx_rec->next;
962             ++(*int_ptr);
963         }else{                           /* else, out of memory */
964             eprintf(9,"cm_get_mbxx_llist: ran out of memory\n");
965             while (tptr = *mbxx_llist_ptr) {      /* free the llist */
966                 *mbxx_llist_ptr = (*mbxx_llist_ptr)->next;
967                 free(tptr);
968             }
969             *int_ptr = 0;
```

```
970     log_status(&cm_activity->cm_get_mbx_list.failure,fsmname,NULL);
971     return(E_CM_INSUFFMEM);
972 }
973 }
974 log_status(&cm_activity->cm_get_mbx_list.success,fsmname,NULL);
975 return(I_CM_OK);           /* nothing to report */
976 }

979 /* |||||||||||||||||| Construct a linked list of fsm names and return
980 * | cm_get_fsm_list | a pointer to the caller. If mbxname == NULL then
981 * |||||||||||||||||| return all cm clients; else return list_type
982 * | clients of the specified mbx.
983 */
984 Int cm_get_fsm_list(mbxname, list_type, fsm_list_ptr, int_ptr)
985 char *mbxname, list_type;
986 struct fsm_list_type **fsm_list_ptr;
987 Int *int_ptr;
988 { struct fsm_list_type *ptr;
989   fsm_rec *tmp_fsm_rec;
990   mbx_rec *tmp_mbx_rec;
991   struct client_chain *tmp_client_rec;
992   Int return_status;

994   if (!cm_activity) cm_inl();           /* Initialize cmm structures */
995   *fsm_list_ptr = NULL;                /* Initialize ptr */
996   *int_ptr = 0;                       /* Initialize nr entries on list */
997   if (!(tmp_fsm_rec = cm_fsm_list))
998     return(I_CM_OK);                 /* no fsm's in cm */
999   tmp_client_rec = NULL;
1000  if (mbxname) {                     /* Is a mbx name provided ? */
1001    if (return_status = cm_validate_mbx(mbxname,1)) { /* fake the mbxsize */
1002      eprintf(9,"cm_get_fsm_list: Invalid mbxname\n");
1003      log_status(&cm_activity->cm_get_fsm_list.failure,NULL,NULL);
1004      return(return_status);           /* mbx error */
1005    }
1006    if (!(tmp_mbx_rec = cm_mbx_find(mbxname))) { /* ck if mbx exists */
1007      eprintf(9,"cm_get_fsm_list: mbxname doesn't exist in cm\n");
1008      log_status(&cm_activity->cm_get_fsm_list.failure,NULL,mbxname);
1009      return(E_CM_MBXNOTDECL);
1010    }
1011    switch (list_type) {
1012      case 'R':
1013      case 'r': tmp_client_rec = tmp_mbx_rec->readerlist;
1014          eprintf(9,"cm_get_fsm_list: for mbx '%s', readerlist\n",mbxname);
1015          break;
1016      case 'W':
1017      case 'w': tmp_client_rec = tmp_mbx_rec->writerlist;
1018          eprintf(9,"cm_get_fsm_list: for mbx '%s', writerlist\n",mbxname);
1019          break;
1020      default : eprintf(9,"cm_get_fsm_list: invalid list_type '%c' for mbx '%s'\n",list_type,mbxname);
1021    }
1022  }
```

```
1021         log_status(&cm_activity->cm_get_fsm_list.failure, NULL, mbxname);
1022         return(E_CM_MBXACCESS);
1023     }
1024     tmp_fsm_rec = NULL;           /* so only tmp_client_rec is used */
1025 }else                         /* no mbx name... get all fsm's */
1026     eprintf(9, "cm_get_fsm_list: no mbxname specified\n");

1028 /* build the list of fsm names */
1029 while ((tmp_fsm_rec) || (tmp_client_rec)) {
1030     tptr = (struct fsm_list_type *) malloc (sizeof(struct fsm_list_type));
1031     if (tptr) {
1032         bclr(tptr,sizeof(struct fsm_list_type));
1033         if (mbxname)
1034             strcpy(tptr->fsmname,tmp_client_rec->who->fsmname);
1035         else strcpy(tptr->fsmname,tmp_fsm_rec->fsmname);
1036         tptr->next = *fsm_list_ptr;
1037         *fsm_list_ptr = tptr;
1038         if (mbxname)
1039             tmp_client_rec = tmp_client_rec->next;
1040         else tmp_fsm_rec = tmp_fsm_rec->next;
1041         ++(*int_ptr);
1042     }else{                      /* else, out of memory */
1043         eprintf(9, "cm_get_fsm_list: ran out of memory\n");
1044         while (tptr = *fsm_list_ptr) {           /* free the list */
1045             *fsm_list_ptr = (*fsm_list_ptr)->next;
1046             free(tptr);
1047         }
1048         *int_ptr = 0;
1049         log_status(&cm_activity->cm_get_fsm_list.failure, NULL, mbxname);
1050         return(E_CM_INSUFFMEM);
1051     }
1052 }
1053 log_status(&cm_activity->cm_get_fsm_list.success, NULL, mbxname);
1054 return(I_CM_OK);                  /* nothing to report */
1055 }

1058 /* ////////////////////////////// Since the size of the statistics areas is static,
1059 * | cm_get_cm_stats | the user must provide pointers to space in the
1060 * | //////////////////| user data area into which the statistics will be
1061 * *
1062 * *                      copied. This avoids malloc overhead in case the
1063 * *                      user wishes to call this routine multiple times.
1064 */
1064 int cm_get_cm_stats(activity_ptr, client_ptr)
1065 cm_activity_stats *activity_ptr;
1066 cm_client_stats *client_ptr;
1067 {
1068     if (!cm_activity) cm_ini();           /* Initialize cmn structures */
1069     if (CM_GET_STATS) {                 /* collecting stats? */
1070         bcopy(cm_activity, activity_ptr, sizeof(cm_activity_stats));
1071         bcopy(cm_clients, client_ptr, sizeof(cm_client_stats));
```

```

1072     eprintf(9,"cm_get_cm_stats: stats transferred to user area\n");
1073 }else{
1074     eprintf(9,"cm_get_cm_stats: cm is not logging stats\n");
1075     bclr(activity_ptr, sizeof(cm_activity_stats));
1076     bclr(client_ptr, sizeof(cm_client_stats));
1077 }
1078 log_status(&cm_activity->cm_get_cm_stats.success,NULL,NULL);
1079 return(I_CM_OK);                                /* nothing to report */
1080 }

1083 /* ===== Since the size of the statistics area is static,
1084 * | cm_get_fsm_stats | the user must provide a pointer to space in the
1085 * ===== user data area into which the statistics will be
1086 * copied. This avoids malloc overhead in case the
1087 * user wishes to call this routine multiple times.
1088 */
1089 Int cm_get_fsm_stats(fsm, ptr)
1090 char *fsm;
1091 cm_fsm_stats_rec *ptr;
1092 { fsm_rec *tmp_fsm_rec;
1093 mbx_rec *tmp_mbx_rec;
1094 Int return_status;

1096 if (!cm_activity) cm_init();                  /* Initialize cmn structures */
1097 if (return_status = cm_validate_fsm(fsm)) {
1098     eprintf(9,"cm_get_fsm_stats: invalid fsm name '%s'\n",fsm);
1099     log_status(&cm_activity->cm_get_fsm_stats.failure,NULL,NULL);
1100     return(return_status);                      /* fsm error */
1101 }
1102 if (!tmp_fsm_rec = cm_fsm_find(fsm)){        /* ck if fsm exists */
1103     eprintf(9,"cm_get_fsm_stats: fsm '%s' doesn't exist in cm\n",fsm);
1104     log_status(&cm_activity->cm_get_fsm_stats.failure,fsm,NULL);
1105     return(E_CM_FSMNOTINCM);
1106 }
1107 ptr->nr_read_mbx = tmp_fsm_rec->read.nr_mbxes;
1108 ptr->nr_reads     = tmp_fsm_rec->read.nr_accesses;
1109 ptr->read_time    = tmp_fsm_rec->read.when;
1110 if (tmp_fsm_rec->read.mbxhandle != 0) {
1111     tmp_mbx_rec = cm_mbxhandle_find(tmp_fsm_rec->read.mbxhandle);
1112     (tmp_mbx_rec) ? strcpy(ptr->mbx_read,tmp_mbx_rec->mbxname)
1113                  : sprintf(ptr->mbx_read,"%d",tmp_fsm_rec->read.mbxhandle);
1114 }else
1115     sprintf(ptr->mbx_read,<none>);
1116 ptr->nr_write_mbx = tmp_fsm_rec->write.nr_mbxes;
1117 ptr->nr_writes    = tmp_fsm_rec->write.nr_accesses;
1118 ptr->write_time   = tmp_fsm_rec->write.when;
1119 if (tmp_fsm_rec->write.mbxhandle != 0) {
1120     tmp_mbx_rec = cm_mbxhandle_find(tmp_fsm_rec->write.mbxhandle);
1121     (tmp_mbx_rec) ? strcpy(ptr->mbx_write,tmp_mbx_rec->mbxname)
1122                  : sprintf(ptr->mbx_write,"%d",tmp_fsm_rec->write.mbxhandle);

```

```
1123 }else
1124     sprintf(ptr->mbx_write,<none>);
1125     ptr->nr_updates = tmp_fsm_rec->nr_updates;
1126     log_status(&cm_activity->cm_get_fsm_stats.success,fsm,NULL);
1127     return(I_CM_OK);
1128 }

1131 /* |::::::::::::::::::::::::::| Since the size of the statistics area is static,
1132 * | cm_get_mbx_stats | the user must provide a pointer to space in the
1133 * |::::::::::::::::::::::::::| user data area into which the statistics will be
1134 * |
1135 * |                                copied. This avoids malloc overhead in case the
1136 * |                                user wishes to call this routine multiple times.
1137 */
1138 Int cm_get_mbx_stats(mbx, ptr)
1139 char *mbx;
1140 cm_mbx_stats_rec *ptr;
1141 fsm_rec *tmp_fsm_rec;
1142 mbx_rec *tmp_mbx_rec;
1143 int return_status;

1144 if (!cm_activity) cm_init();           /* Initialize cmm structures */
1145 if (return_status = cm_validate_mbx(mbx,1)) { /* fake a mbxsize */
1146     eprintf(9,"cm_get_mbx_stats: invalid mbx name '%s'\n",mbx);
1147     log_status(&cm_activity->cm_get_mbx_stats.failure,NULL,NULL);
1148     return(return_status);           /* fsm error */
1149 }
1150 if (!(tmp_mbx_rec = cm_mbx_find(mbx))){ /* ck if fsm exists */
1151     eprintf(9,"cm_get_mbx_stats: mbx '%s' doesn't exist in cm\n",mbx);
1152     log_status(&cm_activity->cm_get_mbx_stats.failure,mbx,NULL);
1153     return(E_CM_MBXNOTDECL);
1154 }

1155 ptr->handle      = tmp_mbx_rec->handle;
1156 ptr->declaredlength = tmp_mbx_rec->declaredlength;
1157 ptr->msglength    = tmp_mbx_rec->msglength;
1158 ptr->read_fsms    = tmp_mbx_rec->readers.nr_fsms;
1159 ptr->read_accesses = tmp_mbx_rec->readers.nr_accesses;
1160 ptr->read_time     = tmp_mbx_rec->readers.when;
1161 (tmp_mbx_rec->readers.who) ? strcpy(ptr->reader,tmp_mbx_rec->readers.who->fsmname)
1162                               : sprintf(ptr->reader,<none>);
1163 ptr->write_fsms   = tmp_mbx_rec->writers.nr_fsms;
1164 ptr->write_accesses = tmp_mbx_rec->writers.nr_accesses;
1165 ptr->write_time    = tmp_mbx_rec->writers.when;
1166 (tmp_mbx_rec->writers.who) ? strcpy(ptr->writer,tmp_mbx_rec->writers.who->fsmname)
1167                               : sprintf(ptr->writer,<none>);
1168 (tmp_mbx_rec->readerlist->exclusive)
1169     ? strcpy(ptr->xreader,tmp_mbx_rec->readerlist->who->fsmname)
1170     : sprintf(ptr->xreader,<none>);
1171 (tmp_mbx_rec->writerlist->exclusive)
1172     ? strcpy(ptr->xwriter,tmp_mbx_rec->writerlist->who->fsmname)
1173     : sprintf(ptr->xwriter,<none>);
```

```
1174 : sprintf(ptr->xwriter,"<none>");  
1175 log_status(&cm_activity->cm_get_mbx_stats.success,NULL,mbx);  
1176 return(I_CM_OK);  
1177 }
```

1 /* cm_utils.c - contains the following utility routines (in the order
2 shown). These are primarily used by the common memory
3 interface routines (in file cm_funcs.c), but are also
4 accessible to the user.

8 eprintf - displays variable-length arg list using fmt
9 IF level <= DEBUG_LEVEL. As written, is
10 specific to Turbo C.

12 cm_free_mbclient_llst - FREE the mbx's entire client list in preparation for
13 deleting the mbx.

15 cm_free_fsm_mbx_llst - FREE the fsm's entire mbx list in preparation for
16 deleting the fsm entry.

18 cm_free_update_llst - Free all memory allocated to an update list.

20 cm_fsm_fInd - Check if the fsm is on the cm's client list.
21 If found, return ptr; else return NULL.

23 cm_mbclient_fInd - Check if the mbx is already in the cm.
24 If found, return ptr; else return NULL.

26 cm_mbclient_del - Remove a mailbox entry from the cm_mbclient_llst.

28 cm_fsm_del - Remove an fsm entry from the cm_fsm_llst.

30 cm_mbchandle_fInd - Find the mbx record associated w/ this mbchandle.
31 Return NULL if not found.

33 find_mbclient - Check if the fsm is on this mailbox's client list.
34 If found, return ptr to client_chain record.
35 Else, return NULL.

37 add_mbclient - Add an fsm as a mbx client.

39 del_mbclient - Remove an fsm from the respective mbx client list.

41 find_fsm_mbx - Find a mbx (by ptr) on the list kept by the fsm.

43 add_fsm_mbx - Add a mbx to the list kept by the fsm.

45 del_fsm_mbx_entry - Find a mbx (by ptr) on the list kept by the fsm
46 and remove it from the list.

48 cm_validate_fsm - validate the fsm name.

50 cm_validate_mbx - validate the mbx name and size.

```
52 cm_validate_access      - validate the mbx access code.

54 find_update_rec        - Look for an update record in the update list
55 of this fsm that references mbxhandle. If found,
56 return pointer, else return NULL.

58 add_update_rec         - Add an update record to the list for this fsm.

60 del_update_rec         - If this mailbox is on the update list of this
61 fsm (waiting to be read) then delete it.

63 notify_clients          - Notify all clients on the access list of this
64 mbx that the mbx has been updated. If the client
65 already has an update notice for this mbx, do not
66 make another entry.

68 clear_declarer         - If new structures were allocated for this
69 mbx declaration, FREE them and return.

71 cm_ini                  - Initialize the variables and data structures
72 required by the common memory manager. This
73 routine is called the 1st time a user makes a cm
74 call. This call is triggered because cm_activity
75 is NULL. The application can call cm_ini directly in
76 order to force the initialization at a known point.
77 This would be desirable, for example, if the
78 application wanted to change CM_DEBUG_LEVEL or
79 CM_GET_STATS. Function can be called multiple times
80 without detrimental results (ostensibly, this would
81 be for the purpose of displaying the library
82 version number).

84 log_status                - Routine to log cm usage statistics.

86 cm_get_statusname        - Return a pointer to string that gives the status name
87 associated with the status code.

89 */

91 #include <stdio.h>
92 #include "cm_globals.h"

95 #include <stdarg.h> /* as written, this routine is specific to Turbo C */
96 /* |||||||||||||||||| displays variable-length arg list using fmt
97 * |     eprIntf      | IF level <= DEBUG_LEVEL
98 * ||||||||||||||||||
99 */

100 void eprIntf(int level, char *fmt, ...)
101 { va_list argptr;
```

```
103  If (level <= CM_DEBUG_LEVEL) {
104      va_start(argptr, fmt);
105      vprIntf(fmt, argptr);
106      va_end(argptr);
107  }
108 }

111 /* | cm_free_mbx_client_llist | FREE the mbx's entire client list
112 * | cm_free_mbx_client_llist | in preparation for deleting the mbx
113 * | cm_free_mbx_client_llist |
114 */
115 static void cm_free_mbx_client_llist (llist)
116     struct client_chain *llist;
117 { struct client_chain *tptr;

119     while (tptr = llist) {
120         llist = llist->next;
121         free(tptr);
122     }
123 }

126 /* | cm_free_fsm_mbx_llist | FREE the fsm's entire mbx list in
127 * | cm_free_fsm_mbx_llist | preparation for deleting the fsm entry
128 * | cm_free_fsm_mbx_llist |
129 */
130 static void cm_free_fsm_mbx_llist (llist)
131     struct mbx_decl_chain *llist;
132 { struct mbx_decl_chain *tptr;

134     while (tptr = llist) {
135         llist = llist->next;
136         free(tptr);
137     }
138 }

141 /* | cm_free_update_llist | free all memory allocated to an update list
142 * | cm_free_update_llist |
143 * | cm_free_update_llist |
144 */
145 void cm_free_update_llist (llist)
146     struct update_llist *llist;
147 {
148     struct update_llist *tptr;

150     while (tptr = llist) {
151         llist = tptr->next;
152         free(tptr);
153     }
```

```
154 }

157 /* ||||||||||||||||| Check If the fsm is on the cm's client list.
158 * | cm_fsm_find | If found, return ptr; else return NULL.
159 * |||||||||||||||||
160 */
161 static fsm_rec *cm_fsm_find (fsm)
162     register char *fsm;
163 {
164     register fsm_rec *tptr;

166     if (tptr = cm_fsm_llst)
167         while (tptr)
168             if (strcmp(fsm, tptr->fsmname) == 0)
169                 return (tptr);
170             else tptr = tptr->next;
171     return (NULL);
172 }

175 /* ||||||||||||||||| Check If the mbx is already in the cm.
176 * | cm_mbx_find | If found, return ptr; else return NULL.
177 * |||||||||||||||||
178 */
179 static mbx_rec *cm_mbx_find (mbxname)
180     char *mbxname;
181 {
182     register mbx_rec *tptr;

184     if (tptr = cm_mbx_llst)
185         while (tptr)
186             if (strcmp(mbxname, tptr->mbxname) == 0)
187                 return (tptr);
188             else tptr = tptr->next;
189     return (NULL);
190 }

193 /* ||||||||||||||||| Remove a mailbox entry from the cm_mbx_llst
194 * | cm_mbx_del | 
195 * |||||||||||||||||
196 */
197 static int cm_mbx_del(entry)
198     mbx_rec *entry;
199 {
200     mbx_rec *tptr;

202     free(entry->data);                                /* free the data mailbox */
203     if (cm_mbx_llst == entry) {                         /* If it's the 1st entry, adj ptrs */
204         cm_mbx_llst = entry->next;
```

```
205     free(entry);
206     eprintf(6,"cm_mbx_del: mbx entry deleted from top of list\n");
207     return;
208 }
209 tptr = cm_mbx_list;           /* look for this entry's predecessor */
210 while (tptr)
211     if (tptr->next == entry) {           /* found it. adjust the pointer */
212         tptr->next = entry->next;
213         free(entry);
214         eprintf(6,"cm_mbx_del: mbx entry deleted\n");
215         return;
216     } else tptr = tptr->next;
217 printf("%ccm_mbx_del: FATAL ERROR - entry not found! Call network team.");
218 exit(0x40);
219 }

222 /* |||||||||||||||||| Remove an fsm entry from the cm_fsm_list
223 * |      cm_fsm_del   |
224 * ||||||||||||||||||
225 */
226 static int cm_fsm_del(entry)
227 fsm_rec *entry;
228 {
229     fsm_rec *tptr;

231     cm_free_update_list(entry->update_top);
232     if (cm_fsm_list == entry) {           /* If it's the 1st entry, adj ptrs */
233         cm_fsm_list = entry->next;
234         free(entry);
235         eprintf(6,"cm_fsm_del: fsm entry deleted from top of list\n");
236         return;
237     }
238     tptr = cm_fsm_list;           /* look for this entry's predecessor */
239     while (tptr)
240         if (tptr->next == entry) {           /* found it. adjust the pointer */
241             tptr->next = entry->next;
242             free(entry);
243             eprintf(6,"cm_fsm_del: fsm entry deleted\n");
244             return;
245         } else tptr = tptr->next;
246     printf("%ccm_fsm_del: FATAL ERROR - entry not found! Call network team.");
247     exit(0x40);
248 }

251 /* |||||||||||||||||| Find the mbx record associated w/ this mbxhandle.
252 * | cm_mbxhandle_find |  Return NULL if not found.
253 * ||||||||||||||||||
254 */
255 static mbx_rec *cm_mbxhandle_find (mbxhandle)
```

```
256 int mbxhandle;
257 {
258     register mbx_rec *tptr;
259
260     if (tptr = cm_mbx_list)
261         while (tptr)
262             if (tptr->handle == mbxhandle)
263                 return (tptr);
264             else tptr = tptr->next;
265     return (NULL);
266 }
267
268 /* ===== Check if the fsm is on this mailbox's client list.
269  * | find_mbx_client | If found, return ptr to client_chain record.
270  * | ===== | Else, return NULL.
271  * */
272 static struct client_chain *find_mbx_client (list,fsm)
273     struct client_chain *list;
274     register char *fsm;
275 {
276     struct client_chain *tptr;
277
278     if (tptr = list)
279         while (tptr)
280             if (tptr->who == fsm)
281                 return (tptr);
282             else tptr = tptr->next;
283     return(NULL);
284 }
285
286 /* ===== Add an fsm as a mbx client.
287  * | add_mbx_client |
288  * | ===== |
289  * */
290 static int add_mbx_client (fsm, clientlist, exclusive)
291     struct client_chain **clientlist;
292     boolean exclusive;
293     char *fsm;
294 {
295     struct client_chain *tmp_client_rec;
296
297     eprintf(7, "add_mbx_client: attempt to add %s\n", fsm);
298     tmp_client_rec = (struct client_chain *) malloc (sizeof(struct client_chain));
299     if (!tmp_client_rec) return(E_CM_INSUFFMEM); /* verify that malloc worked */
300     tmp_client_rec->who = fsm;
301     tmp_client_rec->next = *clientlist;
302     tmp_client_rec->exclusive = exclusive;
303     *clientlist = tmp_client_rec;
304     return(I_CM_OK);
305 }
```

307 }

```

310 /* |::::::::::::::::::| Remove an fsm from the respective mbx client list.
311 * | del_mbx_client |
312 * |::::::::::::::::::|
313 */
314 static void del_mbx_client (client, clientlist, nr_clients)
315 struct client_chain *client, **clientlist;
316 int *nr_clients;
317 { struct client_chain *tptr;
318
319 eprintf(7,"del_mbx_client: attempt to add %s\n",client->who->fsmname);
320 --(*nr_clients); /* decrement client count */
321 if (client == *clientlist) /* If It's the 1st entry, adj ptrs */
322     *clientlist = client->next;
323     free(client);
324     eprintf(9,"del_mbx_client: fsm client entry deleted from top of list\n");
325     return;
326 }
327 tptr = *clientlist; /* look for this entry's predecessor */
328 while (tptr)
329     if (tptr->next == client) /* found it. adjust the pointer */
330         tptr->next = client->next;
331         free(client);
332         eprintf(9,"del_mbx_client: fsm client entry deleted\n");
333         return;
334     } else tptr = tptr->next;
335 printf("%cdel_mbx_client: FATAL ERROR - client entry not found! Call network team.");
336 exit(0x40);
337 }
338
339 /* |::::::::::::::::::| Find a mbx (by ptr) on the list kept by the fsm
340 * | find_fsm_mbx |
341 * |::::::::::::::::::|
342 */
343 static struct mbx_deci_chain *find_fsm_mbx (list, mbx)
344 struct mbx_deci_chain *list;
345 mbx_rec *mbx;
346 {
347     while (list)
348         if (list->mbx == mbx)
349             return(list);
350         else list = list->next;
351     return(NULL);
352 }
353
354 /* |::::::::::::::::::| Add a mbx to the list kept by the fsm
355 * | add fsm mbx |
356 */

```

```
358 * =====
359 */
360 static int add_fsm_mbx (llist, mbx)
361     struct mbx_decl_chain **llist;
362     mbx_rec *mbx;
363 { struct mbx_decl_chain *tptr;
364
365     if (find_fsm_mbx(*llist,mbx))
366         return(I_CM_OK); /* already exists, don't need to add it */
367     eprintf(7,"add_fsm_mbx: attempt to add %s\n",mbx->mbxname);
368     tptr = (struct mbx_decl_chain *) malloc (sizeof(struct mbx_decl_chain));
369     if (!tptr) return(E_CM_INSUFFMEM); /* verify that malloc worked */
370     tptr->mbx = mbx;
371     tptr->next = *llist;
372     *llist = tptr;
373     return(I_CM_OK);
374 }
375
376 /* ===== Find a mbx (by ptr) on the llist kept by the fsm
377 * | del_fsm_mbx_entry | and remove it from the llist
378 * =====
379 */
380 static void del_fsm_mbx_entry (llist, mbx)
381 struct mbx_decl_chain **llist;
382 mbx_rec *mbx;
383 { struct mbx_decl_chain *tptr, *previous;
384
385     eprintf(7,"del_fsm_mbx_entry: looking for mbx %s\n",mbx->mbxname);
386     tptr = *llist;
387     previous = NULL;
388     while (tptr) {
389         eprintf(9,"del_fsm_mbx_entry: comparing %s\n",tptr->mbx->mbxname);
390         if (tptr->mbx == mbx)
391             if (previous == NULL) {
392                 *llist = tptr->next;
393                 free (tptr);
394                 return;
395             }else {
396                 previous->next = tptr->next;
397                 free (tptr);
398                 return;
399             }
400         }
401     else {
402         previous = tptr;
403         tptr = tptr->next;
404     }
405 }
406     eprintf(7,"del_fsm_mbx_entry: couldn't find mbx on llist\n");
407 }
```

```
410 /* |||||||||||||||||||| validate the fsm name
411 * | cm_validate_fsm |
412 * |||||||||||||||||
413 */
414 Int cm_validate_fsm(fsm)
415 char * fsm;
416 { Int l;

418 l = strlen(fsm);
419 eprintf(9,"cm_validate_fsm: fsm = \"%s\", length = %d characters\n",fsm,l);
420 If ((l > MAXFSMNAMELENGTH) || (l < 1))
421     return (E_CM_FSMNAME);
422 return(I_CM_OK);
423 }

426 /* |||||||||||||||| validate the mbx name and size
427 * | cm_validate_mbx |
428 * |||||||||||||||||
429 */
430 Int cm_validate_mbx(mbxname,mbxsize)
431 char * mbxname;
432 Int mbxsize;
433 {
434     Int l;

436 l = strlen(mbxname);           /* validate mbx name */
437 eprintf(9,"cm_validate_mbx: mbxname = \"%s\", length = %d characters\n",mbxname,l);
438 If ((l > MAXMBXNAMELENGTH) || (l < 1))
439     return (E_CM_MBXNAME);

441 /* validate mbx size - must be greater than 0 bytes long */
442 eprintf(9,"cm_validate_mbx: mbxsize = %d\n",mbxsize);
443 If (mbxsize < 1)
444     return (E_CM_MBXSIZE);
445 return(I_CM_OK);
446 }

448 /* |||||||||||||||| validate the mbx access code
449 * | cm_validate_access |
450 * |||||||||||||||||
451 */
452 Int cm_validate_access(mbxaccess)
453 Int mbxaccess;
454 {
455     Int l;

457 eprintf(9,"cm_validate_access: mbxaccess = (hex) %x\n",mbxaccess);
458 l = mbxaccess & (CM_READ_ACCESS | CM_WRITE_ACCESS | CM_XREAD_ACCESS | CM_XWRITE_ACCESS);
459 If ((mbxaccess = 0) || (l != mbxaccess))
```

```

460     return (E_CM_MBXACCESS);           /* unrecognized code */
461  If ((mbxaccess & (CM_READ_ACCESS | CM_XREAD_ACCESS)) == (CM_READ_ACCESS | CM_XREAD_ACCESS))
462      return (E_CM_MBXACCBOOTH);        /* can't request both concurrently */
463  If ((mbxaccess & (CM_WRITE_ACCESS | CM_XWRITE_ACCESS)) == (CM_WRITE_ACCESS | CM_XWRITE_ACCESS))
464      return (E_CM_MBXACCBOOTH);        /* can't request both concurrently */
465  return(1_CM_OK);
466 }

469 /* |||||||||||||||||| Look for an update record in the update list
470 * | find_update_rec | of this fsm that references mbxhandle.  If found,
471 * |||||||||||||| return pointer, else return NULL.
472 */
473 static struct update_list *find_update_rec(fsm,mbxhandle)
474 Int mbxhandle;
475 fsm_rec *fsm;
476 {
477     struct update_list *tptr;

479     If (tptr = fsm->update_top)           /* a list exists */
480         while (tptr)
481             If (tptr->mbxhandle == mbxhandle) {
482                 sprintf(9,"find_update_rec: FOUND for %s, mbxhandle %d\n",
483                         fsm->fsmname, mbxhandle);
484                 return(tptr);
485             }else tptr = tptr->next;
486         sprintf(9,"find_update_rec: NOT FOUND for %s, mbxhandle %d\n",
487                         fsm->fsmname, mbxhandle);
488         return(NULL);
489     }

492 /* |||||||||||||| Add an update record to the list for this fsm
493 * | add_update_rec | 
494 * ||||||||||||||
495 */
496 static void add_update_rec(fsm,mbxhandle)
497 Int mbxhandle;
498 fsm_rec *fsm;
499 {
500     struct update_list *tptr;

502     tptr = (struct update_list *) malloc (sizeof(struct update_list));
503     If (!tptr) {                         /* verify that malloc worked */
504         printf("%cadd_update_rec: FATAL ERROR malloc failed. Out of Memory?",7);
505         exit(0x40);
506     }
507     tptr->mbxhandle = mbxhandle;
508     tptr->next = NULL;
509     If (fsm->update_top)
510         fsm->update_top->next = tptr;           /* adjust NEXT ptr for old BOTTOM */

```



```
562 static void notify_clients(clientlist, mbxhandle)
563 struct client_chain *clientlist;
564 int mbxhandle;
565 {
566     eprintf(9, "notify_clients: for mbxhandle %d\n", mbxhandle);
567     if (clientlist)                                /* client list exists ? */
568         while (clientlist) {
569             if (!find_update_rec(clientlist->who, mbxhandle)) /* update exist ? */
570                 add_update_rec(clientlist->who, mbxhandle); /* NO - add it */
571             clientlist = clientlist->next;                  /* check next client */
572         }
573     }
574
575
576 /* |||||||||||||||||| If new structures were allocated for this
577 * | clear_declare | mbx declaration, FREE them and return.
578 * ||||||||||||||||||
579 */
580 static int clear_declare (new_fsm, new_mbx, tmp_fsm_rec, tmp_mbx_rec, err_code)
581 boolean new_fsm, new_mbx;
582 fsm_rec *tmp_fsm_rec;
583 mbx_rec *tmp_mbx_rec;
584 int err_code;
585 {
586     if (new_fsm)                                /* FREE the fsm rec */
587         eprintf(6, "clear_declare: free the fsm record\n");
588     if (tmp_fsm_rec->read.mbx_list) {           /* FREE mbx lists first */
589         eprintf(6, "clear_declare: free fsm reader list\n");
590         cm_free_fsm_mbx_list(tmp_fsm_rec->read.mbx_list);
591     }
592     if (tmp_fsm_rec->write.mbx_list) {
593         eprintf(6, "clear_declare: free fsm writer list\n");
594         cm_free_fsm_mbx_list(tmp_fsm_rec->write.mbx_list);
595     }
596     free(tmp_fsm_rec);
597 }
598 if (new_mbx) {
599     eprintf(6, "clear_declare: free the mbx record\n");
600     if (tmp_mbx_rec->data) {                   /* FREE the data mbx, too */
601         eprintf(6, "clear_declare: free data mbx\n");
602         free(tmp_mbx_rec->data);
603     }
604     if (tmp_mbx_rec->readerlist) {            /* FREE the client lists first */
605         eprintf(6, "clear_declare: free mbx readerlist\n");
606         cm_free_mbx_client_list(tmp_mbx_rec->readerlist);
607     }
608     if (tmp_mbx_rec->writerlist) {
609         eprintf(6, "clear_declare: free mbx writerlist\n");
610         cm_free_mbx_client_list(tmp_mbx_rec->writerlist);
611     }
612     free(tmp_mbx_rec);
```

```
613     -cm_mbchandle;
614 }
615 return (err_code);
616 }

619 /* ===== Initialize the variables and data structures required
620  * | cm_ini | by the common memory manager. This routine is called
621  * ===== the 1st time a user makes a cm call. This is triggered
622  *          because cm_activity == NULL;
623 *
624 * The following values, when assigned to CM_DEBUG_LEVEL, result in the
625 * display of debugging data corresponding to that level via routine eprintf.
626 * The debugging levels "build" upon each other. That is, selecting a
627 * debug level also displays those levels below it (ie, those that have
628 * a lower debug level number).
629 *
630 *      0 - no debugging data is displayed
631 *      1 - display the common memory library version (compiled into lib)
632 *      2 -
633 *      3 -
634 *      4 -
635 *      5 - when a mbx is written or read
636 *      6 - when fsm/mbx are added/deleted by cm,
637 *          or when cm_declare failed.
638 *      7 - when an fsm/mbx client is added/deleted
639 *      8 - when an fsm has update records added/deleted
640 *      9 - everything (includes 0-8 and more)
641 *
642 */
643 void cm_ini()
644 /* The initial global value of CM_DEBUG_LEVEL is set to -1 in file cm_globa.h.
645  * If the application program changes the value to be >= 0, then cm_ini will not
646  * change it again. If it is still <0, cm_ini will turn off debugging statements */
647 { if (CM_DEBUG_LEVEL < 0) CM_DEBUG_LEVEL = 0;
648   eprintf(1,"cm_ini: using common memory library version %s\n",CM_VERSION);
649   if (cm_activity)           /* ck if we've been here before */
650     eprintf(9,"cm_ini: function called again AFTER Init already performed\n");
651   else {
652     eprintf(9,"cm_ini: Init in progress\n");
653     cm_activity = (cm_activity_stats *) malloc (sizeof(cm_activity_stats));
654     cm_clients = (cm_client_stats *) malloc (sizeof(cm_client_stats));
655     if ((!cm_activity) || (!cm_clients)) { /* verify that malloc worked */
656       printf("%cm_ini: FATAL ERROR malloc failed. Out of Memory ?",7);
657       exit(0x40);
658     }
659     bclr(cm_activity,sizeof(cm_activity_stats));
660     bclr(cm_clients,sizeof(cm_client_stats));
661   }
662 }
```

```
665 /* ===== Routine to log cm usage statistics.
666 * | log_status |
667 * =====
668 */
669 static void log_status(func,fsm, mbx)
670 base_stats *func;
671 char *fsm, *mbx;
672 {
673     if (CM_GET_STATS) {
674         eprintf(9,"log_status: In progress\n");
675         strcpy(func->fsm,fsm); /* save the fsm name */
676         strcpy(func->mbx,mbx); /* save the mbx name */
677         time(&func->when); /* note the time and date */
678         ++func->nr_times; /* Increment nr times this function called */
679     }
680 }
681
682
683 /* ===== return ptr to string that gives the status name
684 * | cm_get_statusname | associated with the status code.
685 * =====
686 */
687 char *cm_get_statusname (code)
688 int code;
689 { char * status;
690
691     switch (code) {
692         case I_CM_OK : status = "I_CM_OK"; break;
693         case I_CM_MBXACTV : status = "I_CM_MBXACTV"; break;
694         case I_CM_DUPMBX : status = "I_CM_DUPMBX"; break;
695         case I_CM_MOREDATA : status = "I_CM_MOREDATA"; break;
696         /* case E_CM_FATALERR : */
697         case E_CM_INSUFFMEM : status = "E_CM_FATALERR or E_CM_INSUFFMEM";
698             break;
699         case E_CM_MBXERR : status = "E_CM_MBXERR"; break;
700         case E_CM_MBXNOREAD : status = "E_CM_MBXNOREAD"; break;
701         case E_CM_MBXNOXREAD : status = "E_CM_MBXNOXREAD"; break;
702         case E_CM_MBXNOWRITE : status = "E_CM_MBXNOWRITE"; break;
703         case E_CM_MBXNOXWRITE : status = "E_CM_MBXNOXWRITE"; break;
704         case E_CM_MBXSIZE : status = "E_CM_MBXSIZE"; break;
705         case E_CM_MBXACCESS : status = "E_CM_MBXACCESS"; break;
706         case E_CM_MBXACCBOTH : status = "E_CM_MBXACCBOTH"; break;
707         case E_CM_MBXNAME : status = "E_CM_MBXNAME"; break;
708         case E_CM_MBXNOTDECL : status = "E_CM_MBXNOTDECL"; break;
709         case E_CM_FSMERR : status = "E_CM_FSMERR"; break;
710         case E_CM_FSMNAME : status = "E_CM_FSMNAME"; break;
711         case E_CM_FSMNOTINCM : status = "E_CM_FSMNOTINCM"; break;
712         default : status = "unknown code"; break;
713     }
714
715     return (status);
716 }
```

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715 }


```
1 /* sfuns.c - miscellaneous functions
2 *
3 * asciidump(b,n) dump n bytes starting at b, in ascii.
4 *
5 * bclr(b,n)      clears (sets = 0) n bytes starting at address b.
6 *
7 * bcopy(s,t,n)   copies n bytes from s to t.
8 *
9 * binstr(n,str) convert byte to binary representation in string str.
10 *
11 * cvtup(c)       converts a lowercase ascii string to upper.
12 *
13 * hexdump(b,n)  dump n bytes starting at b, in hex.
14 *
15 * is_ascii(c)    returns TRUE if c is a printable ascii character.
16 *
17 * pause(fmt,...) routine works like "printf", and accepts a variable-
18 * length arg list after fmt. After it displays the
19 * user-specified data, it displays "Press any character
20 * to continue...", waits for a kbd entry, then
21 * outputs <CRLF> and returns.
22 */

25 /*||||||||||||||||||| display n bytes in ascii, starting w/ address b.
26 * | asciidump      | 'l' columns per line.
27 * |||||||||||||||||
28 */
29 void asciidump(b,n)
30     register char *b;
31     register int n;
32 { register int l = 20;

34     while (n--) {
35         if ((*b > 0x1f) && (*b < 0x7f))
36             printf ("%3c", *b++);
37         else printf (" %02x", 0xff & *b++);
38         if (--l == 0) {
39             printf ("\n");
40             l = 20;
41         }
42     }
43     if (l < 20) printf ("\n");
44 }

47 /*||||||||||||||||||| set n bytes = 0, starting w/ address b.
48 * | bclr           |
49 * |||||||||||||||||
50 */
51 void bclr(b, n)
```

```
52     register char *b;
53     register int n;
54 {
55     while (n--) *b++ = 0;
56 }

59 /* ||||||||||||||||||      copy n bytes from s to t
60 * |      bcopy      |
61 * ||||||||||||||||||
62 */
63 void bcopy(s, t, n)
64     register char *s, *t;
65     register int n;
66 {
67     while (n--) *t++ = *s++;
68 }

71 /* ||||||||||||||      convert byte n to its binary representation
72 * |      binstr      |
73 * ||||||||||||||
74 *
75 *
76 *
77 */
78 char * binstr(n,str)
79 unsigned char n;
80 char *str;                                /* must be at least 9 bytes long */
81 { int i;

83     for (i = 0; i < 8; i++) {
84         str[i] = (n & 0x80) ? '1' : '0';
85         n = n << 1;
86     }
87     str[8] = 0;                                /* null - terminate the string */
88     return (str);
89 }

92 /* ||||||||||||||      convert a lowercase ASCII string to uppercase.
93 * |      cvtup      |
94 * ||||||||||||||
95 */
96 void cvtup (s)
97     register char *s;
98 {
99     while (*s) {
100         if (*s >= 'a' && *s <= 'z') *s -= 'a' - 'A';
101         s++;
102     }

```

103 }

```
106 /* |::::::::::::::::::::::::::| display n bytes in hex, starting w/ address b.
107 * | hexdump | 'l' columns per line.
108 * |::::::::::::::::::::::::::|
109 */
110 void hexdump(b,n)
111     char *b;
112     int n;
113 { int group, l, max;
114     int max_per_line = 15;
115     int group_size = 5;
116
117     while (n) {
118         max = max_per_line;
119         if (max > n) max = n;
120         group = group_size;
121         for (l=1; l <= max; l++) {
122             printf (" %02x",*b++ & 0xff);
123             if (l==group) {
124                 printf (" ");
125                 group = group_size;
126             }
127         }
128         if (l <= max_per_line)
129             for (l = l; l <= max_per_line; l++) {
130                 printf(" "); /* space over to ascii field */
131                 if (l==group) {
132                     printf (" ");
133                     group = group_size;
134                 }
135             }
136         printf(" "); /* gap between hex and ascii areas */
137         b -= max; /* point to start of section */
138         for (l=1; l <= max; l++) {
139             printf("%c", is_ascii(*b) ? *b : 5);
140             b++;
141             if (l==group) {
142                 printf ("%c",176);
143                 group = group_size;
144             }
145         }
146         printf("\n");
147         n -= max;
148     }
149 }
150
151 /* |::::::::::::::::::::::::::| returns TRUE if c is a printable ascii character.
152 * | is_ascii |
```

```
154 * =====
155 */
156 int Is_ascii(c)
157     char c;
158 {
159     if ((c > 0x1f) && (c < 0x7f))
160         return (1);
161     return (0);
162 }

164 #include <stdio.h>           /* as written, this routine is specific to Turbo C */
165 #include <stdarg.h>
166 /* =====
167 * | pause |
168 * =====
169 */
170 void pause(char *fmt, ...)
171 { va_list argptr;

173     va_start(argptr, fmt);
174     vprintf(fmt,argptr);
175     va_end(argptr);
176     printf("Press any character to continue... ");
177     getch();
178     printf("\n");
179 }
```

GLOSSARY

AMRF - acronym, refers to the Automated Manufacturing Research Facility of the National Bureau of Standards.

cm - abbreviation for "common memory". See also common memory.

cmm - abbreviation for "common memory manager". See also common memory manager.

common memory - a term used to generically identify both local and global common memory. See also local common memory and global common memory.

common memory manager - collection of functions that establish the mailboxes on the local host and manage associated data structures to assure and provide proper access to them. The common memory manager also gathers and provides utilization statistics.

DOS - abbreviation for "Disk Operating System", a single-user operating system developed by Microsoft Corporation for use with the IBM PC class of machine. It is marketed by various companies under the names PC DOS and MS DOS.

finite state machine - a term assigned to user application programs that have a finite number of clearly defined processing states. Examples of such states are: data acquisition, data reduction, and data reporting. In the context of this documentation, "fsm" is intended to mean "user application program".

fsm - abbreviation for "finite state machine". See also finite state machine.

global common memory - two or more local common memories combine to form a global common memory. This is accomplished with the introduction of a network interface process (NIP) at each computer system that has a local common memory. The NIP becomes another client of its local common memory with all implied READ/WRITE privileges. NIPs exchange common memory mailgrams with each other using network services, propagating these mailgrams globally and creating the global common memory.

local common memory - a contiguous area of physical memory accessible to two or more distinct processes within a single computer system. This physical memory is

divided into a collection of READ and WRITE mailbox areas. The data in these areas, called mailgrams, is available to all other applications on the computer system.

mailbox - a contiguous area of high-speed memory assigned and managed by the common memory manager. Messages (called mailgrams) can be placed into a mailbox by one or more writer applications and copied from the mailbox by one or more reader applications.

mailgram - term used to identify the contents of a mailbox. That is, a collection of contiguous bytes stored in a mailbox. The data representation (binary, ASCII, etc.) for mailgrams is determined by the application process.

mbx - abbreviation for "mailbox". See also mailbox.

NBS - acronym, refers to the National Bureau of Standards, located in Gaithersburg, Maryland.

PC - abbreviation for "personal computer". This term is used to identify all classes of personal computers that are compatible with the IBM PC and use the DOS operating system.

process - term used to identify a user application program. It is used interchangably with "fsm".

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